

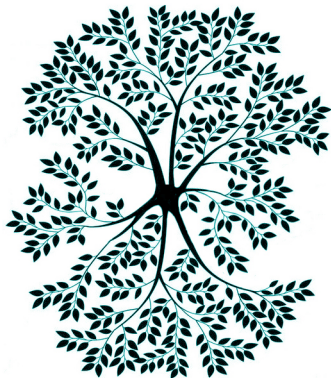


Categorified cyclic operads (in nature)

Category Theory 2018
University of Azores

Jovana Obradović

& Pierre-Louis Curien



1. Categorification
2. Cyclic operads
3. Categorified cyclic operads
 - 3.1 The coherence theorem
 - 3.2 Categorified cyclic operads “in nature”
 - 3.3 Polytopes of categorified cyclic operads

CATEGORIFICATION

Relevant examples of categorified structures

sets

functions

equalities

categories

functors

coherent isomorphisms

Relevant examples of categorified structures

sets

functions

equalities

categories

functors

coherent isomorphisms

Coherence of symmetric monoidal categories



S. Mac Lane

Categories for the Working Mathematician

Springer, 1997

$$\beta_{f,g,h} : (fg)h \rightarrow f(gh) \quad \gamma_{f,g} : fg \rightarrow gf$$

Relevant examples of categorified structures

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Coherence of categorified non-symmetric skeletal operads



K. Došen, Z. Petrić

Weak Cat-operads

Logical Methods in Computer Science, 2009

$$\beta_{f,g,h}^{i,j} : (f \circ_i g) \circ_j h \rightarrow f \circ_i (g \circ_{j-i+1} h) \quad \theta_{f,g,h}^{i,j} : (f \circ_i g) \circ_j h \rightarrow (f \circ_j h) \circ_{i+n-1} g$$

CYCLIC OPERADS

Cyclic operads: definition

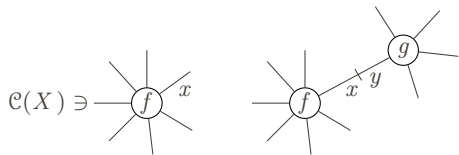
$$\mathcal{C} : \mathbf{Bij}^{op} \rightarrow \mathbf{Set}$$

$$\mathcal{C}(X) \ni \begin{array}{c} \diagup \\ | \\ \textcircled{f} \\ | \\ \diagdown \end{array} x$$

Cyclic operads: definition

$$\mathcal{C} : \mathbf{Bij}^{op} \rightarrow \mathbf{Set}$$

$$x \circ_y : \mathcal{C}(X) \times \mathcal{C}(Y) \rightarrow \mathcal{C}(X \setminus \{x\} \cup Y \setminus \{y\})$$

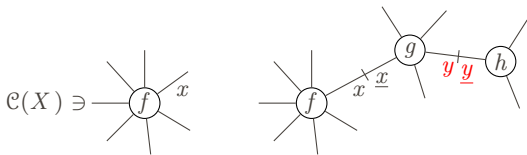


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$$(f \circ_{x \underline{x}} g) \circ_{y \underline{y}} h = f \circ_{x \underline{x}} (g \circ_{y \underline{y}} h) \quad f \circ_{x \circ_y} g = g \circ_{y \circ_x} f$$

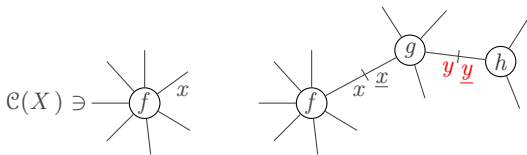


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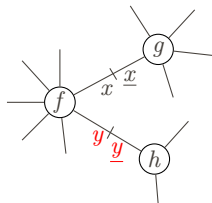
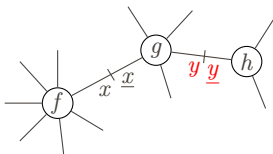
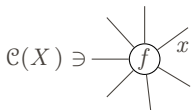
$$f^{\sigma_1} \circ_{\sigma_1^{-1}(x)} \circ_{\sigma_2^{-1}(y)} g^{\sigma_2} = (f \circ_{x \circ_y} g)^{\sigma}$$

