

# On Morita Equivalence of Categories

Bertalan Pécsi

october 2007.

## Heteromorphisms

**Def. 1.1.** Category  $\mathbb{H}$  is a *bridge*  $\mathbb{A} \rightleftharpoons \mathbb{B}$ , if

- $\mathbb{A}, \mathbb{B}$  are *disjoint, full* subcategories of  $\mathbb{H}$ ,
- $\text{Ob}\mathbb{A} \cup \text{Ob}\mathbb{B} = \text{Ob}\mathbb{H}$ .

$\mathbb{H}$  is *directed bridge*  $\mathbb{A} \Rightarrow \mathbb{B}$ , if moreover

- $(b \dashv a)_{\mathbb{H}} = \emptyset$  for all  $a \in \text{Ob}\mathbb{A}, b \in \text{Ob}\mathbb{B}$ .

Let  $\mathbb{H}, \mathbb{K}$  be bridges  $\mathbb{A} \rightleftharpoons \mathbb{B}$ .

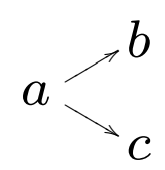
A functor  $T : \mathbb{H} \rightarrow \mathbb{K}$  is a *bridge morphism*, if

- $T \upharpoonright_{\mathbb{A}} = \text{id}_{\mathbb{A}}$  and  $T \upharpoonright_{\mathbb{B}} = \text{id}_{\mathbb{B}}$ .

**Prop. 1.2.** Directed bridges  $\mathbb{A} \Rightarrow \mathbb{B}$  are just the *profunctors* (i.e. functors  $\mathbb{A}^{op} \times \mathbb{B} \rightarrow \text{Set}$ ), corresponding (*bridge morphisms*  $\iff$  *nat. transformations*).

## Examples:

Let  $\mathbb{A}$  be arbitrary, and  $\mathbb{B} \leq \mathbb{A}$  be a full subcat.

bridge	heteromorphisms
$\text{Set} \Rightarrow \text{Grp}$	functions $S \rightarrow G$
$\mathbb{A}b \times \mathbb{A}b \Rightarrow \mathbb{A}b$	bilinear morphisms $A \times B \rightarrow C$
$\mathbb{A} \Rightarrow \mathbb{A} \times \mathbb{A}$	cones 
$\text{Set} \Rightarrow \text{Set}^{op}$	relations between $A$ and $B$
$\mathbb{B} \rightleftarrows \mathbb{A}$	copies of arrows $b \rightarrow a$ & $a \rightarrow b$

A bridge between **monoids** is just a category with 2 objects.

## Profunctors

**Theorem 1.3.** Let  $\mathbb{L} : \mathbb{A} \rightrightarrows \mathbb{B}$  be a profunctor.

•  $\mathbb{L}$  is induced by a functor  $\mathbb{A} \rightarrow \mathbb{B}$

iff  $\mathbb{B} \leq \mathbb{L}$  is reflective.

•  $\mathbb{L}$  is induced by a functor  $\mathbb{B} \rightarrow \mathbb{A}$

iff  $\mathbb{A} \leq \mathbb{L}$  is coreflective.

•  $\mathbb{L}$  is induced by an adjunction  $\mathbb{A} \dashv \mathbb{B}$

iff  $\mathbb{B} \leq \mathbb{L}$  is reflective and  $\mathbb{A} \leq \mathbb{L}$  is coreflective.

**Def. 1.4.** *Composition* of  $\mathbb{F} : \mathbb{A} \rightrightarrows \mathbb{B}$  and  $\mathbb{G} : \mathbb{B} \rightrightarrows \mathbb{C}$  :

$$(a \mid c)_{\mathbb{F}.\mathbb{G}} := \{ \langle f, g \rangle \mid a \xrightarrow{f} b \xrightarrow{g} c, b \in \text{Ob}\mathbb{B} \} / \sim$$

where  $\langle f\beta, g \rangle \sim \langle f, \beta g \rangle$  for  $\beta \in \text{Mor}\mathbb{B}$ .

$\rightsquigarrow$  the bicategory  $\mathbb{Prof}$

(of categories, profunctors, bridge morphisms).

**Note.** A bridge  $\mathbb{H} : \mathbb{A} \rightleftarrows \mathbb{B}$  is determined by its

*parts*  $\mathbb{H}^> : \mathbb{A} \rightrightarrows \mathbb{B}$ ,  $\mathbb{H}^< : \mathbb{B} \rightrightarrows \mathbb{A}$ ,

and compositions  $\mathbb{H}^> \cdot \mathbb{H}^< \rightarrow \mathbb{A}$ ,  $\mathbb{H}^< \cdot \mathbb{H}^> \rightarrow \mathbb{B}$ .

## Equivalences

**Def. 2.1.**  $\mathbb{H} : \mathbb{A} \rightleftarrows \mathbb{B}$  is an *equivalence bridge*, if  
 $\forall a \in \text{Ob}\mathbb{A} \exists b \in \text{Ob}\mathbb{B} : a \cong b$  in  $\mathbb{H}$ , and  
 $\forall b \in \text{Ob}\mathbb{B} \exists a \in \text{Ob}\mathbb{A} : a \cong b$  in  $\mathbb{H}$ .

