



Laboratoire
Méthodes
Formelles

Jean Goubault-Larrecq

Noetherian spaces, wqos, and their statures

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PARIS-SACLAY



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Outline

- ❖ Noetherian spaces and wqos
- ❖ A computer scientist's view
- ❖ Sobrifications of Noetherian spaces, and their representations
- ❖ Statures of Noetherian spaces and maximal order types of wqos

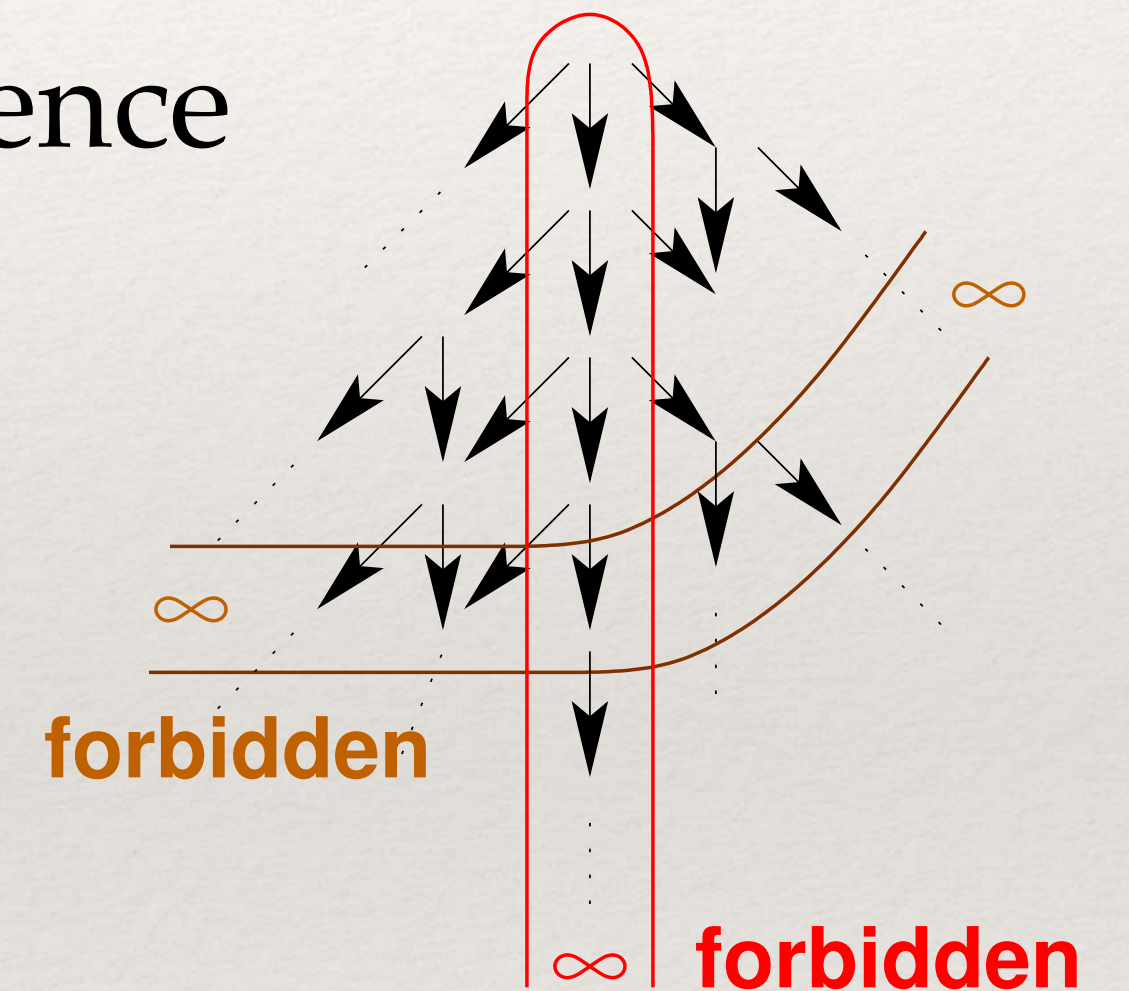
Noetherian spaces and wqos

Well-quasi-orders

- ❖ **Fact.** The following are equivalent for a quasi-ordering \leq :
 - (1) Every sequence $(x_n)_{n \in \mathbb{N}}$ is **good**: $x_m \leq x_n$ for some $m < n$
 - (2) Every sequence $(x_n)_{n \in \mathbb{N}}$ is **perfect**: has a monotone subsequence
 - (3) \leq is **well-founded** and has **no infinite antichain**.

❖ **Defn.** Such a quasi-ordering \leq is called a **well-quasi-order (wqo)**.

- ❖ Applications:
 - classification of graphs (Kuratowski, Robertson-Seymour)
 - verification (computer science)
 - model theory (logic: Fraïssé, Jullien, Pouzet)



Examples

- ❖ \mathbb{N} , with its usual ordering — More generally, any **total well-founded** order
- ❖ Every **finite** set, with any quasi-ordering
- ❖ Finite disjoint **sums**, finite **products** of wqos are wqo
- ❖ **Images** of wqos by monotonic maps are wqo (in particular quotients)
- ❖ **Inverse images** of wqos by order-reflecting maps are wqo (in particular subsets)
- ❖ **Higman's Lemma.** Let $X^* = \{\text{finite words over alphabet } X\}$
ordered by word embedding \leq_* . Then $X \text{ wqo} \Leftrightarrow X^* \text{ wqo}$
- ❖ **Kruskal's Theorem.** Let $\mathcal{T}(X) = \{\text{finite trees with } X\text{-labeled vertices}\}$
ordered by homeomorphic tree embedding \leq_T . Then $X \text{ wqo} \Leftrightarrow \mathcal{T}(X) \text{ wqo}$.
- ❖ And so on.

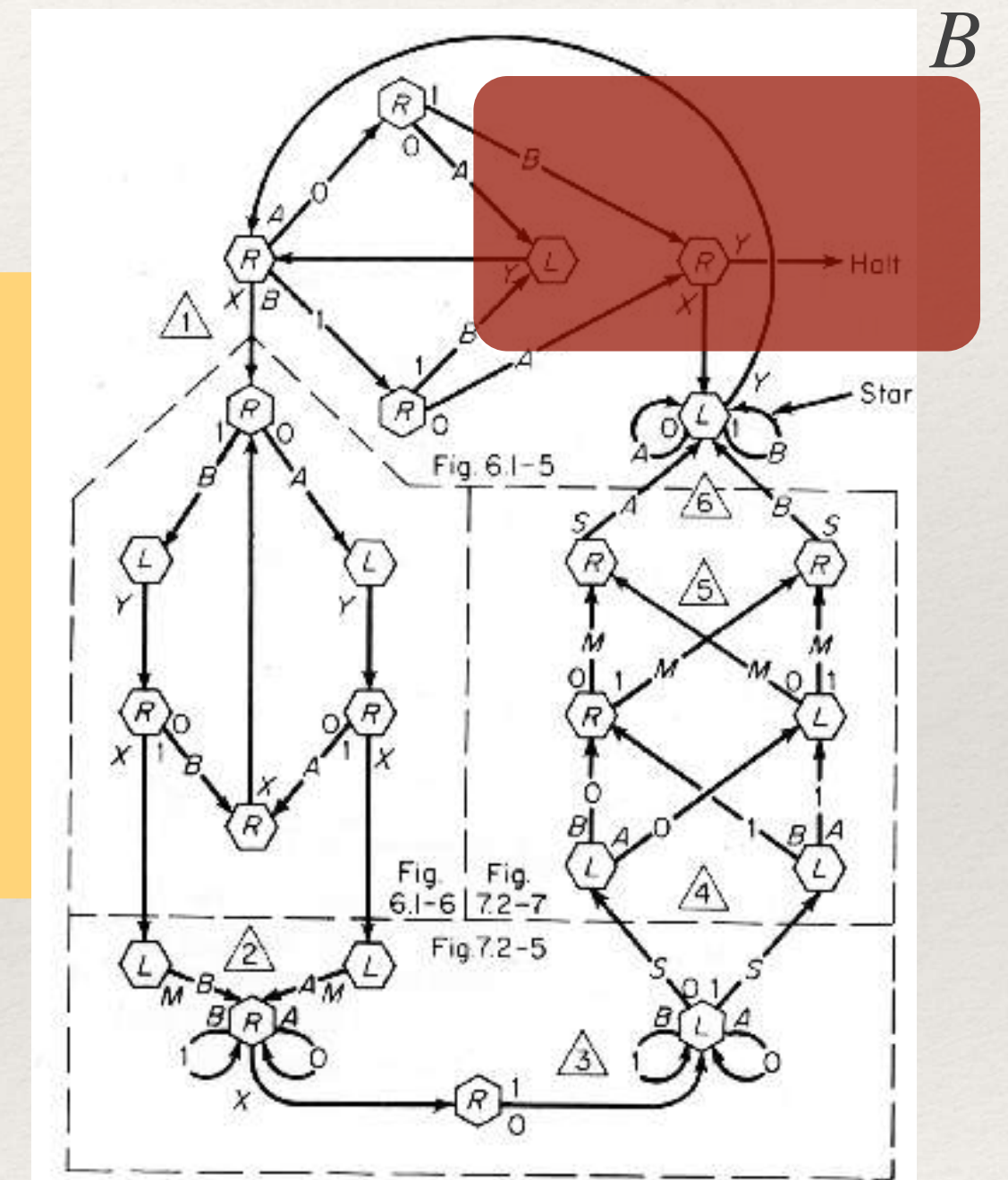
A computer scientist's view

Transition systems

- ❖ A **transition system** is just a directed graph (not necessarily finite)
Vertices are **configurations** (of a computer system, say)
Computation proceeds in steps $C \rightarrow C'$ (along edges)

- ❖ **Reachability:** Given a starting configuration C_0 ,
and a set B of configurations,
can one reach B from C_0 ?
(i.e., is there a $C \in B$ such that $C_0 \rightarrow^* C$?)

- ❖ Decidable for finite transition systems (in polynomial time)
- ❖ In general **undecidable**: consider the graph of configurations
of a universal Turing machine, $B = \{\text{accepting configurations}\}$