Cyclic operads: definition

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Parallel associativity

$$(f \circ_{x \circ \underline{x}} g) \circ_{y \circ \underline{y}} h = (g \circ_{\underline{x} \circ x} f) \circ_{y \circ \underline{y}} h = g \circ_{\underline{x} \circ x} (f \circ_{y \circ \underline{y}} h) = (f \circ_{y \circ \underline{y}} h) \circ_{x \circ \underline{x}} g$$

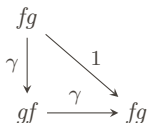
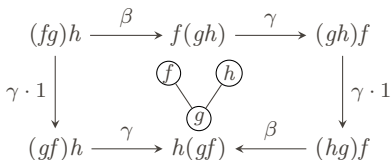
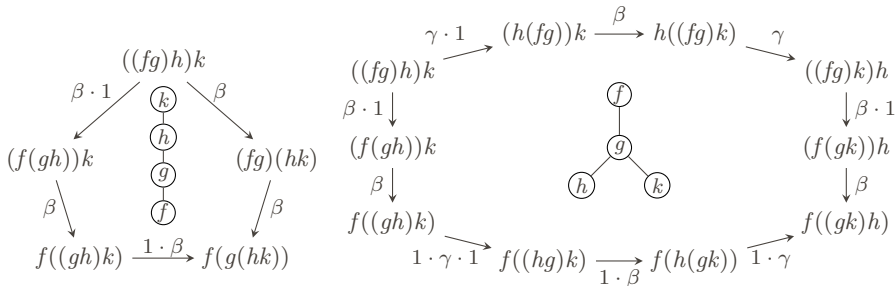
CATEGORIFIED CYCLIC OPERADS

Categorified cyclic operads: the definition

$$\mathcal{C} : \mathbf{Bij}^{op} \rightarrow \mathbf{Cat}$$

$$x \circ_y : \mathcal{C}(X) \times \mathcal{C}(Y) \rightarrow \mathcal{C}(X \setminus \{x\} \cup Y \setminus \{y\})$$

$$\beta_{f,g,h}^{x,\underline{x};y,\underline{y}} : (f \circ_{x\underline{x}} g) \circ_{y\underline{y}} h \rightarrow f \circ_{x\underline{x}} (g \circ_{y\underline{y}} h) \quad \gamma_{f,g}^{x,y} : f \circ_{x\underline{x}} g \rightarrow g \circ_{y\underline{y}} f$$



$$\beta_{f,g,h}^\sigma = \beta_{f^{\sigma 1}, g^{\sigma 2}, h^{\sigma 3}}$$

$$\gamma_{f,g}^\sigma = \gamma_{f^{\sigma 1}, g^{\sigma 2}}$$

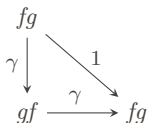
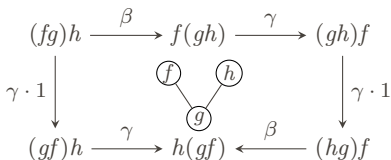
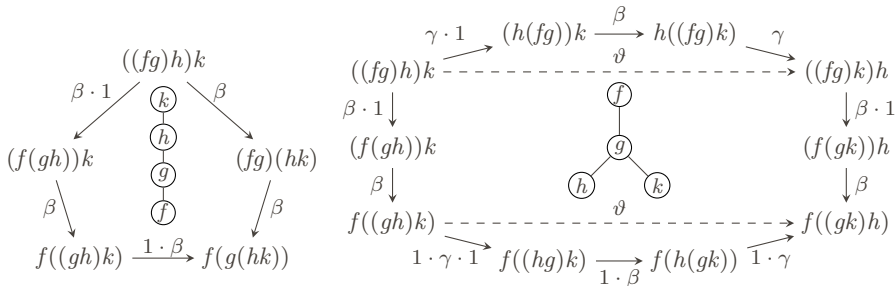
$$(\varphi \cdot \psi)^\sigma = \varphi^{\sigma 1} \cdot \psi^{\sigma 2}$$

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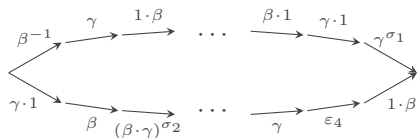
$$\gamma_{f,g}^\sigma = \gamma_{f^{\sigma 1}, g^{\sigma 2}}$$

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The coherence theorem

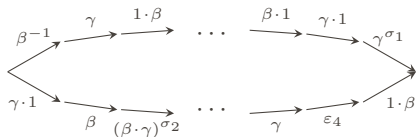
The coherence theorem: a formal language

Every diagram of canonical arrows in $\mathcal{C}(X)$ commutes.



