

Positive Semidefinite Rank

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Section 1

What is it

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$$\begin{bmatrix} 2 & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 2 & 1 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 1 & 1 & 1 \end{bmatrix}$$

Additional matrices shown above the main equation:

$$\begin{bmatrix} 1/2 & -1/2 \\ -1/2 & 1 \end{bmatrix} \begin{bmatrix} 1/2 & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix}$$

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The smallest size of a semidefinite factorization is defined to be the positive semidefinite rank of M , $\text{rank}_{\text{psd}}(M)$

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Section 2

Why I do I care

Semidefinite Representations

A semidefinite representation of size k of a polytope P is a description

$$P = \left\{ x \in \mathbb{R}^n \mid \exists y \text{ s.t. } A_0 + \sum A_i x_i + \sum B_i y_i \succeq 0 \right\}$$

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This tells us how hard it is to optimize over P using semidefinite programming.

The Square

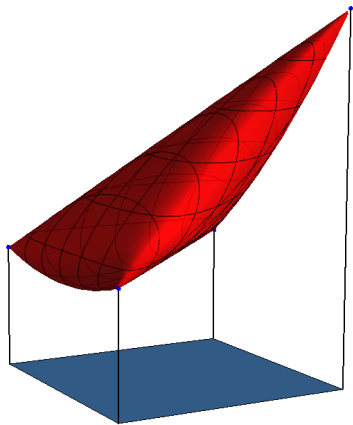
The 0/1 square is the projection onto x_1 and x_2 of

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$$\begin{array}{c|c|c|c|c|c|c|c} 0 & 1 & 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \end{array}$$

$$\begin{array}{l} x \geq 0 \\ y \geq 0 \\ z \geq 0 \\ 1 - x \geq 0 \\ 1 - y \geq 0 \\ 1 - z \geq 0 \end{array} \left[\begin{array}{c} \\ \\ \\ \\ \\ \\ \end{array} \right]$$

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	0	0	1	0	1	1	0	1	
	0	0	0	1	0	1	1	1	
$x \geq 0$	[0	1	0	0	1	0	1	1
$y \geq 0$		0	0	1	0	1	1	0	1
$z \geq 0$		0	0	0	1	0	1	1	1
$1 - x \geq 0$		0	1	0	0	1	0	1	1
$1 - y \geq 0$		0	0	1	0	1	1	0	1
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Theorem (G.-Parrilo-Thomas 2011)

A polytope P has a semidefinite representation of size k if and only if $\text{rank}_{\text{psd}}(S_P) \leq k$.

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- ▶ What is the psd rank of the travelling salesman polytope?

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Many interesting open questions are raised. The two most important:

- ▶ What is the psd rank of the travelling salesman polytope?
- ▶ Can psd lifts do much better than linear lifts?

The end



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Positive semidefinite rank.

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Thank you