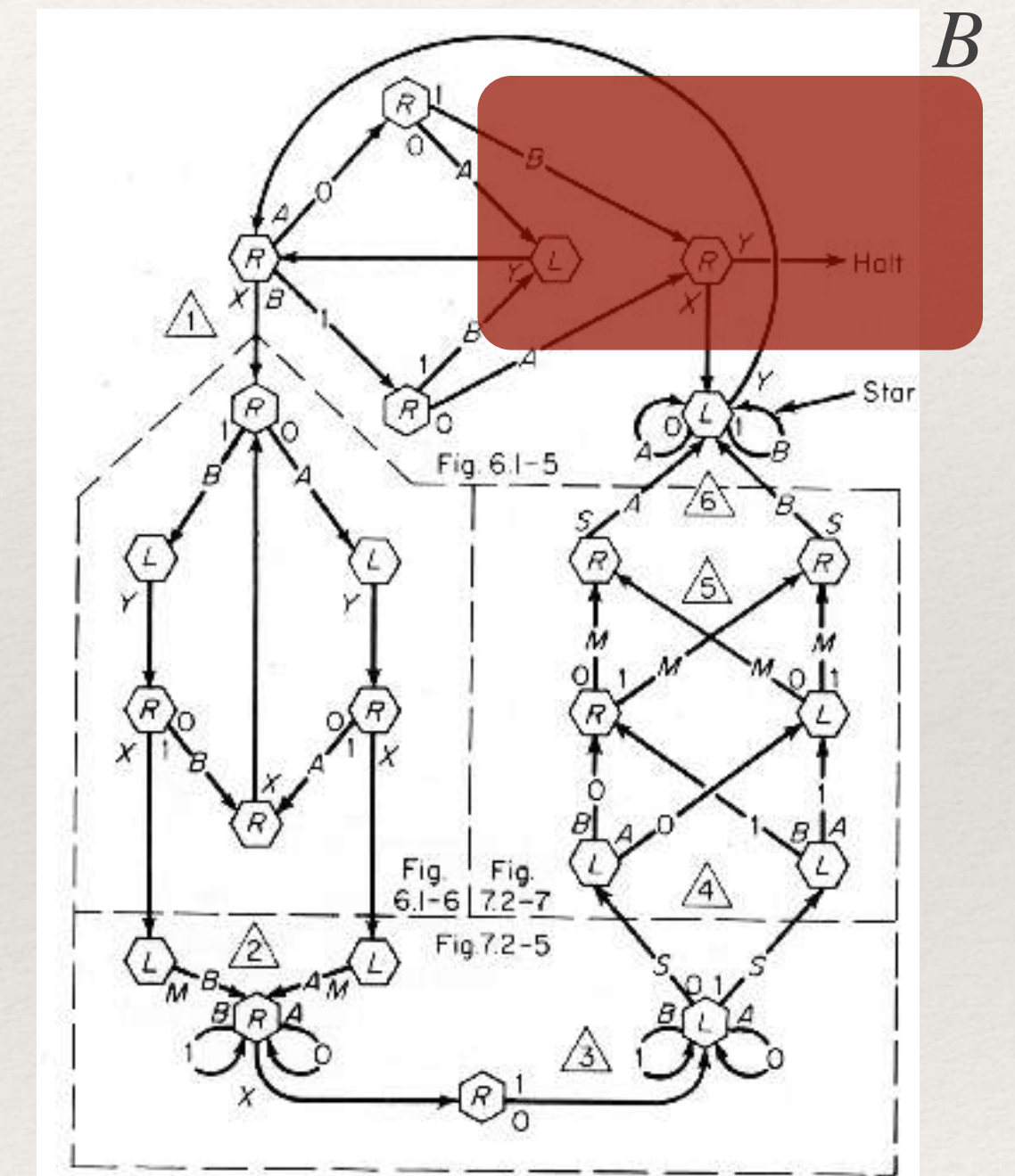


Verification

- ❖ In practice, a transition system is a model of some computer system (e.g., a program)
- ❖ and B is the set of **bad configurations**, typically where some property of interest is violated.

Illustration:

$B =$

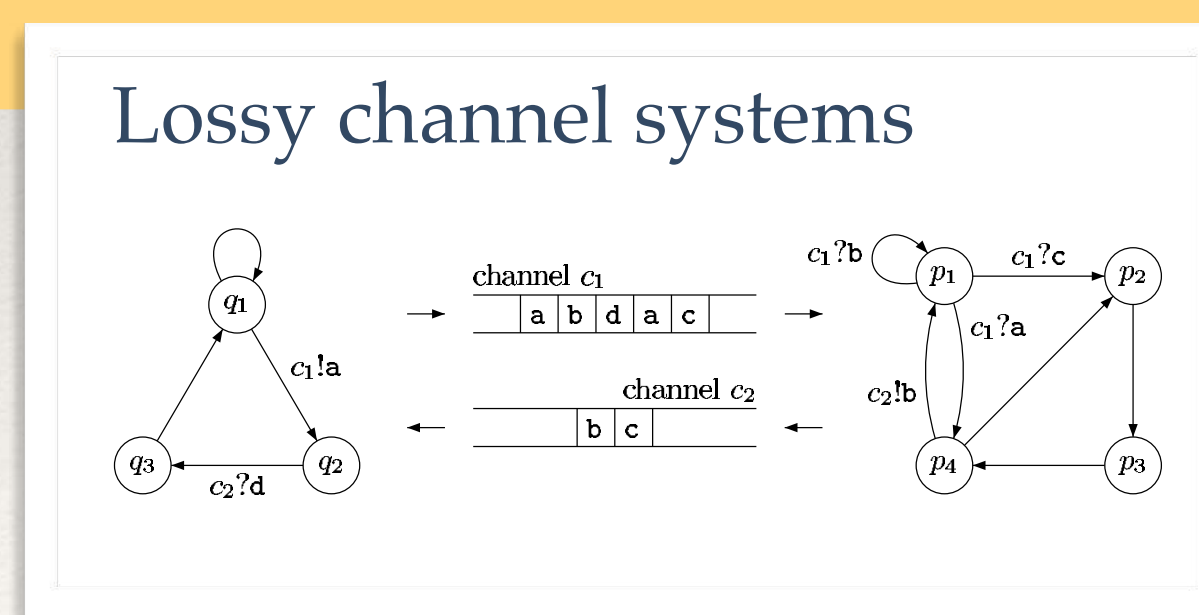
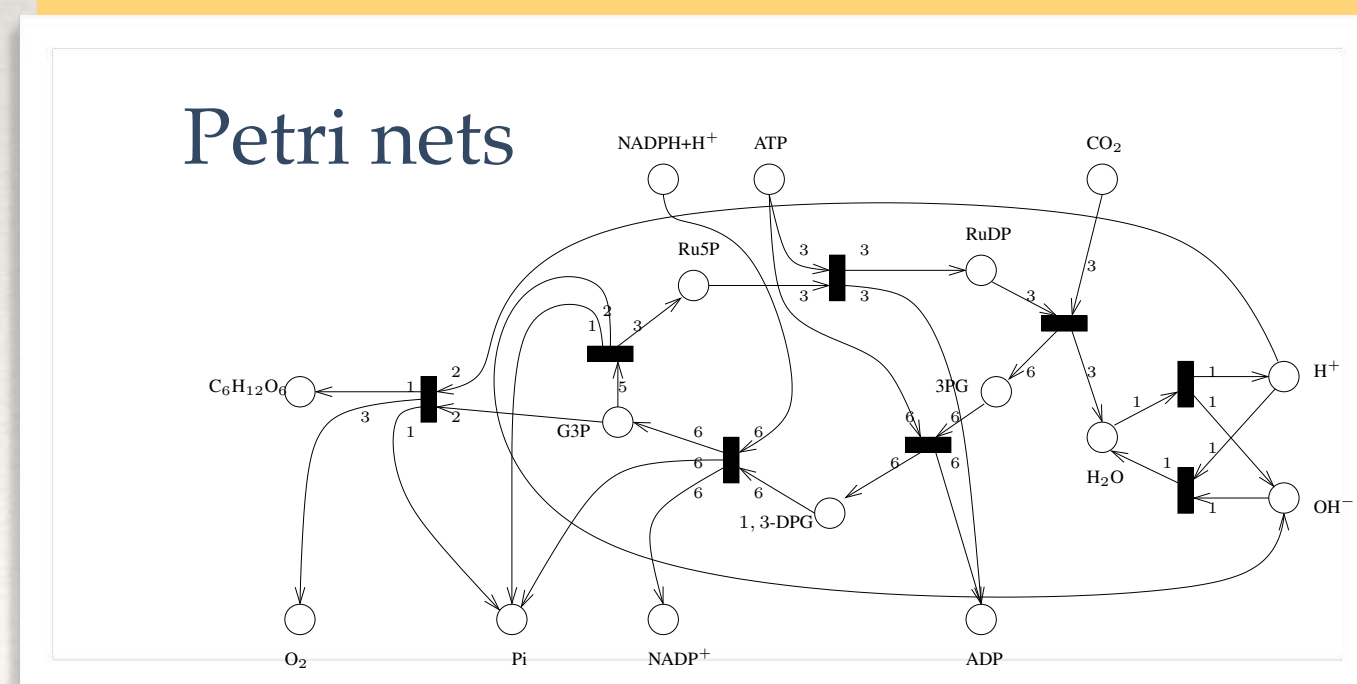
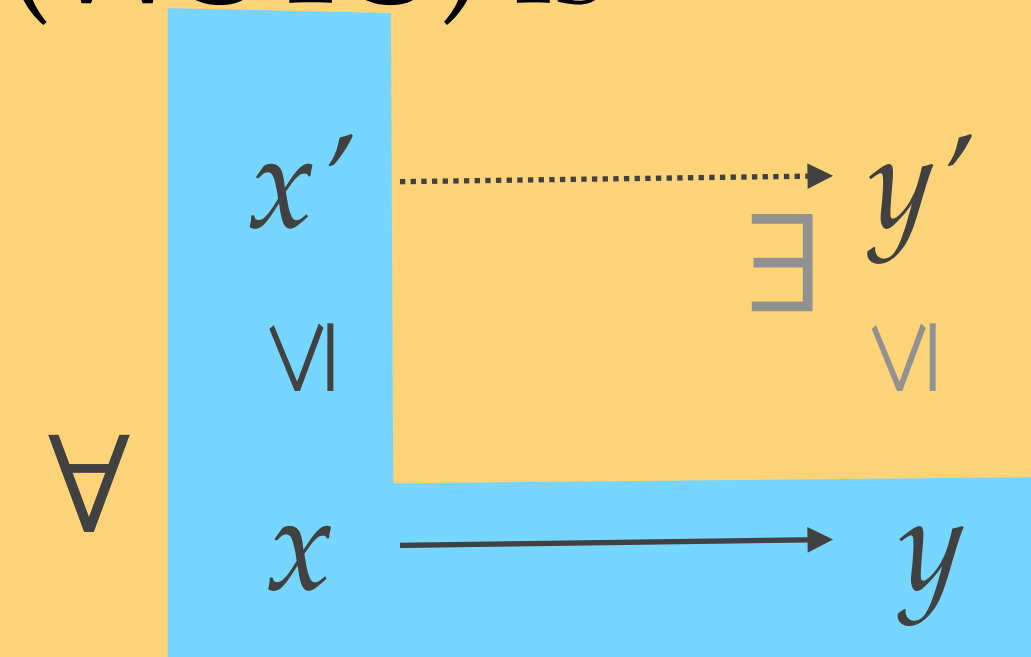


Well-structured transition systems

(Abdulla, Čerāns, Jonsson & Tsay 2000,
Finkel & Schnoebelen 2001)

- ❖ A very interesting class of (infinite) transition systems where **coverability** (a special form of reachability) is **decidable**

