

Representability Relative to a Doctrine

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Many classical notions in CT are representability notions

- ① Limits = representability of cone functors.
- ② Adjunctions = representability of $\mathcal{A}(L-, A)$.
- ③ Monoidal structures = representability of promonoidal structures.
- ④ ...

Weakened representability \Rightarrow weakened notions.

Example: Weak Limits

$F : \mathcal{A}^{op} \rightarrow \text{Set}$ is **weakly representable** if there is an epimorphism $\mathcal{A}(-, A) \rightarrow F$ for some A .

For a **weak limit** of $D : \mathcal{D} \rightarrow \mathcal{A}$, choose F to be the **cone functor**

$\text{Cone}(D) : X \mapsto$ the set of D -cones with vertex X

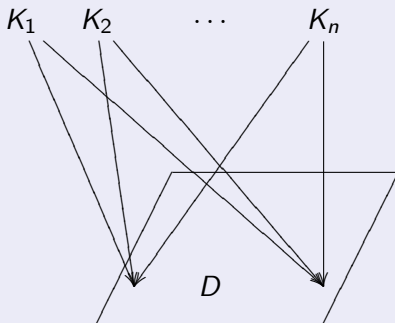
Hence an epimorphism

$$\mathcal{A}(-, A) \rightarrow \text{Cone}(D)$$

meaning: there is a distinguished cone for D with vertex A through which any other factors (not necessarily uniquely).

Example: A Finite Plurilimit of a Finite Diagram D (P.Karazeris, J.Rosický, JV, JPAA, 2005)

There exists a **finite set of distinguished cones**



for D through which any other factors **uniquely up to a zig-zag**:

$$\text{Cone}(D) \cong \text{colim}_i \mathcal{A}(-, K_i)$$

Plurirepresentability

$F : \mathcal{A}^{op} \rightarrow \text{Set}$ is **plurirepresentable** if there is a natural isomorphism

$$F \cong \text{colim}_i \mathcal{A}(-, K_i)$$

for some finite diagram $K : \mathcal{K} \rightarrow \mathcal{A}$.

Many Other Such Notions

- 1 Multirepresentability:

$$F \cong \coprod_i \mathcal{A}(-, K_i)$$

for some finite discrete diagram $K : \mathcal{K} \rightarrow \mathcal{A}$.

- 2 ...

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Common Important Features

Given \mathcal{A} , form a category $\mathbb{C}(\mathcal{A})$ that

- ① contains \mathcal{A}
- ② is contained in $[\mathcal{A}^{op}, \text{Set}]$

$$\mathcal{A} \cdots \longrightarrow \mathbb{C}(\mathcal{A}) \cdots \longrightarrow [\mathcal{A}^{op}, \text{Set}]$$

to measure the “degree of representability” of (say) cone functors for $D : \mathcal{D} \rightarrow \mathcal{A}$.

The Goals of the Talk

- 1 To give a **uniform environment** where weak notions can be studied.
- 2 To show that **weak notions** abound: in domain theory, in general algebra, ...
- 3 **Weak limits** have connections to **honest limits in free cocompletions**.

For this, it is convenient to work in **enriched categories**.
In fact, it does not make the reasoning any harder.

Motivation: Representability as a Factorization

Let \mathcal{I} be the **one-morphism category**.

For $F : \mathcal{A}^{op} \rightarrow \mathbf{Set}$, denote by $\lceil F \rceil : \mathcal{I} \rightarrow [\mathcal{A}^{op}, \mathbf{Set}]$ the **name** of F , i.e., $\lceil F \rceil(*) = F$.

F is **representable** if there is a **factorization**

$$\begin{array}{ccc}
 \mathcal{I} & & \\
 \vdots & \searrow \lceil F \rceil & \\
 \mathcal{A} & \xrightarrow{Y_{\mathcal{A}}} & [\mathcal{A}^{op}, \mathbf{Set}]
 \end{array}$$

to within an isomorphism.