The coherence theorem: a formal language

Every diagram of canonical arrows in $\mathcal{C}(X)$ commutes.



For the syntax $\text{Free}_{\mathcal{C}}$ of $\beta\gamma\sigma$ -diagrams & the coherence theorem:



P.-L. Curien, J. Obradović
Categorified cyclic operads
arXiv:1706.06788v2, 2018

The proof scheme

Coherence of categorified non-skeletal cyclic operads with symmetries

All $\beta\gamma\sigma$ -diagrams in $\mathcal{C}(X)$ commute.

Došen & Petrić: Coherence of categorified skeletal operads without symmetries

For an arbitrary categorified operad \mathcal{O} ,
all $\beta\theta$ -diagrams in $\mathcal{O}(n)$ commute.

The proof scheme

Coherence of categorified non-skeletal cyclic operads with symmetries

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1st reduction:
removing symmetries

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All $\beta\gamma$ -diagrams in $\mathcal{C}(X)$ commute.

2nd reduction:
removing cyclicity

All $\beta\vartheta$ -diagrams in $\mathcal{C}(X)$ commute.

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$$(f \circ_{x \circ y} g)^\sigma \rightsquigarrow f^{\sigma_1} \circ_{\sigma_1^{-1}(x)} \circ_{\sigma_2^{-1}(y)} g^{\sigma_2}$$

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2nd reduction:
removing cyclicity

All $\beta\vartheta$ -diagrams in $\mathcal{C}(X)$ commute.

3rd reduction:
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$$\mathbf{Bij} \simeq \Sigma$$

Došen & Petrić: Coherence of categorified skeletal operads without symmetries

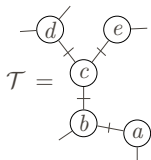
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2nd reduction: removing cyclicity

$$\underline{\text{Free}}_{\underline{c}}(X) \ni (\underline{a} \square_{\underline{x}} \underline{b}) \square_{\underline{y}} \underline{c} \iff \left(\begin{array}{c} \textcircled{c} \\ | \\ \underline{y} \\ | \\ \textcircled{b} \\ | \\ \underline{x} \\ | \\ \textcircled{a} \end{array} \right), (\underline{a} \underline{b}) \underline{c} \in \underline{\mathbb{T}}_{\underline{c}}(X)$$

2nd reduction: removing cyclicity

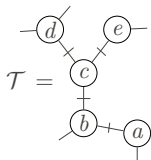
$$\underline{\text{Free}}_{\underline{c}}(X) \ni (\underline{a} \ \underline{x} \square \underline{x} \ \underline{b}) \ \underline{y} \square \underline{y} \ \underline{c} \iff \left(\begin{array}{c} \textcircled{c} \\ | \\ \underline{y} \\ | \\ \textcircled{b} \\ | \\ \underline{x} \\ | \\ \textcircled{a} \end{array} , (\underline{a} \ \underline{b}) \underline{c} \right) \in \underline{\mathbb{T}}_{\underline{c}}(X)$$



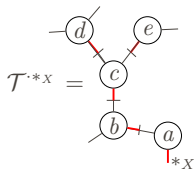
$$a(b(d(ec)))$$

2nd reduction: removing cyclicity

$$\underline{\text{Free}}_{\underline{c}}(X) \ni (\underline{a} \ x \square_x \ \underline{b}) \ y \square_y \ \underline{c} \iff \left(\begin{array}{c} \textcircled{c} \\ | \\ \underline{y} \\ | \\ \textcircled{b} \\ | \\ \underline{x} \\ | \\ \textcircled{a} \end{array} , (\underline{a} \ \underline{b}) \underline{c} \right) \in \underline{\mathbb{T}}_{\underline{c}}(X)$$