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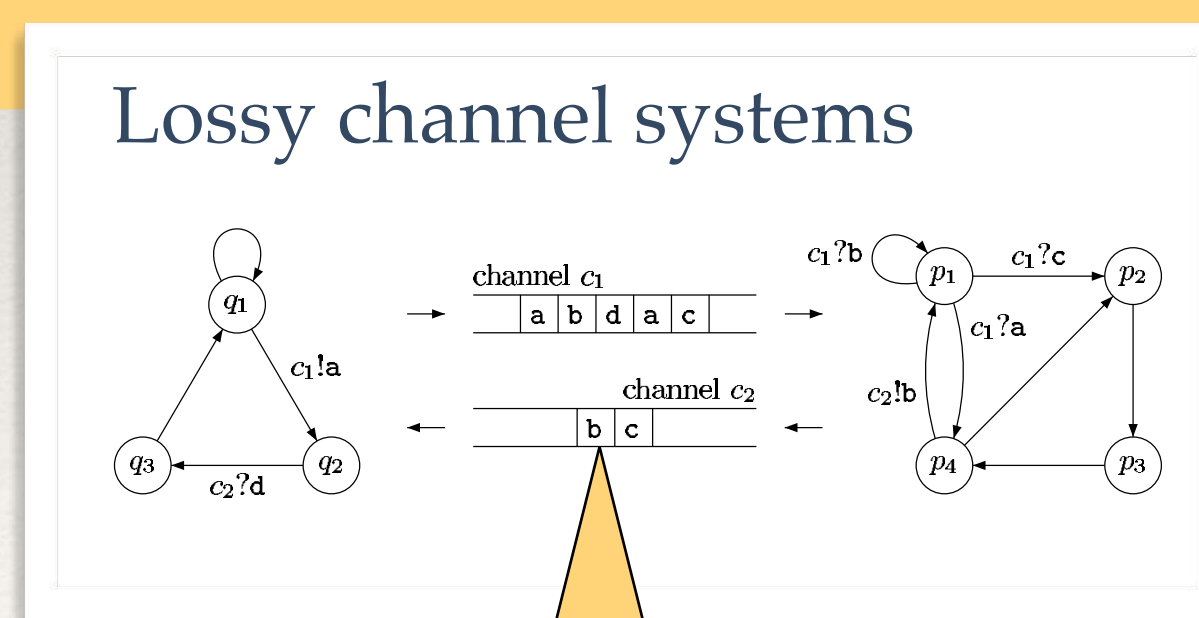
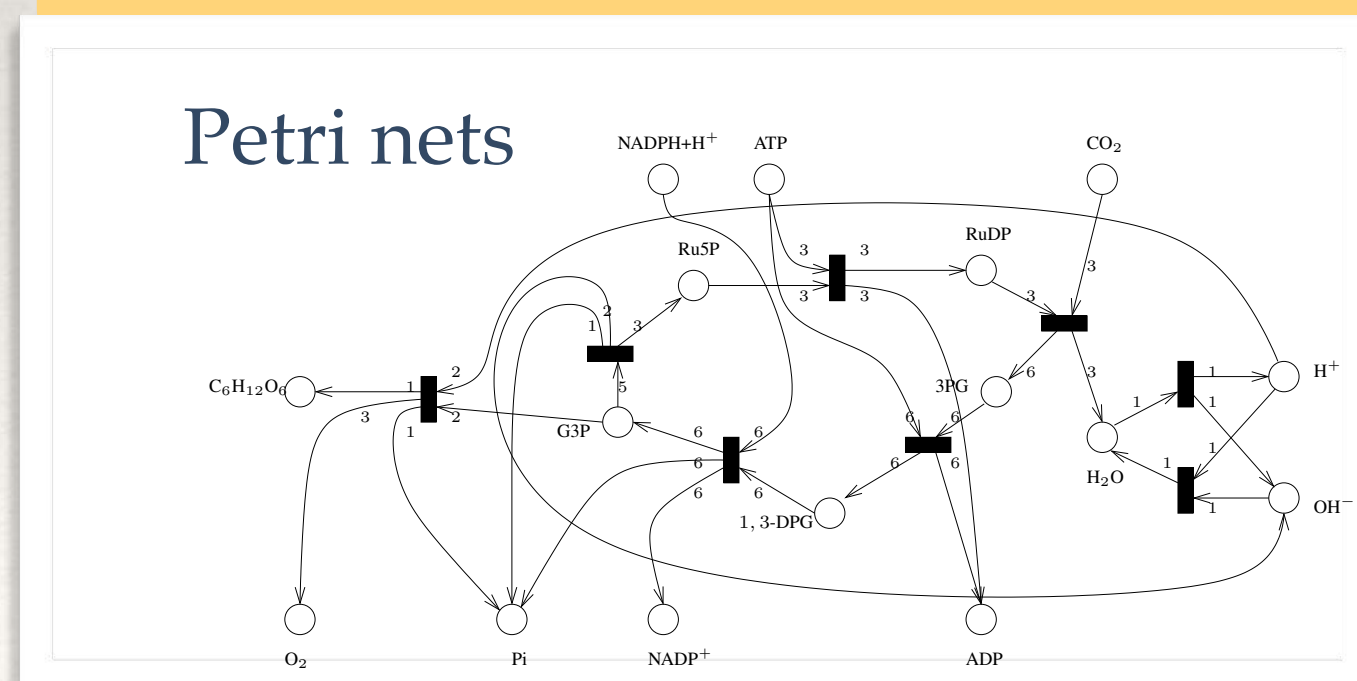
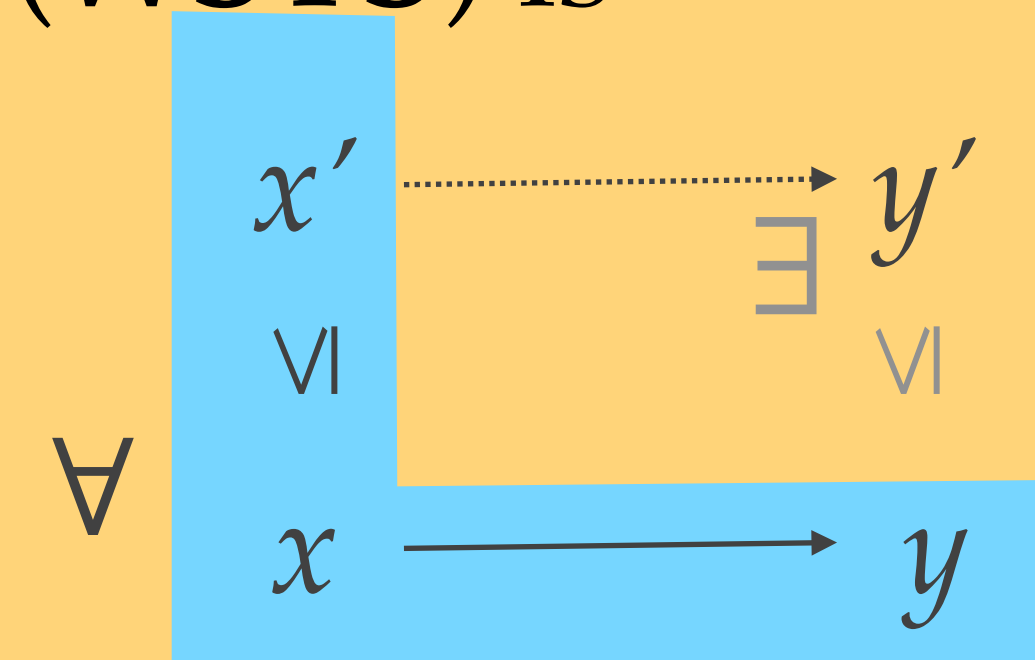
... and many other examples

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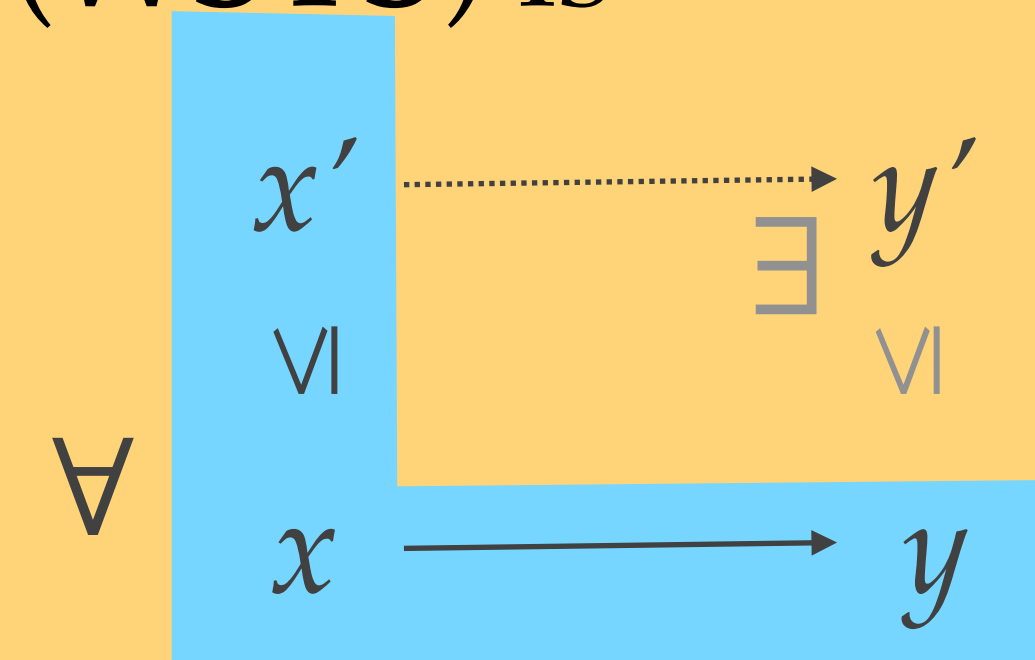
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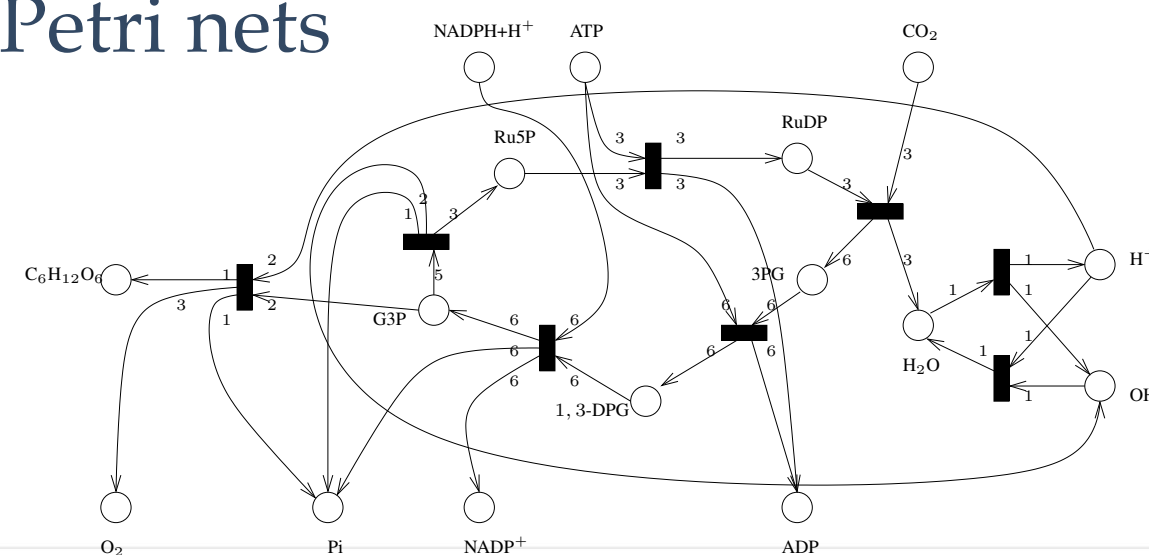
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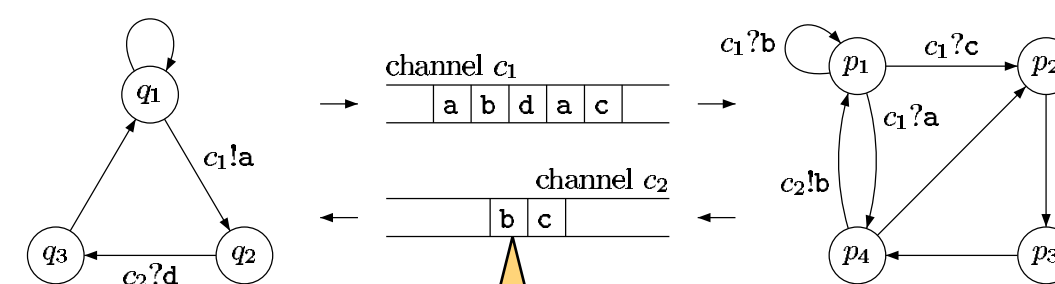
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Petri nets



Lossy channel systems



- ❖ In order to understand them, we need...