The Weakening Strategy

In the diagram

$$\begin{array}{ccc}
 \mathcal{I} & & \\
 \vdots & \searrow \lceil F \rceil & \\
 \mathcal{A} & \xrightarrow{Y_{\mathcal{A}}} & [\mathcal{A}^{op}, \text{Set}]
 \end{array}$$

replace

- 1 \mathcal{A} by $\mathbb{C}(\mathcal{A})$ (with a **fully faithful** $\gamma_{\mathcal{A}} : \mathcal{A} \rightarrow \mathbb{C}(\mathcal{A})$).
- 2 $Y_{\mathcal{A}} : \mathcal{A} \rightarrow [\mathcal{A}^{op}, \text{Set}]$ by a **fully faithful**
 $\widetilde{\gamma}_{\mathcal{A}} : \mathbb{C}(\mathcal{A}) \rightarrow [\mathcal{A}^{op}, \text{Set}], X \mapsto \mathbb{C}(\mathcal{A})(\gamma_{\mathcal{A}} -, X)$.
- 3 \mathcal{I} by a **general indexing category** \mathcal{M} .
- 4 $\lceil F \rceil$ by a **general functor** $G : \mathcal{M} \rightarrow [\mathcal{A}^{op}, \text{Set}]$.
- 5 Set by a well-behaved base **monoidal category** \mathcal{V} .

Definition

A **doctrine** on \mathcal{V} -CAT is a pair (\mathbb{C}, γ) consisting of a pseudofunctor $\mathbb{C} : \mathcal{V}\text{-CAT} \rightarrow \mathcal{V}\text{-CAT}$ and a pseudonatural $\gamma : \text{Id} \rightarrow \mathbb{C}$ such that for each \mathcal{A} :

- ① $\gamma_{\mathcal{A}} : \mathcal{A} \rightarrow \mathbb{C}(\mathcal{A})$ is **fully faithful**.
- ② $\tilde{\gamma}_{\mathcal{A}} : \mathbb{C}(\mathcal{A}) \rightarrow [\mathcal{A}^{op}, \mathcal{V}]$, $X \mapsto \mathbb{C}(\mathcal{A})(\gamma_{\mathcal{A}} -, X)$, is **fully faithful** (i.e., $\gamma_{\mathcal{A}}$ is **dense**).

Examples of Doctrines

- ① (Id, id) .
- ② Any free cocompletion $\gamma_{\mathcal{A}} : \mathcal{A} \rightarrow \mathbb{C}(\mathcal{A})$ under a class \mathbb{C} of colimits.
- ③ $\mathcal{V} = \text{Set}$, $\mathbb{Q}(\mathcal{A}) =$ quotients of representables.

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Definition

A functor $G : \mathcal{M} \rightarrow [\mathcal{A}^{op}, \mathcal{V}]$ is **representable relative to** (\mathbb{C}, γ) , if there exists a factorization

$$\begin{array}{ccc}
 \mathcal{M} & & \\
 \text{rep}(G) \downarrow & \searrow G & \\
 \mathbb{C}(\mathcal{A}) & \xrightarrow{\widetilde{\gamma}_{\mathcal{A}}} & [\mathcal{A}^{op}, \mathcal{V}]
 \end{array}$$

to within an isomorphism. The isomorphism $\alpha : G \rightarrow \widetilde{\gamma}_{\mathcal{A}} \cdot \text{rep}(G)$ is called the **representation**.

This means:

$$\alpha_{M,A} : (GM)(A) \cong \mathbb{C}(\mathcal{A})(\gamma_{\mathcal{A}} A, \text{rep}(G)M)$$

holds **naturally** in M and A .

Examples of Representability Relative to (\mathbb{C}, γ)

$\mathcal{V} = \text{Set}$, $\mathcal{M} = \mathcal{I}$, $G = \lceil F \rceil$ for $F : \mathcal{A}^{op} \rightarrow \mathcal{V}$.