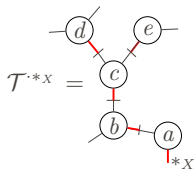
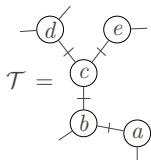


$$a(b(d(ec)))$$



2nd reduction: removing cyclicity

$$\underline{\text{Free}}_{\underline{c}}(X) \ni (\underline{a} \ \underline{x} \square \underline{x} \ \underline{b}) \ \underline{y} \square \underline{y} \ \underline{c} \iff \left(\begin{array}{c} \textcircled{c} \\ \downarrow \underline{y} \\ \textcircled{b} \\ \downarrow \underline{x} \\ \textcircled{a} \end{array} , (\underline{a} \ \underline{b}) \underline{c} \right) \in \underline{\mathbb{T}}_{\underline{c}}(X)$$



$$a(b(d(ec)))$$

$$\downarrow 1_a \cdot 1_b \cdot \gamma_{d,ec}$$

$$a(b((ec)d))$$

$$\downarrow 1_a \cdot 1_b \cdot \gamma_{e,c} \cdot 1_d$$

$$a(b((ce)d))$$

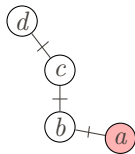
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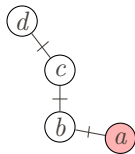
$$((ba)c)d \xrightarrow{\beta_{ba,c,d}} (ba)(cd)$$



2nd reduction: removing cyclicity

$$\underline{\text{Free}}_{\underline{c}}(X) \ni (\underline{a} \ \underline{x} \ \underline{x} \ \underline{b}) \ \underline{y} \ \underline{y} \ \underline{c} \Leftrightarrow \left(\begin{array}{c} \textcircled{c} \\ | \\ \textcircled{b} \\ | \\ \textcircled{a} \end{array} \right), (\underline{a} \ \underline{b}) \ \underline{c} \in \underline{\mathbf{T}}_{\underline{c}}(X)$$

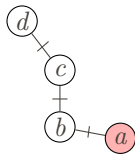
$$\begin{array}{ccc} ((ba)c)d & \xrightarrow{\beta_{ba,c,d}} & (ba)(cd) \\ \gamma_{b,a} \cdot 1_c \cdot 1_d \downarrow & & \downarrow \gamma_{b,a} \cdot 1_{cd} \\ ((ab)c)d & & (ab)(cd) \end{array}$$



2nd reduction: removing cyclicity

$$\underline{\text{Free}}_{\underline{c}}(X) \ni (\underline{a} \ \underline{x} \square \underline{x} \ \underline{b}) \ \underline{y} \square \underline{y} \ \underline{c} \iff \left(\begin{array}{c} \textcircled{c} \\ | \\ \textcircled{b} \\ | \\ \textcircled{a} \end{array} \right), (\underline{a} \ \underline{b}) \underline{c} \in \underline{\mathbf{T}}_{\underline{c}}(X)$$

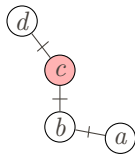
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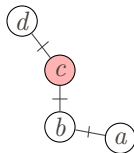
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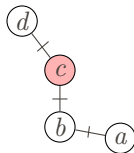
$$\begin{array}{ccc} ((ba)c)d & \xrightarrow{\beta_{ba,c,d}} & (ba)(cd) \\ \downarrow \gamma_{ba,c} \cdot 1_d & & \downarrow \gamma_{ba,cd} \\ (c(ba))d & & (cd)(ba) \end{array}$$



2nd reduction: removing cyclicity

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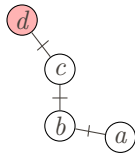
$$\begin{array}{ccc} ((ba)c)d & \xrightarrow{\beta_{ba,c,d}} & (ba)(cd) \\ \gamma_{ba,c} \cdot 1_d \downarrow & & \downarrow \gamma_{ba,cd} \\ (c(ba))d & \xrightarrow{\vartheta_{c,ba,d}} & (cd)(ba) \end{array}$$



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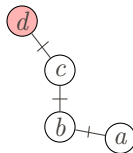
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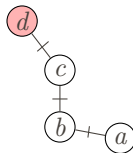
$$\begin{array}{ccc} ((ba)c)d & \xrightarrow{\beta_{ba,c,d}} & (ba)(cd) \\ \gamma_{(ba)c,d} \downarrow & & \downarrow \gamma_{ba,cd} \\ d((ba)c) & & (cd)(ba) \end{array}$$



2nd reduction: removing cyclicity

$$\underline{\text{Free}}_{\underline{c}}(X) \ni (\underline{a} \ \underline{x} \ \underline{x} \ \underline{b}) \ \underline{y} \ \underline{y} \ \underline{c} \iff \left(\begin{array}{c} \textcircled{c} \\ | \\ \textcircled{b} \\ | \\ \textcircled{a} \end{array} \right), (\underline{a} \ \underline{b}) \ \underline{c} \in \underline{\mathbb{T}}_{\underline{c}}(X)$$