... and many letters can spontaneously vanish from communication queues (needed for decidability... and rather realistic)

More on well-quasi-orders

Wqos

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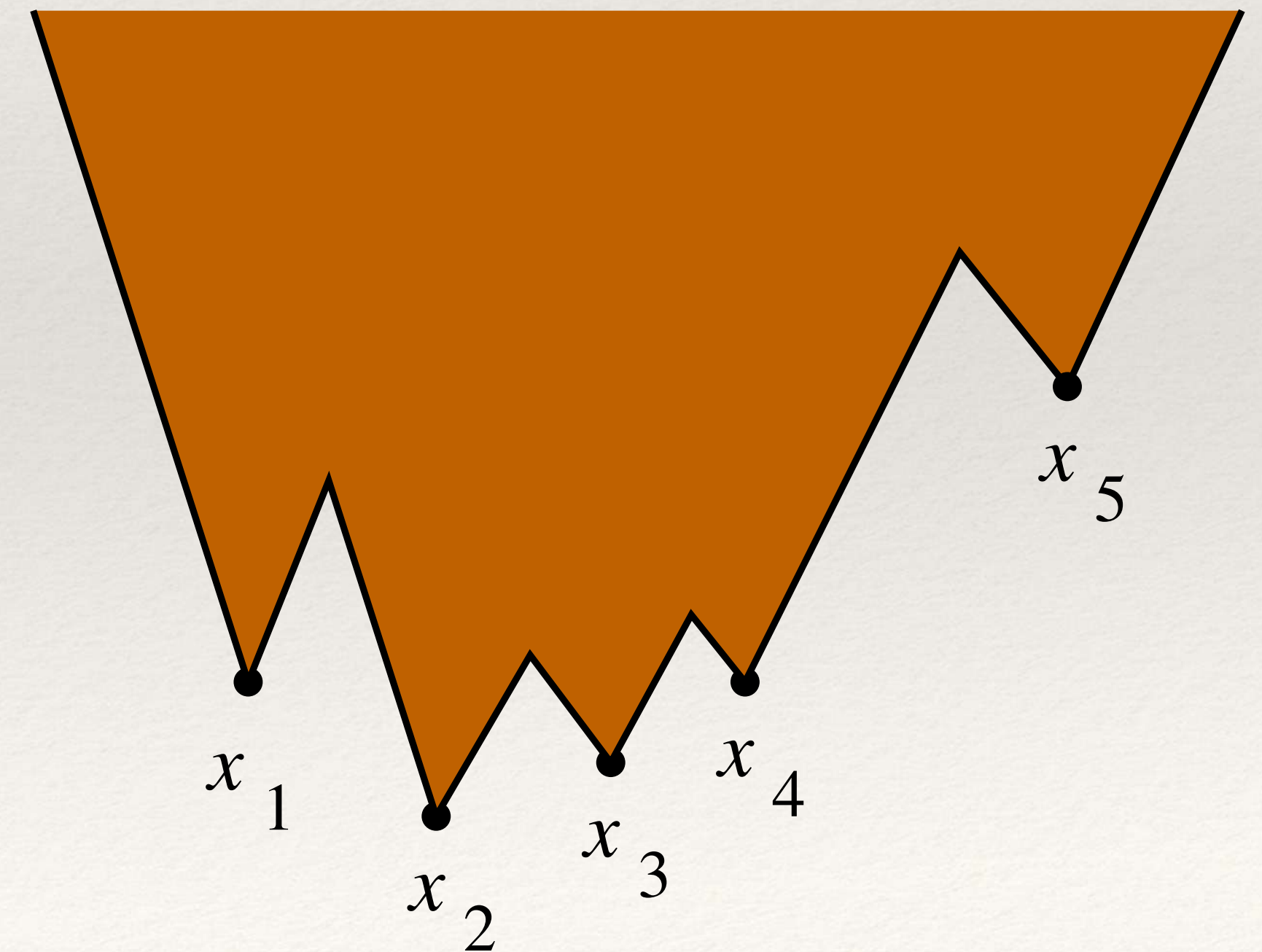
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- (4) Every **upwards-closed** subset is the upwards-closure $\uparrow \{x_1, \dots, x_n\}$ of a **finite set**



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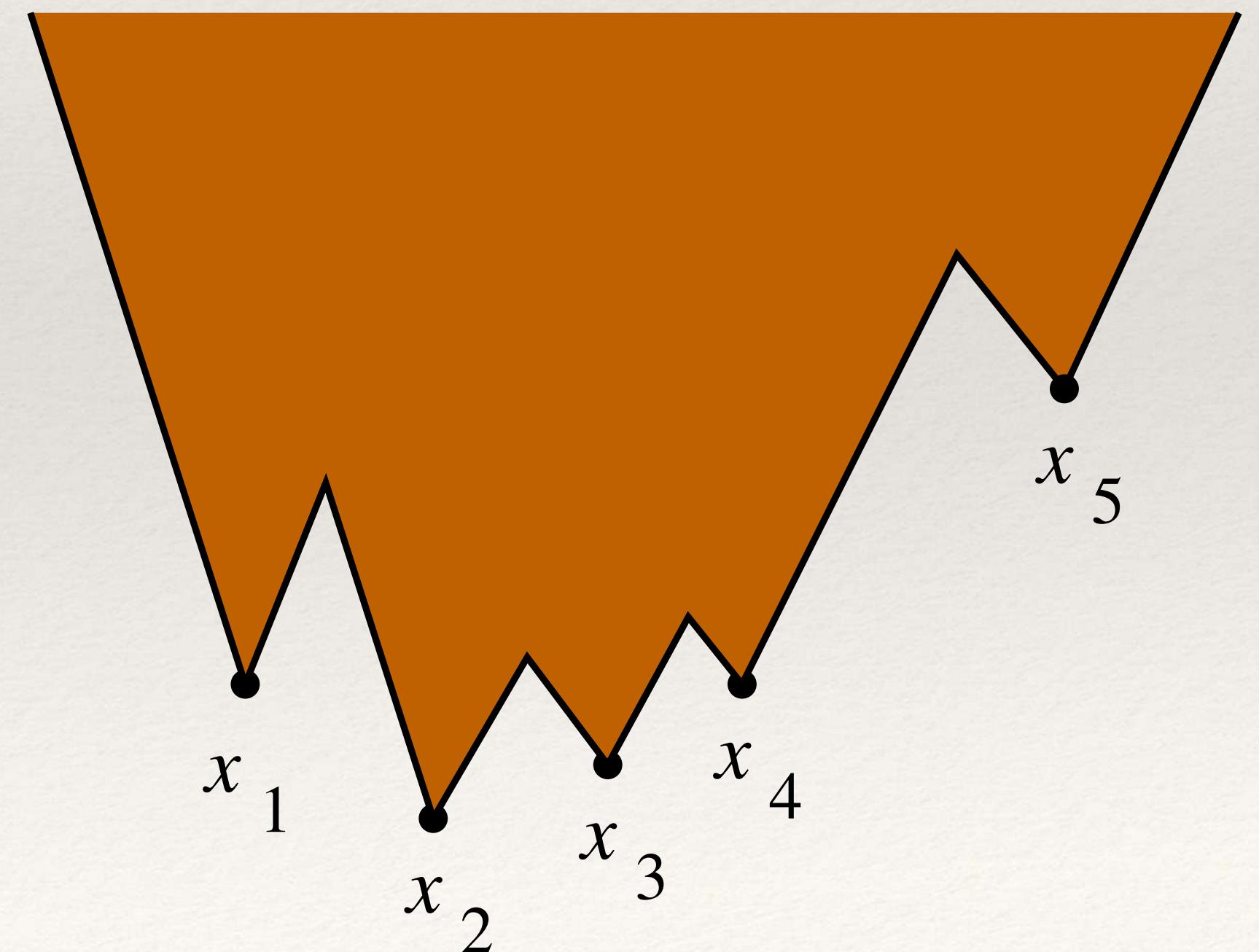
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If $x \leq y$ and x is in the set, then so is y



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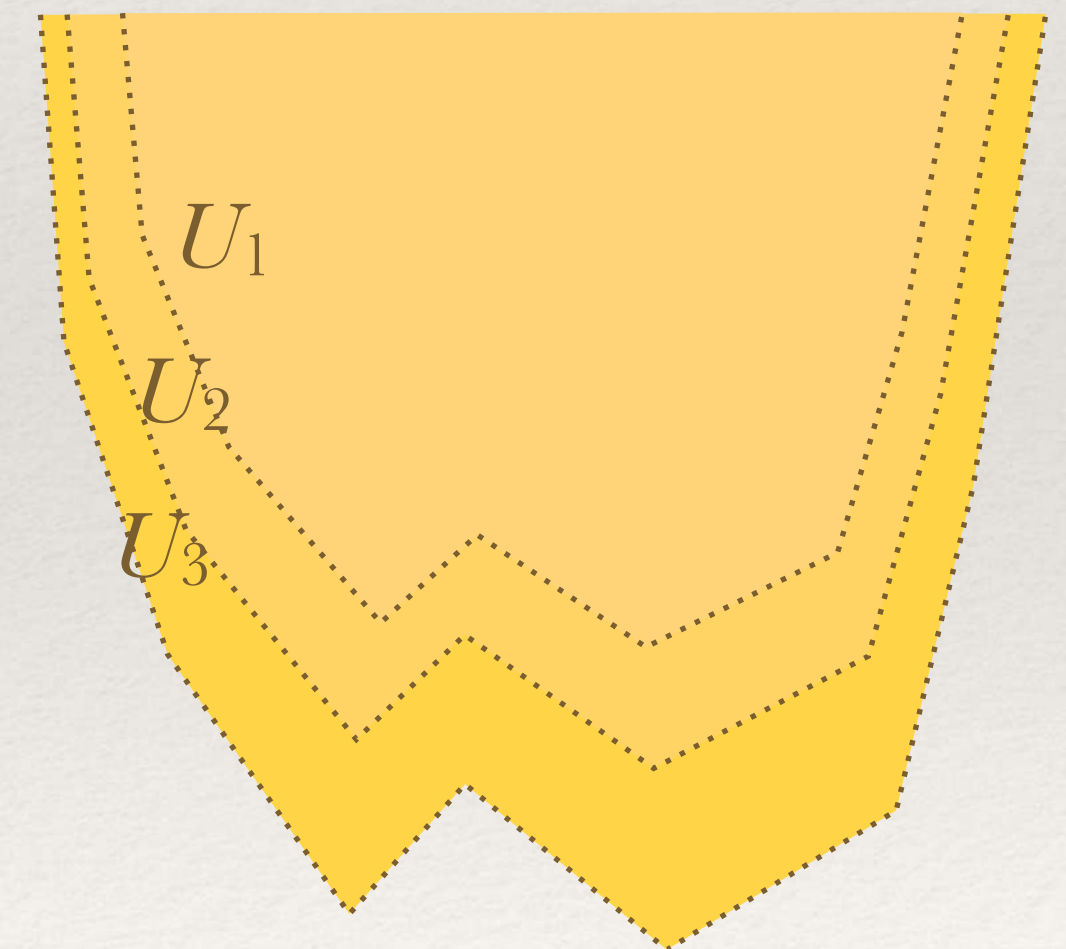
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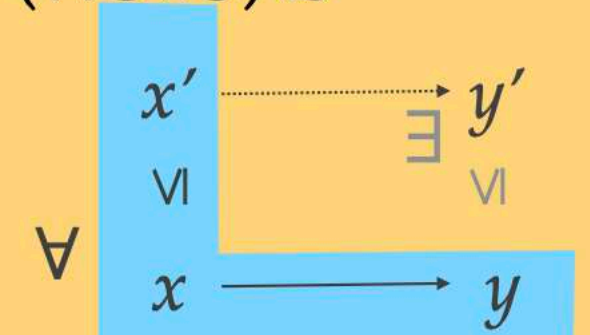
(i.e., all the sets U_n are equal from some rank on)



Coverability

- ❖ **Coverability** is the special case of reachability where the set B is **upwards-closed**