**Theorem 2.2.**  $\mathbb{A} \simeq \mathbb{B}$  iff  $\exists \mathbb{A} \rightleftarrows \mathbb{B}$  equiv. bridge.

**Note.** Axiom of choice is used in constructing a functor from an equivalence bridge.

(cf. Makkai: "Avoiding the Axiom of Choice...")

**Def. 2.3.**  $\mathbb{M} : \mathbb{A} \rightleftarrows \mathbb{B}$  is a *Morita bridge*, if every morphism is composition of heteromorphisms.

**Def. 2.4** (*Idempotent completion*).

- $\text{Ob}(\mathbb{A}^{id}) := \{e \in \text{Mor} \mathbb{A} \mid e^2 = e\},$
- $(e \mid f)_{\mathbb{A}^{id}} := \{\alpha \mid e\alpha f = \alpha\}.$

**Theorem 2.5.** The followings are equivalent:

a) There are profunctors  $\underset{\mathbb{A} \Rightarrow \mathbb{B}}{\mathbb{F}}, \underset{\mathbb{B} \Rightarrow \mathbb{A}}{\mathbb{G}},$

such that  $\mathbb{F} \cdot \mathbb{G} \cong \mathbb{A}$  and  $\mathbb{G} \cdot \mathbb{F} \cong \mathbb{B}.$

b) There is a Morita bridge  $\mathbb{M} : \mathbb{A} \rightleftarrows \mathbb{B}.$

c)  $\mathbb{A}^{id} \simeq \mathbb{B}^{id}.$

**Theorem 2.5.** The followings are equivalent:

a) There are profunctors  $\mathbb{F} : \mathbb{A} \Rightarrow \mathbb{B}$ ,  $\mathbb{G} : \mathbb{B} \Rightarrow \mathbb{A}$ ,

such that  $\mathbb{F} \cdot \mathbb{G} \cong \mathbb{A}$  and  $\mathbb{G} \cdot \mathbb{F} \cong \mathbb{B}$ .

b) There is a Morita bridge  $\mathbb{M} : \mathbb{A} \rightleftharpoons \mathbb{B}$ .

c)  $\mathbb{A}^{id} \simeq \mathbb{B}^{id}$ .

*Proof.*

a)  $\Rightarrow$  b):  $\mathbb{M} := \mathbb{F} \cup \mathbb{G}$  can be made a bridge by

**Lemma 2.6.** Let  $f : A \rightarrow B$ ,  $g : B \rightarrow A$  be an equivalence in a bicategory, with isomorphisms  $\varphi : fg \rightarrow 1_A$  and  $\psi : gf \rightarrow 1_B$ .

Then  $\exists \psi' : f \cdot \psi' = \varphi \cdot f$  and  $\psi' \cdot g = g \cdot \varphi$ .  
 $gf \rightarrow 1_B$        $fgf \rightarrow f$        $gfg \rightarrow g$

b)  $\Rightarrow$  c): Consider  $\mathbb{M}^{id}$ .

c)  $\Rightarrow$  b): Let  $\mathbb{H} : \mathbb{A}^{id} \rightleftharpoons \mathbb{B}^{id}$  be an equivalence bridge. Set  $\mathbb{M} := \mathbb{H} \downarrow_{\mathbb{A} \cup \mathbb{B}}$ . It is a Morita bridge.

b)  $\Rightarrow$  a): Consider the parts of  $\mathbb{M}$ :

$\mathbb{F} := \mathbb{M}^>$  and  $\mathbb{G} := \mathbb{M}^<$ . They give an equivalence by

**Lemma 2.7.** Let  $\mathbb{K} : \mathbb{A} \rightleftharpoons \mathbb{B}$  be a bridge, with surjective composition  $\chi : \mathbb{K}^> \cdot \mathbb{K}^< \rightarrow \mathbb{A}$ . Then  $\chi$  is isomorphism.

□

## A next level in abstraction

**Def. 3.1.** In a bicategory,

$\langle \begin{matrix} f & g \\ A \rightarrow B & B \rightarrow A \end{matrix}, \begin{matrix} \varphi & \psi \\ fg \rightarrow 1_A & gf \rightarrow 1_B \end{matrix} \rangle$  is a *bridge*, if

$$\cdot \begin{matrix} f \cdot \psi = \varphi \cdot f \\ fgf \rightarrow f \end{matrix} \quad \text{and} \quad \begin{matrix} \psi \cdot g = g \cdot \varphi \\ gfg \rightarrow g \end{matrix}$$

### Examples:

Let  $\mathbb{B}imod$  be the bicategory of (rings, bimodules).

Then  $\begin{matrix} R(R^{1 \times n}) \\ R^{n \times n} \end{matrix}$  with  $\begin{matrix} R^{n \times n}(R^{n \times 1}) \\ R \end{matrix}$

is a bridge in  $\mathbb{B}imod$ .

By lemma 2.6 every *equivalence* in a bicategory can be made a *bridge*.

Note that lemma 2.7 also holds in  $\mathbb{B}imod$ .

### Question.

Search for more *examples* of bridges.



Bertalan Pécsi  
ELTE University, Budapest

[aladar@renyi.hu](mailto:aladar@renyi.hu)  
<http://www.renyi.hu/~aladar>