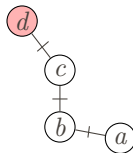
$$\begin{array}{ccc} ((ba)c)d & \xrightarrow{\beta_{ba,c,d}} & (ba)(cd) \\ \gamma_{(ba)c,d} \downarrow & & \downarrow \gamma_{ba,cd} \\ d((ba)c) & & (cd)(ba) \\ 1_d \cdot \gamma_{ba,c} \downarrow & & \downarrow \gamma_{c,d} \cdot 1_{ba} \\ d(c(ba)) & & (dc)(ba) \end{array}$$



2nd reduction: removing cyclicity

$$\underline{\text{Free}}_{\underline{c}}(X) \ni (\underline{a} \ \underline{x} \ \underline{x} \ \underline{b}) \ \underline{y} \ \underline{y} \ \underline{c} \Leftrightarrow \left(\begin{array}{c} \textcircled{c} \\ | \\ \textcircled{b} \\ | \\ \textcircled{a} \end{array} \right), (\underline{a} \ \underline{b}) \ \underline{c} \in \underline{\mathbf{T}}_{\underline{c}}(X)$$

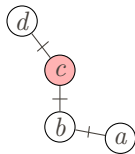
$$\begin{array}{ccc} ((ba)c)d & \xrightarrow{\beta_{ba,c,d}} & (ba)(cd) \\ \gamma_{(ba)c,d} \downarrow & & \downarrow \gamma_{ba,cd} \\ d((ba)c) & & (cd)(ba) \\ 1_d \cdot \gamma_{ba,c} \downarrow & & \downarrow \gamma_{c,d} \cdot 1_{ba} \\ d(c(ba)) & \xrightarrow{\beta_{d,g,ca}^{-1}} & (dc)(ba) \end{array}$$



2nd reduction: removing cyclicity

$$\underline{\text{Free}}_{\underline{c}}(X) \ni (\underline{a} \ \underline{x} \square \underline{x} \ \underline{b}) \ \underline{y} \square \underline{y} \ \underline{c} \iff \left(\begin{array}{c} \textcircled{c} \\ | \\ \textcircled{b} \\ | \\ \textcircled{a} \end{array} \right), (\underline{a} \ \underline{b}) \underline{c} \in \underline{\mathbf{T}}_{\underline{c}}(X)$$

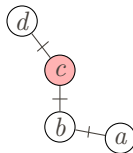
$$(ba)(cd) \xrightarrow{\gamma^{ba,cd}} (cd)(ba)$$



2nd reduction: removing cyclicity

$$\underline{\text{Free}}_{\underline{c}}(X) \ni (\underline{a} \ \underline{x} \ \underline{x} \ \underline{b}) \ \underline{y} \ \underline{y} \ \underline{c} \iff \left(\begin{array}{c} \textcircled{c} \\ | \\ \textcircled{b} \\ | \\ \textcircled{a} \end{array} \right), (\underline{a} \ \underline{b}) \ \underline{c} \in \underline{\mathbf{T}}_{\underline{c}}(X)$$

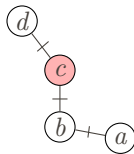
$$\begin{array}{ccc} (ba)(cd) & \xrightarrow{\gamma_{ba,cd}} & (cd)(ba) \\ \gamma_{ba,cd} \downarrow & & \downarrow 1_{ba,cd} \\ (cd)(ba) & & (cd)(ba) \end{array}$$



2nd reduction: removing cyclicity

$$\underline{\text{Free}}_{\underline{c}}(X) \ni (\underline{a} \ \underline{x} \ \underline{x} \ \underline{b}) \ \underline{y} \ \underline{y} \ \underline{c} \Leftrightarrow \left(\begin{array}{c} \textcircled{c} \\ | \\ \textcircled{b} \\ | \\ \textcircled{a} \end{array} \right), (\underline{a} \ \underline{b}) \ \underline{c} \in \underline{\mathbf{T}}_{\underline{c}}(X)$$

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Categorified cyclic operads “in nature”