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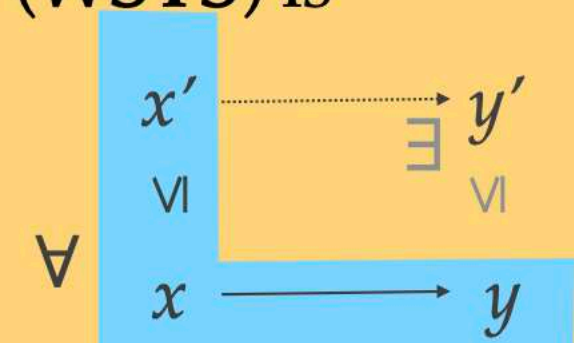
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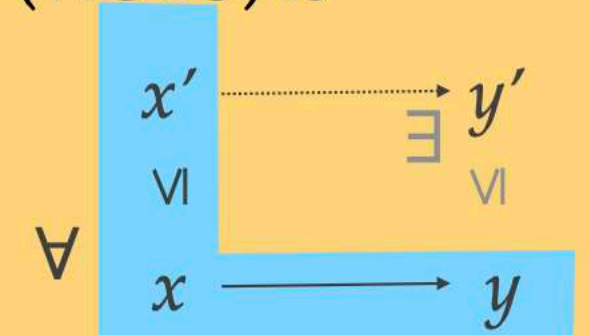
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Then $\text{Pre}^{\leq 0}(U) \subseteq \text{Pre}^{\leq 1}(U) \subseteq \dots \subseteq \text{Pre}^{\leq n}(U) \subseteq \dots$

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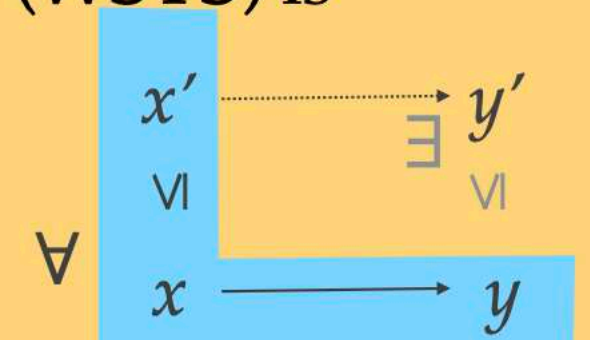
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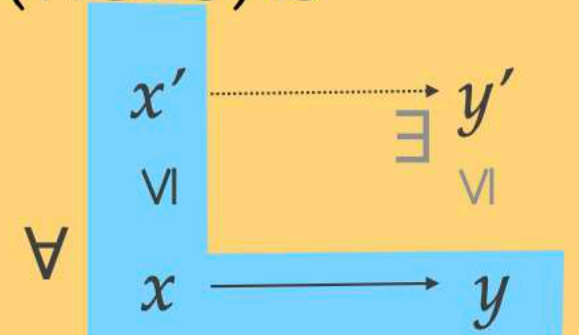
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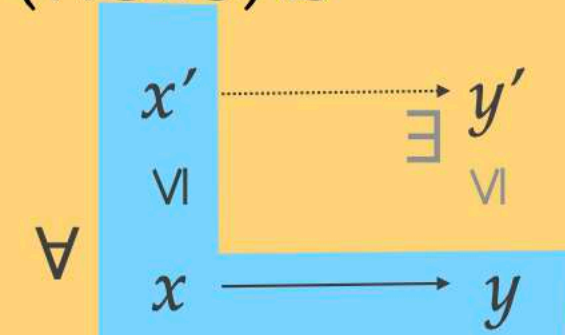
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- ❖ Now note that B is reachable from C_0 iff $C_0 \in \text{Pre}^*(B)$

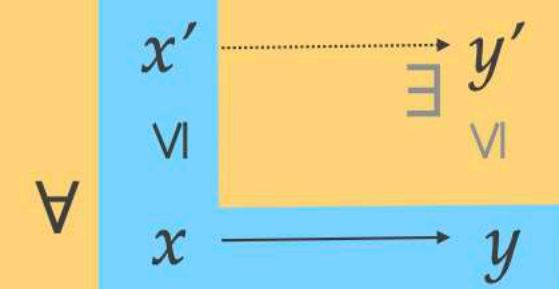
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Coverability is decidable

- ❖ In order to make this argument precise, we really need to reason with **effective** WSTSs, where
 - points are representable
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 - $y \mapsto \{x_1, \dots, x_n\} = \text{Pre}(\uparrow y)$ is computable (so one can compute $\text{Pre}(U)$)