- 1 Representability relative to (Id, id) is the **usual representability**.
- 2 Representability relative to (\mathbb{Q}, γ) — the doctrine of quotients of representables — is the **weak representability**.
- 3 Representability relative to the doctrine of cocompletions under finite colimits is the **plurirepresentability**.
- 4 Etc...

The case $G = \tilde{F} : \mathcal{M} \rightarrow [\mathcal{A}^{op}, \mathcal{V}]$ for $F : \mathcal{A} \rightarrow \mathcal{M}$

Representability of \tilde{F} relative to (\mathbb{C}, γ) is an isomorphism

$$\alpha_{M,A} : \mathbb{C}(\mathcal{A})(\gamma_{\mathcal{A}} A, \text{rep}(\tilde{F})M) \cong (\tilde{F}M)(A) = \mathcal{M}(FA, M)$$

natural in M and A .

This means: F is a **left adjoint along** $\gamma_{\mathcal{A}}$, $F \dashv_{\gamma_{\mathcal{A}}} \text{rep}(\tilde{F})$, studied by Max Kelly, Walter Tholen, ...

$$\begin{array}{ccc}
 \mathcal{M} & & \\
 \text{rep}(\tilde{F}) \downarrow \text{dotted} & \searrow \tilde{F} & \\
 \mathbb{C}(\mathcal{A}) & \xrightarrow{\gamma_{\mathcal{A}}} & [\mathcal{A}^{op}, \mathcal{V}]
 \end{array}$$

$$\begin{array}{ccc}
 & \mathcal{M} & \\
 F \nearrow & & \searrow \text{rep}(\tilde{F}) \\
 \mathcal{A} & \xrightarrow{\gamma_{\mathcal{A}}} & \mathbb{C}(\mathcal{A})
 \end{array}$$

Representability Relative to (\mathbb{C}, γ) of \tilde{F} , for $F : \mathcal{A} \rightarrow \mathcal{M}$

- ① Any \mathcal{V} , (\mathbb{C}, γ) =identity doctrine. Then $F \dashv_{\text{id}} \text{rep}(\tilde{F})$ is an **honest adjunction**:

$$\alpha_{M,A} : \mathcal{A}(\text{id}_{\mathcal{A}} A, \text{rep}(\tilde{F})M) \cong (\tilde{F}M)(A) = \mathcal{M}(FA, M)$$

natural in M and A .

- ② $\mathcal{V}=\text{Set}$, (\mathbb{C}, γ) =free cocompletion under small colimits. Then $F \dashv_{\gamma_{\mathcal{A}}} \text{rep}(\tilde{F})$ asserts a **solution set condition**.

Example: The “Most General” Gabriel-Ulmer Duality

Suppose (\mathbb{C}, γ) is a fixed **free-cocompletion** doctrine.
The obvious correspondence

$$\begin{aligned} G & : \mathcal{A} \rightarrow [\mathcal{B}^{\text{op}}, \mathcal{V}] \\ \text{rep}(G) & : \mathcal{A} \rightarrow \mathbb{C}(\mathcal{B}) \\ \text{Lan}_{\gamma_{\mathcal{A}}}(\text{rep}(G)) & : \mathbb{C}(\mathcal{A}) \rightarrow \mathbb{C}(\mathcal{B}) \end{aligned}$$

for \mathcal{A}, \mathcal{B} small and G representable relative to \mathbb{C} is a part of a **biequivalence** between **certain profunctors** and “ **\mathbb{C} -accessible functors**”.

This biequivalence **restricts to duality** of “theory morphisms” and “ **\mathbb{C} -accessible right adjoints**”.

