Forking in accessible categories

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joint work with M.Lieberman and S. Vasey

Ponta Delgada 2018

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They are the same as μ -abstract elementary classes.

Any μ -accessible category whose morphisms are monomorphisms is a μ -AEC and any μ -AEC is λ^+ -accessible where λ is its LST number (LRV + R. Grossberg and W. Boney 2016).

Accessible categories whose morphisms are monomorphisms cannot be locally presentable.

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Theorem 1. Locally multipresentable categories whose morphisms are monomorphisms coincide with universal μ -AECs.

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A *polyinitial* object is a set \mathcal{I} of objects of a category \mathcal{K} such that for every object M in \mathcal{K} :

- 1. There is a unique $i \in \mathcal{I}$ having a morphism $i \to M$.
- 2. For each $i \in \mathcal{I}$, given $f, g : i \to M$, there is a unique (isomorphism) $h : i \to i$ with fh = g.

We get groups of automorphisms of members of a polyinitial object. In the case of a multinitial objects, they are singletons. An example of a locally polypresentable category whose morphisms are monomorphisms are algebraically closed fields. The polyinitial object is formed by algebraic closures of the multiinitial object in fields.

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Lemma 1. Let \mathcal{K} be a coregular locally μ -presentable category and \mathcal{K}_{reg} be the category having the same objects as \mathcal{K} and regular monomorphisms as morphisms. Then \mathcal{K}_{reg} is locally μ -multipresentable.

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Examples of coregular locally presentable categories: Grothendieck toposes, Grothendieck abelian categories, **Gra** graphs, **Gr** groups, **Bool** Boolean algebras, **Ban** Banach spaces with linear contractions, **Hilb** Hilbert spaces with linear isometries, **CAlg** commutative unital C^* -algebras, etc.

Let \mathcal{K} be a coregular locally presentable category and

$$M_1$$
 \uparrow
 $M_0 \rightarrow M_2$

a span in \mathcal{K}_{reg} . Let

$$\begin{array}{c} M_1 \longrightarrow I \\ \uparrow & \uparrow \\ M_0 \longrightarrow M_2 \end{array}$$

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a pushout in $\mathcal{K}.$ Then a multipushout in $\mathcal{K}_{\textit{reg}}$ is formed by all squares

$$\begin{array}{c} M_1 \longrightarrow Q \\ \uparrow & \uparrow \\ M_0 \longrightarrow M_2 \end{array}$$

where the induced morphism $P \rightarrow Q$ is an epimorphism.

The notion of independence ${\buildrel } {\buildrel }$ in ${\buildrel } {\buildrel } {\buildrel } {\buildrel } {\buildrel }$ in ${\buildrel } {\buildrel } {\buildrel }$ in the choice of squares

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which are declared to be independent. We say that M_1 and M_2 are independent over M_0 in M_3 .

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The notion of independence ${\bf \perp}$ in ${\cal K}$ consists in the choice of squares

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which are declared to be independent. We say that M_1 and M_2 are independent over M_0 in M_3 .

The following properties should be satisfied

- (i) invariance under isomorphisms of squares
- (ii) independence on M_3 ,
- (iii) existence,
- (iv) uniqueness,
- (v) symmetry,
- (vi) closedness under compositions of squares, and
- (vii) accessibility.



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(iii) any span can be completed to an independent square,



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(vii) the category whose objects are morphisms in ${\cal K}$ and whose morphisms are independent squares is accessible.

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More generally, if \mathcal{K} has the notion of independence \downarrow then any $\stackrel{1}{\downarrow}$ with (i-vi) equals to \downarrow .

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More generally, if \mathcal{K} has the notion of independence \downarrow then any $\stackrel{1}{\downarrow}$ with (i-vi) equals to \downarrow .

Theorem 3. Let \mathcal{K} be an accessible category whose morphisms are monomorphism having a notion of independence. Then \mathcal{K} is tame, stable and does not have the order property.

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More generally, if \mathcal{K} has the notion of independence \downarrow then any $\stackrel{1}{\downarrow}$ with (i-vi) equals to \downarrow .

Theorem 3. Let \mathcal{K} be an accessible category whose morphisms are monomorphism having a notion of independence. Then \mathcal{K} is tame, stable and does not have the order property.

Theorem 4. Let \mathcal{K} be a coregular locally presentable category with effective unions. Then \mathcal{K}_{reg} has an independence notion (consisting of pullback squares).