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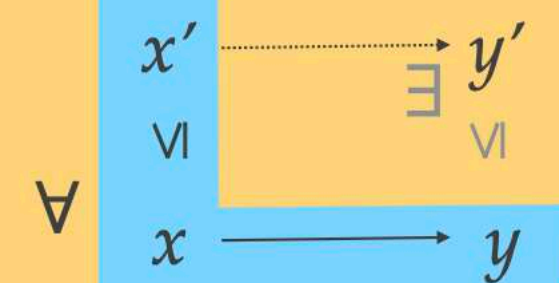
- ❖ **Theorem.** (Abdulla et al. 2000, Finkel&Schnoebelen 2001.)
Coverability is **decidable** on effective WSTSs.

```
fun pre* U =  
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  in  
    if V ⊆ U  
    then U  
    else pre* (U ∪ V)  
  end;  
  
fun coverability (s, B) =  
  s in pre* (B);
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- ❖ Complexity: appalling (**EXPSPACE**-complete for Petri nets, grows faster than Ackermann for lossy channel systems)

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Beyond wqos: Noetherian spaces

Going topological

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X is wqo iff:

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- ❖ **Proposition.** (X, \leq) is wqo iff X is **Noetherian** in its Alexandroff topology.

- ❖ Hence Noetherian spaces generalize wqos

Is the generalization proper?

- ❖ **Yes.** Consider \mathbb{N}_{cof} , the set of natural numbers with the **cofinite topology**, whose closed sets are the finite subsets (plus \mathbb{N})

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❖ Oh, wait, why does \mathbb{N}_{cof} **not** arise from a wqo?

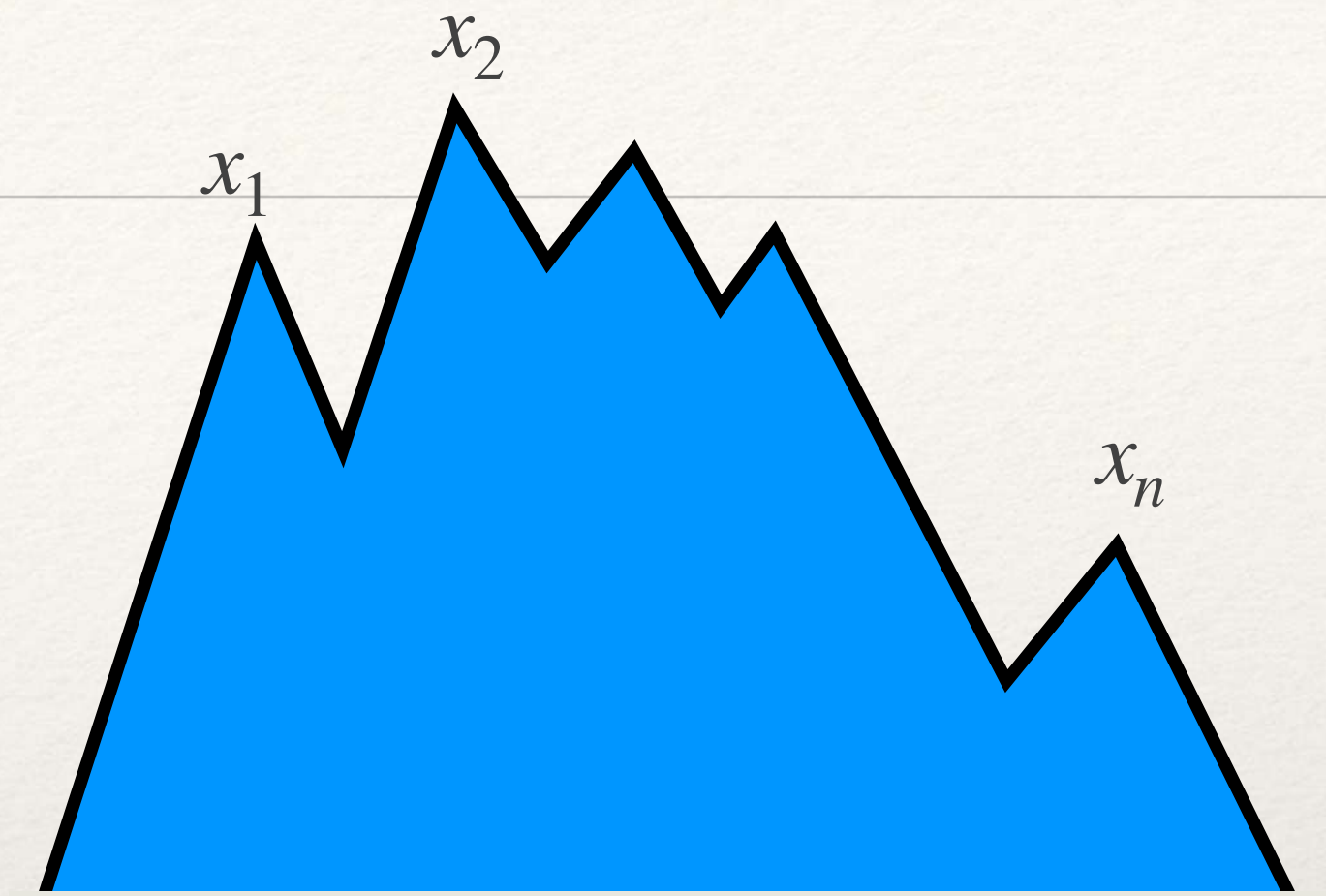
The specialization quasi-ordering

- ❖ Every topological space X has a **specialization quasi-ordering**:
 $x \leq y$ iff every open neighborhood of x contains y
iff x is in the closure of $\{y\}$
- ❖ The specialization quasi-ordering of $(X$ in the Alexandroff topology of \leq) is \leq
- ❖ The specialization quasi-ordering of \mathbb{N}_{cof} is **equality** (\mathbb{N}_{cof} is T_1)
and equality is **never** a wqo on an infinite set

So \mathbb{N}_{cof} is a Noetherian space
that does **not** arise from a wqo

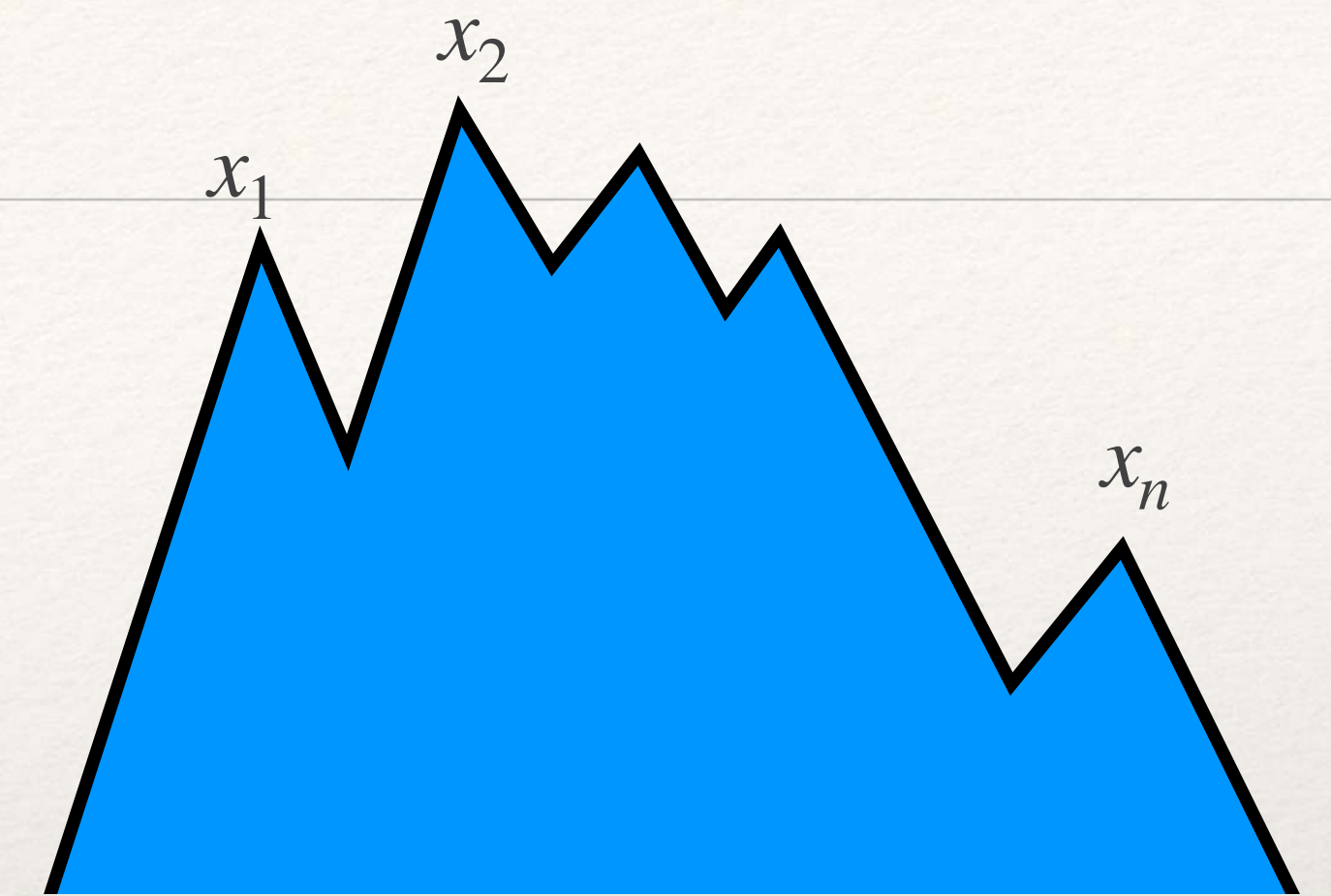
Properties T and W

- ❖ Let (X, \leq) be a quasi-ordered set.
Its **finitary** subsets are $\downarrow \{x_1, \dots, x_n\}$
- ❖ The finitary subsets generate the **upper topology**
It, too, has \leq as specialization quasi-ordering



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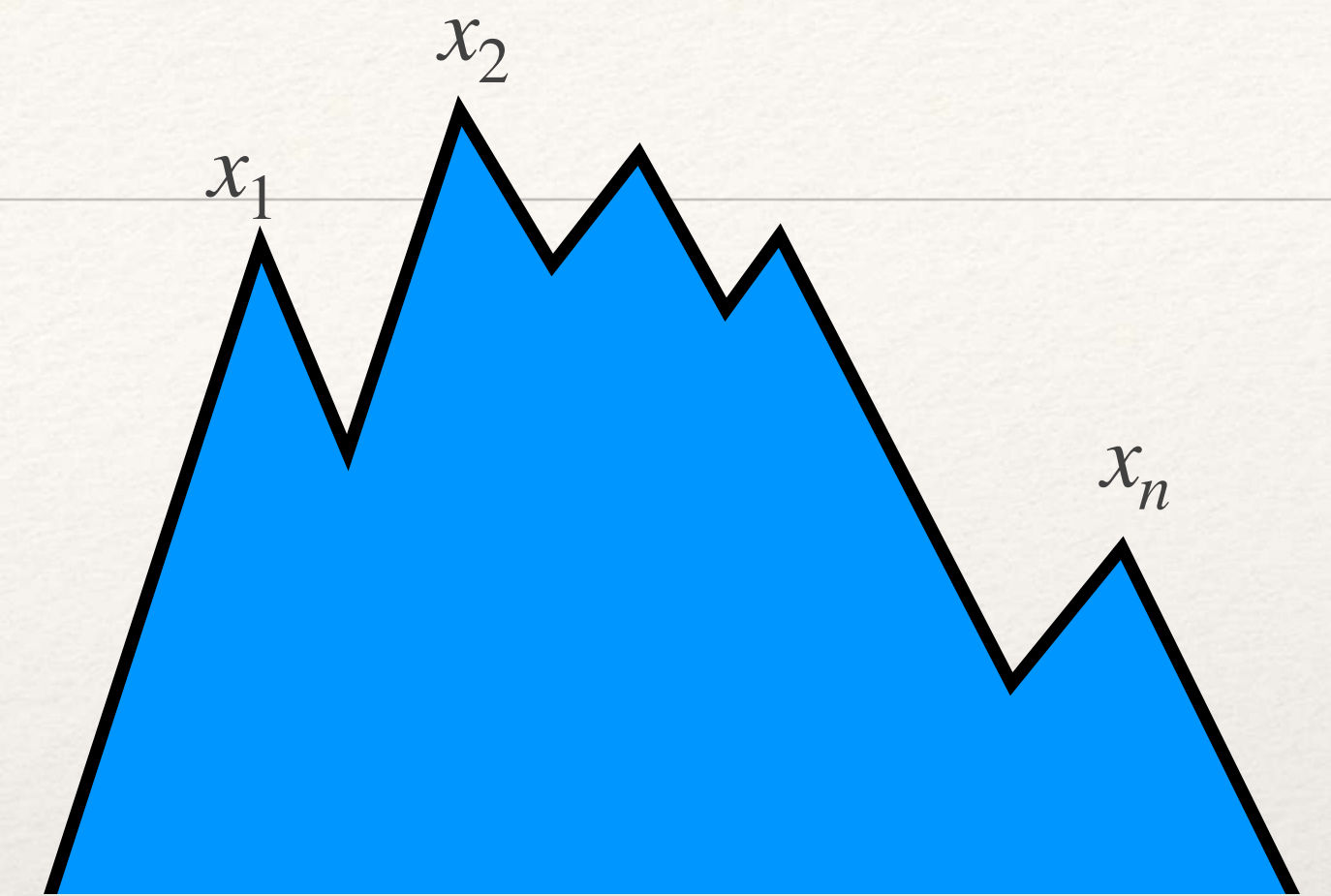
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The upper topology is the **coarsest** topology
with \leq as specialization
The Alexandroff topology is the **finest**.

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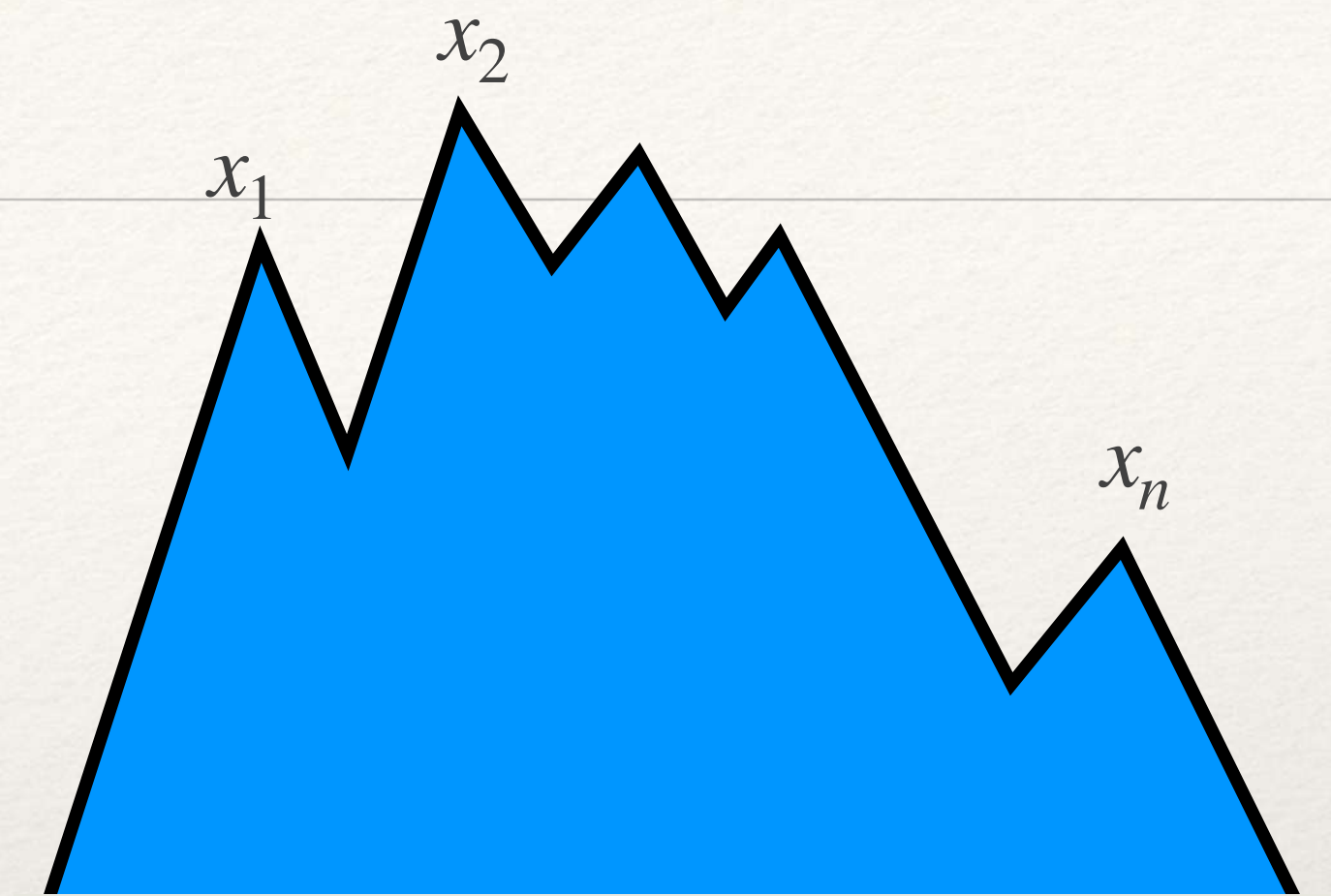


- ❖ **Proposition.** If:
 - X is **well-founded**
 - (Property T) X is **finitary**
 - (Property W) For all $x, y \in X$, $\downarrow x \cap \downarrow y$ is **finitary**then X is **Noetherian** in the upper topology
and the **closed** sets are the **finitary subsets**.

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This turns out to be the general form
of all **sober** Noetherian spaces.

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❖ Let $\mathcal{H}X = \{\text{closed subsets of } X\}$ with the upper topology of \subseteq (Hoare hyperspace of X)

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Proposition. If X is Noetherian, then $\mathcal{H}X$ is Noetherian.

(That is actually an equivalence.)

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- ❖ Equip $\mathbb{P}(X)$ with the **lower Vietoris topology**,
Subbase of closed sets $\square C = \{A \in \mathbb{P}(X) \mid A \subseteq C\}$,
 C closed in X
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Yes, Noetherianness is a **localic** property.
Category of **sober Noetherian** spaces \cong
Locales with **no infinite monotonic chain**

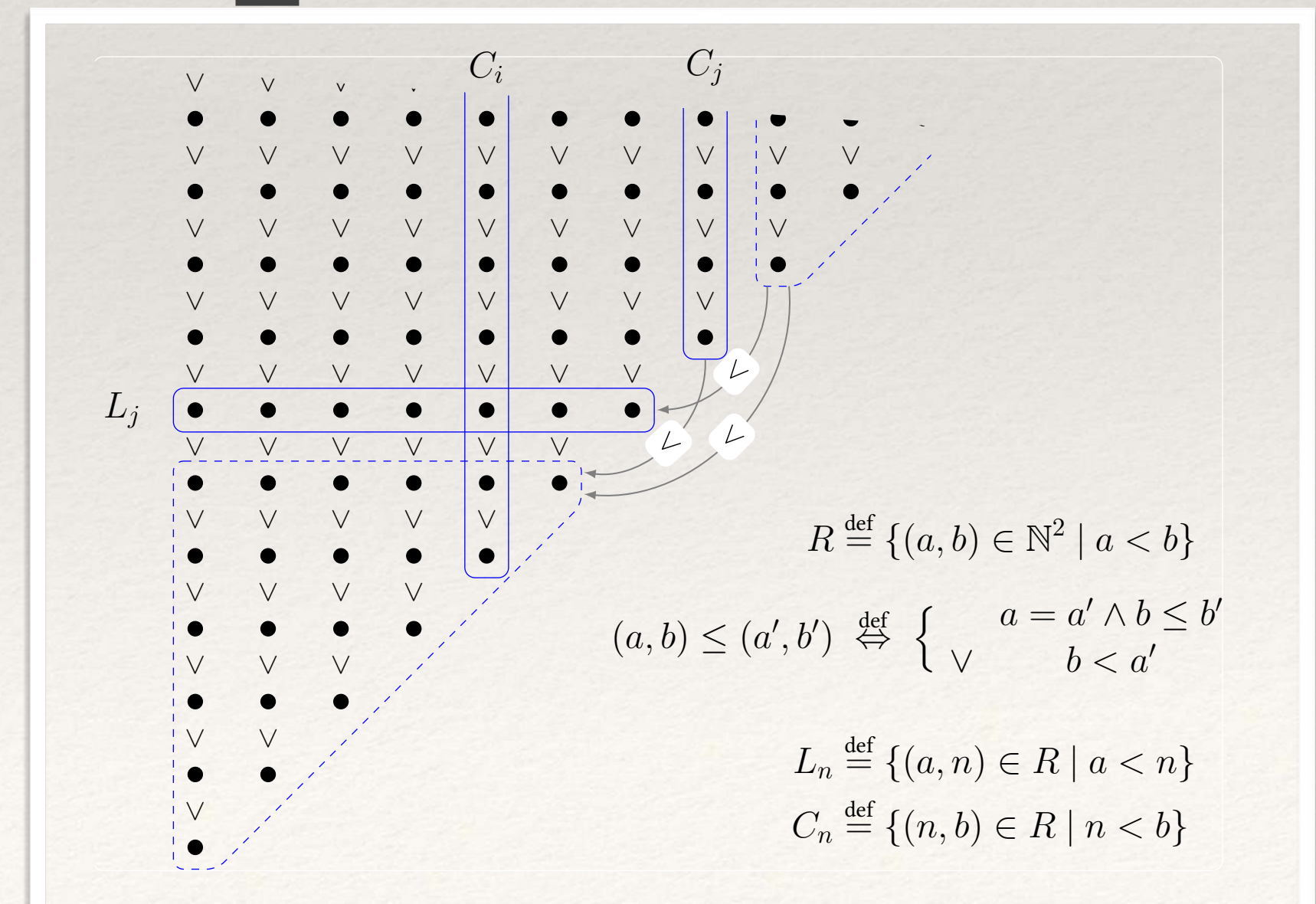
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When X is a wqo (with the Alexandroff topology),

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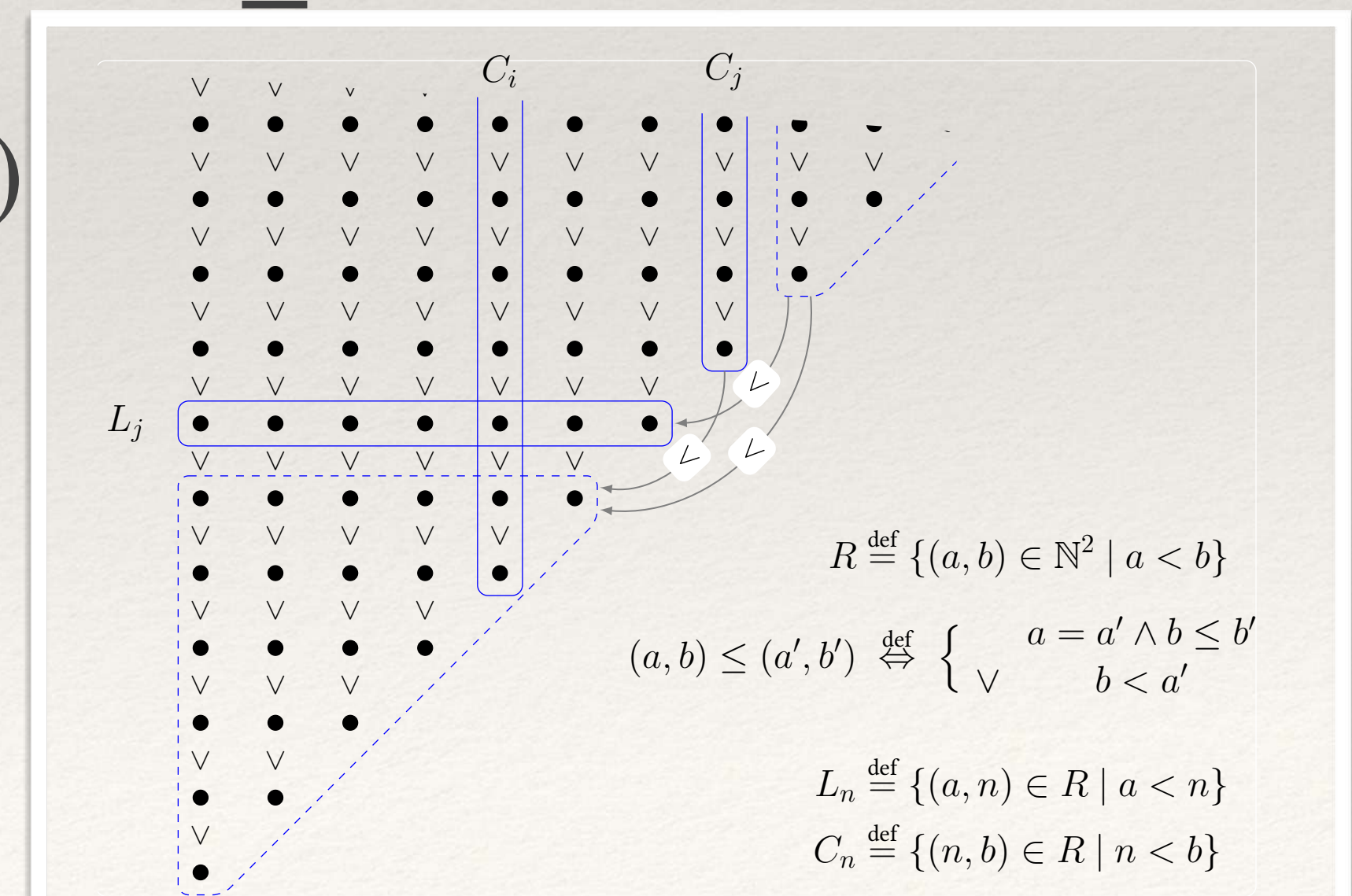
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- ❖ But $(\mathbb{P}(X), \leq^b)$ is **not wqo** for general wqos (X, \leq) (Rado, 1957)



Finite words

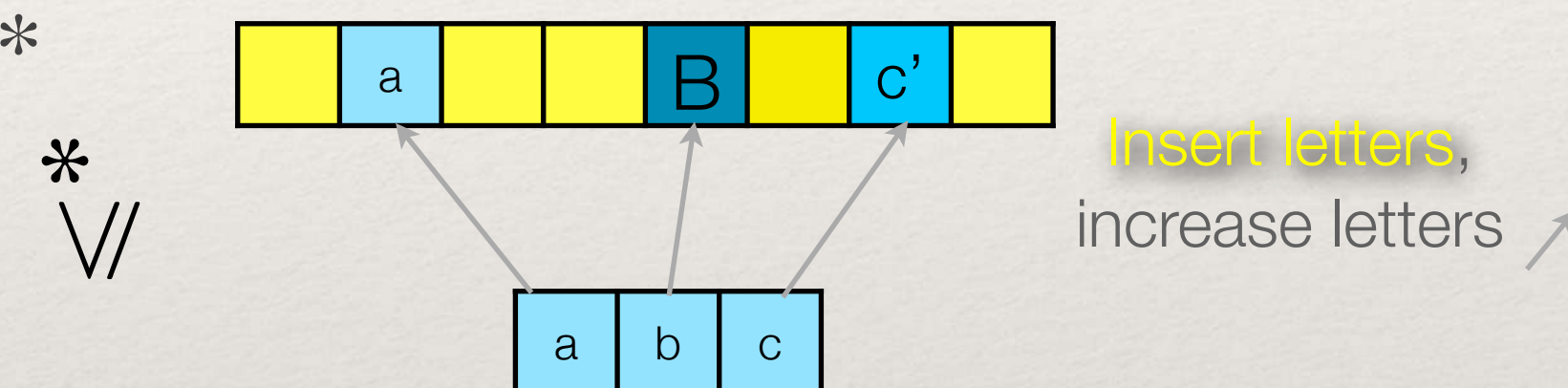
- ❖ Let $X^* = \{\text{finite words on } X\}$ with **word topology**:
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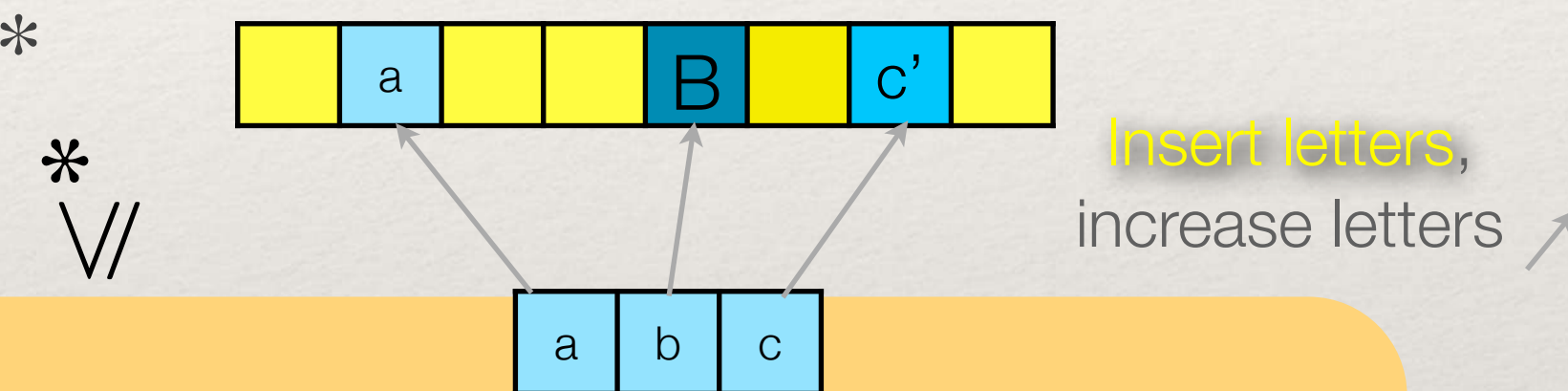
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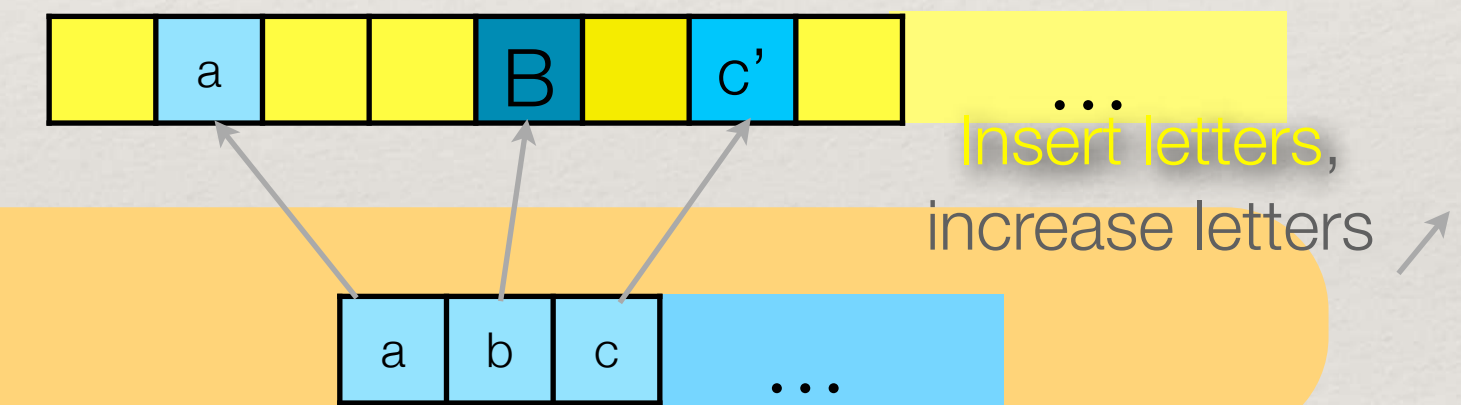


- Theorem (JGL 2013).** X Noetherian iff X^* Noetherian

Generalizes Higman's Lemma (Higman 1952): X wqo iff X^* wqo

Infinite words

- Let $X^{\leq\omega} = \{\text{finite or infinite words on } X\}$ with **asymptotic word topology**:
 subbasic open sets $\langle U_1, \dots, U_n \rangle = X^* U_1 X^* \dots X^* U_n X^{\leq\omega}$,
 and $\langle U_1, \dots, U_n; (\infty)V \rangle = X^* U_1 X^* \dots X^* U_n (X^* V)^\omega$ (U_i, V open in X)
- Specialization quasi-ordering is (infinite) **word embedding**



- Theorem (JGL 2021).** X Noetherian iff $X^{\leq\omega}$ Noetherian
 No equivalent in wqo theory — except if you adopt bqo theory.

Transfinite words

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- ❖ **Theorem** (JGL, Halfon, Lopez 2022, submitted).

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No equivalent in wqo theory — except if you adopt bqo theory...

(Warning: specialization \neq word embedding in general.)

Topological WSTS

-
-
- ❖ So Noetherian spaces go beyond wqos,
but do they have any use?
 - ❖ Of course they do: a reminder of where they come from
 - ❖ An application in verification

The origin of Noetherian spaces

- ❖ The **spectrum** $\text{Spec}(R)$ of a ring R is the set of its **prime ideals** p
- ❖ with the **Zariski topology**, whose closed subsets are $\{p \in \text{Spec}(R) \mid I \subseteq p\}$, where I ranges over the ideals of R
- ❖ **Fact.** The spectrum of a Noetherian ring is Noetherian. (every monotone chain of ideals is stationary)
- ❖ In particular if $R = K[X_1, \dots, X_n]$ for some Noetherian ring, e.g., \mathbb{Z}
- ❖ One can **compute** with ideals, represented by **Gröbner bases** (Buchberger 1976)

An application of Gröbner bases in verification

- ❖ Verification of **polynomial programs**
(Müller-Olm&Seidl 2002)
- ❖ Propagates ideals of $\mathbb{Z}[X_1, \dots, X_n]$
backwards, as in the Pre^* algorithm
($X_1, \dots, X_n =$ variables of the program)
- ❖ Terminates because every monotonic chain $I_0 \subseteq I_1 \subseteq \dots \subseteq I_n \subseteq \dots$
of ideals is **stationary**
- ❖ ... very similar to Pre^* on WSTS, but
the (infinite) transition system underlying a polynomial program
is **not** a WSTS (inclusion between ideals **not** a wqo)