Recollection of (Weighted) Limits

For a **diagram** $D : \mathcal{D} \rightarrow \mathcal{A}$ together with a **weight**
 $W : \mathcal{M} \rightarrow [\mathcal{D}, \mathcal{V}]^{op}$ form a **cylinder functor**

$$\text{Cyl}(W, D) : \mathcal{M} \rightarrow [\mathcal{A}^{op}, \mathcal{V}], M \mapsto [\mathcal{D}, \mathcal{V}]^{op}(\widehat{D}-, WM)$$

where $\widehat{D} : A \mapsto \mathcal{A}(A, D-)$.

A **limit** of D weighted by W is a representation $\{W, D\} : \mathcal{M} \rightarrow \mathcal{A}$
of $\text{Cyl}(W, D)$, i.e., we have a diagram

$$\begin{array}{ccc} \mathcal{M} & & \\ \downarrow \{W, D\} & \searrow \text{Cyl}(W, D) & \\ \mathcal{A} & \xrightarrow{Y_{\mathcal{A}}} & [\mathcal{A}^{op}, \mathcal{V}] \end{array}$$

commutative to within an isomorphism.

(Weighted) Limits Relative to a Doctrine

A **limit relative to (\mathbb{C}, γ)** of $D : \mathcal{D} \rightarrow \mathcal{A}$ weighted by $W : \mathcal{M} \rightarrow [\mathcal{D}, \mathcal{V}]^{op}$ is a representation $\{W, D\}_{(\mathbb{C}, \gamma)} : \mathcal{M} \rightarrow \mathbb{C}(\mathcal{A})$ of $\text{Cyl}(W, D)$ relative to (\mathbb{C}, γ) , i.e., we have a diagram

$$\begin{array}{ccc}
 \mathcal{M} & & \\
 \downarrow \{W, D\}_{(\mathbb{C}, \gamma)} & \searrow \text{Cyl}(W, D) & \\
 \mathbb{C}(\mathcal{A}) & \xrightarrow[\gamma_{\mathcal{A}}]{} & [\mathcal{A}^{op}, \mathcal{V}]
 \end{array}$$

commutative to within an isomorphism.

Or, in elementary terms:

$$\begin{aligned}
 \mathbb{C}(\mathcal{A})(\gamma_{\mathcal{A}} A, \{W, D\}_{(\mathbb{C}, \gamma)} M) &\cong \text{Cyl}(W, D)(M)(A) \\
 &= [\mathcal{D}, \mathcal{V}](WM, \hat{D}A)
 \end{aligned}$$

naturally in M and A .

(Weighted) Limits of Some Class Relative to a Doctrine

Fix a **limit doctrine** (\mathbb{L}, λ) .

That is: for each \mathcal{A} , $\lambda_{\mathcal{A}} : \mathcal{A} \rightarrow \mathbb{L}(\mathcal{A})$ is a free completion under a **class of limits**.^a

- 1 A weight $W : \mathcal{M} \rightarrow [\mathcal{D}, \mathcal{V}]^{op}$, \mathcal{M}, \mathcal{D} small, is an **\mathbb{L} -weight**, if it factors through $\widehat{\lambda}_{\mathcal{D}} : \mathbb{L}(\mathcal{D}) \rightarrow [\mathcal{D}, \mathcal{V}]^{op}$.
- 2 A category \mathcal{A} **has \mathbb{L} -limits relative to (\mathbb{C}, γ)** , if $\{W, D\}_{(\mathbb{C}, \gamma)}$ exists for every \mathbb{L} -weight W and every diagram D .

^a (\mathbb{L}, λ) and (\mathbb{C}, γ) are **independent** of each other.