J. Bénabou

Les distributeurs

Université Catholique de Louvain

Institut de Mathématique Pure et Appliquée, rapport 33, 1973

Bicategory **Prof**

- objects: small categories
- 1-morphisms: profunctors $F : \mathbf{C} \nrightarrow \mathbf{D}$, i.e. functors $F : \mathbf{D}^{op} \times \mathbf{C} \rightarrow \mathbf{Set}$

$$G \circ F = \int^d F(d, -) \times G(-, d)$$

- 2-morphisms: natural transformations

Generalised profunctors as categorified cyclic operads

- \mathbf{D} small category, $(-)^* : \mathbf{D}^{op} \rightarrow \mathbf{D}$ duality

$F : \mathbf{D}^{op} \times \mathbf{D} \rightarrow \mathbf{Set}$ can be considered as $F : \mathbf{D} \times \mathbf{D} \rightarrow \mathbf{Set}$

Generalised profunctors as categorified cyclic operads

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- \mathbf{D}^n -profunctors: $\mathcal{C}(n) = [\mathbf{D}^n, \mathbf{Set}]$

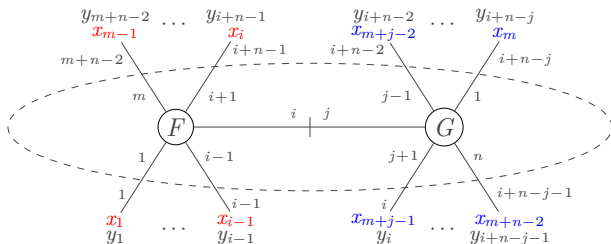
Generalised profunctors as categorified cyclic operads

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$F : \mathbf{D}^{op} \times \mathbf{D} \rightarrow \mathbf{Set}$ can be considered as $F : \mathbf{D} \times \mathbf{D} \rightarrow \mathbf{Set}$

- \mathbf{D}^n -profunctors: $\mathcal{C}(n) = [\mathbf{D}^n, \mathbf{Set}]$
- $i \circ_j : \mathcal{C}(m) \times \mathcal{C}(n) \rightarrow \mathcal{C}(m+n-2)$

$$(F \circ_j G)(y_1, \dots, y_{m+n-2}) = \int^{u,v} F(x_1, \dots, x_{i-1}, u, x_i, \dots, x_{m-1}) \times G(x_m, \dots, x_{m+j-2}, v, x_{m+j-1}, \dots, x_{m+n-2}) \times \mathbf{D}[u, v^*]$$



Feynman category for cyclic operads admits an odd version



R. Kaufmann, B. Ward

Feynman Categories

Astérisque (Société Mathématique de France), Numéro 387, 2017

Cyclic operads are representations of the Feynman category Cyc .

Feynman category for cyclic operads admits an odd version



R. Kaufmann, B. Ward

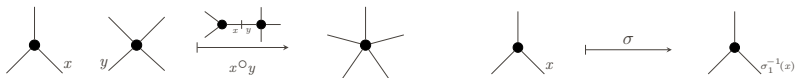
Feynman Categories

Astérisque (Société Mathématique de France), Numéro 387, 2017

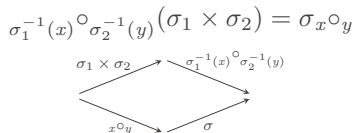
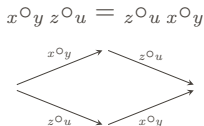
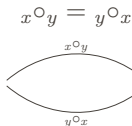
Cyclic operads are representations of the Feynman category Cyc .

Generators-and-relations representation of Cyc :

- objects: sets of corollas
- morphisms:



- relations:



Feynman category for cyclic operads admits an odd version

Cyc admits an **ordered presentation** if there exists

$$sgn : \left\{ \begin{array}{c} \begin{array}{c} \xrightarrow{x \circ y} \\ \xleftarrow{y \circ x} \end{array} \quad , \quad \begin{array}{c} \xrightarrow{x \circ y} \quad \xrightarrow{z \circ u} \\ \xleftarrow{z \circ u} \quad \xleftarrow{x \circ y} \end{array} \quad , \quad \begin{array}{c} \xrightarrow{\sigma_1 \times \sigma_2} \quad \xrightarrow{\sigma_1^{-1}(x) \circ \sigma_2^{-1}(y)} \\ \xleftarrow{x \circ y} \quad \xleftarrow{\sigma} \end{array} \end{array} \right\} \rightarrow \{+, -\}$$

such that the **sign coherence** holds:

*sgn extends multiplicatively to every “polygon”;
every two parallel polygons receive the same sign*

Feynman category for cyclic operads admits an odd version

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Theorem

Cyc can be ordered by assigning $-$ to $x \circ y = y \circ x$ and $+$ to other relations.