```
while (*) {  
  if (*) { x=2; y=3; }  
  else { x=3; y=2; }  
  x = x*y-6; y=0;  
  if (x2-3*x*y==0)  
    while (*) { x=x+1; y=y-1; };  
  x = x2+x*y;  
}
```

Topological WSTS

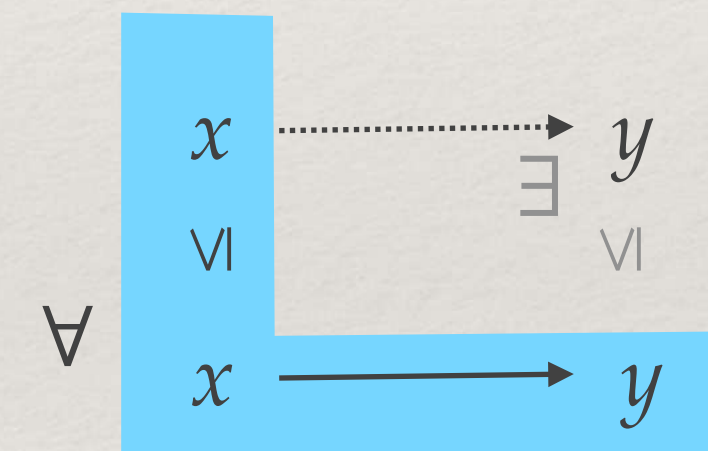
(JGL 2011)

❖ **Definition.** A **topological WSTS** is a transition system (X, \rightarrow) with a **Noetherian topology** \leq on X satisfying **lower semicontinuity**:
for every open subset U , $\text{Pre}(U)$ is open

❖ Namely, replace **wqo** by **Noetherian monotonicity** by **lower semicontinuity**

❖ If the topology is Alexandroff, then **Noetherian=wqo**,
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In particular, **every WSTS** is a topological WSTS



Topological WSTS

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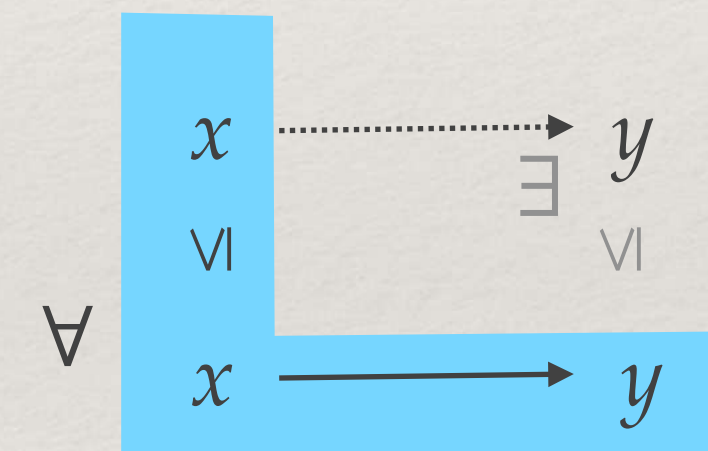
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❖ **Polynomial programs** are topological WSTS
— in the Zariski topology of $\text{Spec}(\mathbb{Z}[X_1, \dots, X_n])$



Topological coverability is decidable

- ❖ **Topological coverability:**

INPUT: an initial configuration x_0 ,
an **open set** U of bad configurations

QUESTION: is there a $x \in U$ such that $x_0 \rightarrow^* x$?

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- ❖ **Theorem (JGL 2011.)** Topological coverability is **decidable** on effective topological WSTSs.

- ❖ The algorithm is the same as with WSTSs.

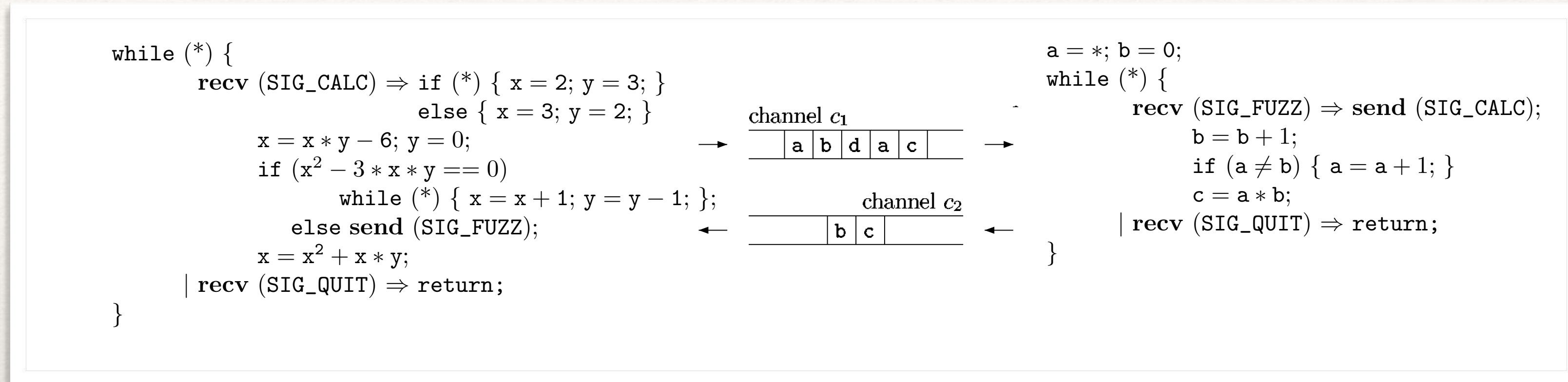
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```
fun pre* U =  
  let V = pre U  
  in  
    if V ⊆ U  
    then U  
    else pre* (U ∪ V)  
end;  
  
fun coverability (s, B) =  
  s in pre* (B);
```