Main Theorem

For any \mathcal{A} , the following are equivalent:

- ① \mathcal{A} has \mathbb{L} -limits relative to (\mathbb{C}, γ) .
- ② $\mathbb{C}(\mathcal{A})$ has \mathbb{L} -limits of representables.
- ③ $\widetilde{\lambda}_{\mathcal{A}} : \mathbb{L}(\mathcal{A}) \rightarrow [\mathcal{A}^{op}, \mathcal{V}]$ is representable relative to (\mathbb{C}, γ) ,
i.e.,

$$\lambda_{\mathcal{A}} \dashv_{\gamma_{\mathcal{A}}} U : \mathbb{L}(\mathcal{A}) \rightarrow \mathbb{C}(\mathcal{A})$$

holds.

The Meaning of $U : \mathbb{L}(\mathcal{A}) \rightarrow \mathbb{C}(\mathcal{A})$

$$U : \text{Ran}_D W \mapsto \{W, \gamma_{\mathcal{A}} D\}$$

for every diagram $D : \mathcal{D} \rightarrow \mathcal{A}$ and every \mathbb{L} -weight $W : \mathcal{D} \rightarrow \mathcal{V}$.

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for every diagram $D : \mathcal{D} \rightarrow \mathcal{A}$ and every \mathbb{L} -weight $W : \mathcal{D} \rightarrow \mathcal{V}$.

Main Theorem when (\mathbb{C}, γ) is a Colimit Doctrine

For any \mathcal{A} , the following are equivalent:

- 1 \mathcal{A} has \mathbb{L} -limits relative to (\mathbb{C}, γ) .
- 2 For every X in $\mathbb{L}(\mathcal{A})$ there exists a pair $W_X : \mathcal{K}_X^{\text{op}} \rightarrow \mathcal{V}$, $J_X : \mathcal{K}_X \rightarrow \mathcal{A}$, with W_X a \mathbb{C} -weight such that there is an isomorphism

$$\mathbb{L}(\mathcal{A})(\lambda_{\mathcal{A}} A, X) \cong \int^{K \in \mathcal{K}_X^{\text{op}}} W_X K \otimes \mathcal{A}(A, J_X K)$$

natural in A .

- 3 $\mathbb{C}(\lambda_{\mathcal{A}}) : \mathbb{C}(\mathcal{A}) \rightarrow \mathbb{CL}(\mathcal{A})$ has a right adjoint that preserves \mathbb{C} -colimits.

Examples

- ① $\mathcal{V} = \text{Set}$, $\mathbb{L} = \text{finite limits}$, $\mathbb{C} = \text{finite colimits}$. \mathcal{A} has \mathbb{L} -limits relative to \mathbb{C} iff it has **finite plurilimits of finite diagrams** (P.Karazeris, J.Rosický, J.V., JPAA, 2005). Exploit the coend formula

$$\mathbb{L}(\mathcal{A})(\lambda_{\mathcal{A}} A, X) \cong \int^{K \in \mathcal{K}_X^{\text{op}}} W_X K \times \mathcal{A}(A, J_X K)$$

- $\mathcal{K}_X = \text{finite category}$
- $\mathbb{L}(\mathcal{A})(\lambda_{\mathcal{A}} A, X) = \text{set of cones for a finite diagram } X \text{ with vertex } A$
- $W_X K = \text{set of distinguished cones for a finite diagram } X \text{ with vertex } K$
- the coend provides the factorizations up to a zig-zag

Examples, cont.

- ② The same \mathbb{L} and \mathbb{C} as above (i.e., finite) but with \mathcal{V} =Abelian groups. **One-object** \mathcal{A} has \mathbb{L} -limits relative to \mathbb{C} iff it is a **left coherent ring**.

A ring A is left coherent iff the **dualization functor**

$$\mathrm{Hom}(-, A) : \mathrm{Mod}\text{-}A \rightarrow A\text{-}\mathrm{Mod}$$

restricts to categories of f.p. A -modules (R.R.Colby, J.Algebra, 1975).

- ③ Any \mathcal{V} , \mathbb{L} =small limits, \mathbb{C} =small colimits. \mathcal{A} has \mathbb{L} -limits relative to \mathbb{C} iff $\mathbb{C}(\mathcal{A})$ (the **category of small presheaves**) has **small limits of representables** (B.Day, S.Lack, JPAA, 2007).