Feynman category for cyclic operads admits an odd version

Cyc admits an **ordered presentation** if there exists

$$sgn : \left\{ \begin{array}{c} \text{Diagram 1: } \begin{array}{c} \xrightarrow{x \circ y} \\ \xleftarrow{y \circ x} \end{array} \quad , \quad \text{Diagram 2: } \begin{array}{c} \xrightarrow{x \circ y} \quad \xrightarrow{z \circ u} \\ \xleftarrow{z \circ u} \quad \xleftarrow{x \circ y} \end{array} \quad , \quad \text{Diagram 3: } \begin{array}{c} \xrightarrow{\sigma_1 \times \sigma_2} \quad \xrightarrow{\sigma_1^{-1}(x) \circ \sigma_2^{-1}(y)} \\ \xleftarrow{x \circ y} \quad \xleftarrow{\sigma} \end{array} \end{array} \right\} \rightarrow \{+, -\}$$

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$(Cyc, sgn) \Rightarrow ||(Cyc, sgn)||^{odd} \mathbf{Ab}$ -enriched Feynman category
for **anti-cyclic operads**

Polytopes of categorified cyclic operads

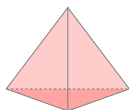
Categorified operads: Hypergraph polytopes



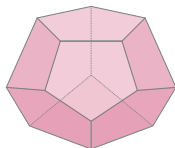
K. Došen, Z. Petrić

Hypergraph polytopes

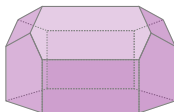
Topology and its Applications 158, pp. 1405–1444, 2011



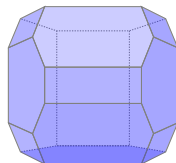
simplex



associahedron



hemiassociahedron



permutohedron



“the more hyperedges, the more truncations”

Hypergraph arrangements of hypercubes

Categorified operads

associahedron K_n

hemiassoiahedron H_n

permutohedron P_n

\cap

hypergraph polytopes

Categorified cyclic operads

K_n -arrangement of hypercubes

H_n -arrangement of hypercubes

P_n -arrangement of hypercubes

\cap

hypergraph arrangements of hypercubes

Hypergraph arrangements of hypercubes

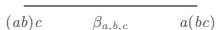
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Categorified cyclic operads

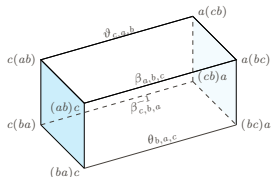
K_n -arrangement of hypercubes

H_n -arrangement of hypercubes

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\cap

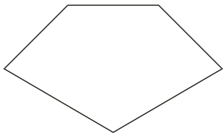
hypergraph arrangements of hypercubes



Hypergraph arrangements of hypercubes

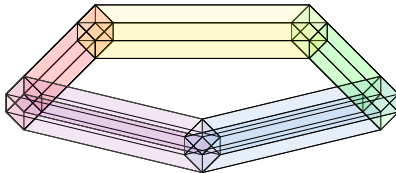
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 \cap
hypergraph polytopes



Categorified cyclic operads

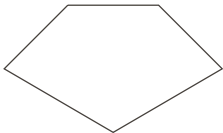
K_n -arrangement of hypercubes
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hypergraph arrangements of hypercubes



Hypergraph arrangements of hypercubes

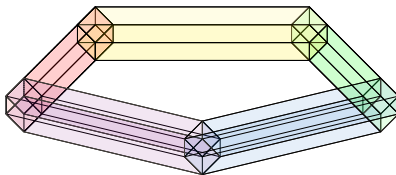
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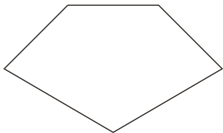


$$\mathcal{A}_{ha}(\mathbf{H}) = \mathcal{A}(\mathbf{H}) \times \mathbf{Q}_{|H|}$$

Hypergraph arrangements of hypercubes

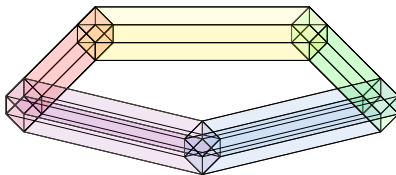
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Thank you!