 ${\cal K}$ has effective unions if whenever we have a pullback

$$M_1 \rightarrow M_3$$
 $\uparrow \qquad \uparrow$
 $M_0 \rightarrow M_2$

and a pushout

$$\begin{array}{c} M_1 \longrightarrow P \\ \uparrow & \uparrow \\ M_0 \longrightarrow M_2 \end{array}$$

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Gra, **Gr**, **Ban**, **Bool** or **CAlg** do not have effective unions. They do not have a notion of independence because they have the order property.

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Corollary 1. Let \mathcal{K} be a coregular locally presentable category with an independence notion \bot in \mathcal{K}_{reg} . Then \bot is given by effective pullback squares.

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Corollary 1. Let \mathcal{K} be a coregular locally presentable category with an independence notion \bot in \mathcal{K}_{reg} . Then \bot is given by effective pullback squares.

In Gra, in an effective pullback square

$$\begin{array}{c} M_1 \rightarrow M_3 \\ \uparrow & \uparrow \\ M_0 \rightarrow M_2 \end{array}$$

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There is another attempt of independence where we include all cross-edges between M_1 and M_2 . This yields \downarrow satisfying (i-vi). By Theorem 2, **Gra** does not have an independence notion.

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 \mathcal{K} does not have chain bounds but the proof of Theorem 2 still goes through – thus this independence is unique.

Based on Malliaris and Shelah 2011, we say that a graph is stable if it does not contain a copy of the half graph (the bipartite graph on $\mathbb{N} \times \mathbb{N}$ such that E(i, j) iff i < j). The category \mathcal{K} of stable graphs is locally \aleph_1 -multipresentable and effective pullback squares do not form an independence notion in \mathcal{K}_{reg} . But we do not know whether \mathcal{K}_{reg} has an independence notion. \mathcal{K} does not have chain bounds and we could not adapt the proof of Theorem 2 to this case.

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Let \mathcal{K} be a locally polypresentable category whose morphisms are monomorphisms having a stable independence notion. Then, for each span, exactly one instance of a polypushout is independent. Moreover, a morphism of spans

$$(\mathrm{id}_{M_0}, h_1, h_2) : (M_0, M_1, M_2) \to (M_0, M_1', M_2')$$

induces a morphism of independent instances of polypushouts. Thus the independence yields a coherent choice of polypushouts. Based on Malliaris and Shelah 2011, we say that a graph is stable if it does not contain a copy of the half graph (the bipartite graph on $\mathbb{N} \times \mathbb{N}$ such that E(i,j) iff i < j). The category \mathcal{K} of stable graphs is locally \aleph_1 -multipresentable and effective pullback squares do not form an independence notion in \mathcal{K}_{reg} . But we do not know whether \mathcal{K}_{reg} has an independence notion. \mathcal{K} does not have chain bounds and we could not adapt the proof of Theorem 2 to this case.

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Theorem 5. Let \mathcal{K} be a coregular locally presentable category where regular monomorphisms are closed under directed colimits. Then \mathcal{K}_{reg} has a stable independence notion iff regular monomorphisms are cofibrantly generated.

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Consequently, **Gr**, **Ban**, **Bool** and **CAIg** do not have a notion of independence. **Gr** do not have enough regular injectives and thus regular monomorphisms cannot be cofibrantly generated. In all other cases, regular injectives do not form an accessible category and thus regular monomorphisms cannot be cofibrantly generated again.

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Consequently, **Gr**, **Ban**, **Bool** and **CAIg** do not have a notion of independence. **Gr** do not have enough regular injectives and thus regular monomorphisms cannot be cofibrantly generated. In all other cases, regular injectives do not form an accessible category and thus regular monomorphisms cannot be cofibrantly generated again.

Another consequence is that regular monomorphisms in **Gra** are not cofibrantly generated. Equivalently, **Gra** does not have enough regular injectives or those are not accessible. **Theorem 6.** Let \mathcal{K} be an accessible category whose morphisms are monomorphisms having the amalgamation property and chain bounds. Let κ be a strongly compact cardinal. If \mathcal{K} does not have the order property then the full subcategory of \mathcal{K} consisting of κ -saturated objects has an independence notion.

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Theorem 6. Let \mathcal{K} be an accessible category whose morphisms are monomorphisms having the amalgamation property and chain bounds. Let κ be a strongly compact cardinal. If \mathcal{K} does not have the order property then the full subcategory of \mathcal{K} consisting of κ -saturated objects has an independence notion.

Theorem 7. Let \mathcal{K} be an accessible category whose morphisms are monomorphisms having an independence notion. Then there exists a regular cardinal κ such that any independent square with M_0 κ -saturated is a pullback square.

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