④ ...

Corollary of the Main Theorem

The following are equivalent:

- 1 Every $\mathbb{C}(\mathcal{A})$ has \mathbb{L} -limits whenever \mathcal{A} has \mathbb{L} -limits.
- 2 Every $\mathbb{C}(\mathcal{A})$ has \mathbb{L} -limits whenever \mathcal{A} has \mathbb{L} -limits relative to \mathbb{C} .
- 3 Every $\mathbb{C}(\mathbb{L}(\mathcal{A}))$ has \mathbb{L} -limits.

These equivalent conditions are satisfied in the presence of a distributive law $\delta : \mathbb{L}\mathbb{C} \rightarrow \mathbb{C}\mathbb{L}$.

Promonoidal Structures

A **promonoidal structure** on \mathcal{A} is given by

$$J : \mathcal{I} \rightarrow [\mathcal{A}^{op}, \mathcal{V}] \quad P : \mathcal{A} \otimes \mathcal{A} \rightarrow [\mathcal{A}^{op}, \mathcal{V}]$$

such that P is “associative” and J is a “unit” (to within isomorphisms) — see B.Day, 1974.

Example

A promonoidal structure (\mathcal{A}, J, P) with J, P **representable** is precisely a **monoidal structure** on \mathcal{A} :

$$\begin{array}{ccc} \mathcal{I} & & \\ \downarrow & \searrow J & \\ \mathcal{A} & \xrightarrow{Y_{\mathcal{A}}} & [\mathcal{A}^{op}, \mathcal{V}] \end{array}$$

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Promonoidal Structures Representable Relative to (\mathbb{C}, γ)

$$\begin{array}{ccc}
 \mathcal{I} & & \mathcal{A} \otimes \mathcal{A} \\
 \downarrow \text{dotted} & \searrow J & \downarrow \text{dotted} \\
 \mathbb{C}(\mathcal{A}) & \xrightarrow[\gamma_{\mathcal{A}}]{} & [\mathcal{A}^{op}, \mathcal{V}] \\
 & & \downarrow P \\
 & & \mathbb{C}(\mathcal{A}) \xrightarrow[\gamma_{\mathcal{A}}]{} [\mathcal{A}^{op}, \mathcal{V}]
 \end{array}$$

are precisely monoidal structures on $\mathbb{C}(\mathcal{A})$, if \mathbb{C} is a doctrine of free cocompletions.

(B.Day, S.Lack, JPAA, 2007 for \mathbb{C} =small colimits).

Example: Flatness and Merging

Suppose \mathcal{A} has \mathbb{L} -limits relative to (\mathbb{C}, γ) and let \mathcal{B} have \mathbb{L} -limits and $\widetilde{\gamma}_{\mathcal{A}}$ -colimits. A functor $H : \mathcal{A} \rightarrow \mathcal{B}$ is called

- ① **\mathbb{L} -flat relative to (\mathbb{C}, γ)** if $\text{Lan}_{\gamma_{\mathcal{A}}} H : \mathbb{C}(\mathcal{A}) \rightarrow \mathcal{B}$ preserves \mathbb{L} -limits of representables.
- ② **merging \mathbb{L} -limits relative to (\mathbb{C}, γ)** if the canonical comparison

$$\widetilde{\gamma}_{\mathcal{A}}(\{W, \gamma_{\mathcal{A}} D\}) * H \rightarrow \{W, HD\}$$

is an isomorphism for every $D : \mathcal{D} \rightarrow \mathcal{A}$ and every \mathbb{L} -weight $W : \mathcal{M} \rightarrow [\mathcal{D}, \mathcal{V}]^{\text{op}}$ (introduced by H.Hu, W.Tholen).

Result: these concepts are **equivalent**.

Hear Panagis' talk for further applications.