Concurrent polynomial programs

(JGL 2011)

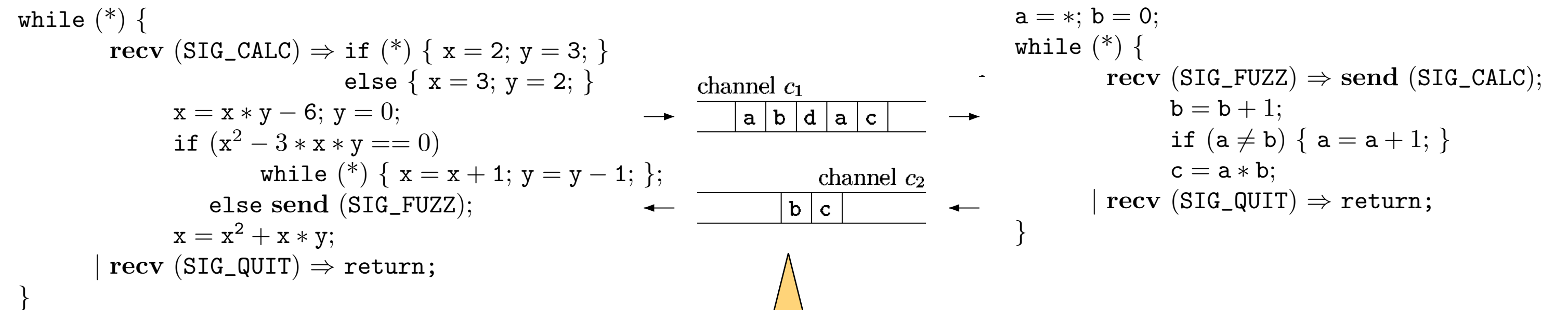
- ❖ Finite networks of polynomial programs P_1, \dots, P_m communicating through lossy communication queues on a finite alphabet Σ



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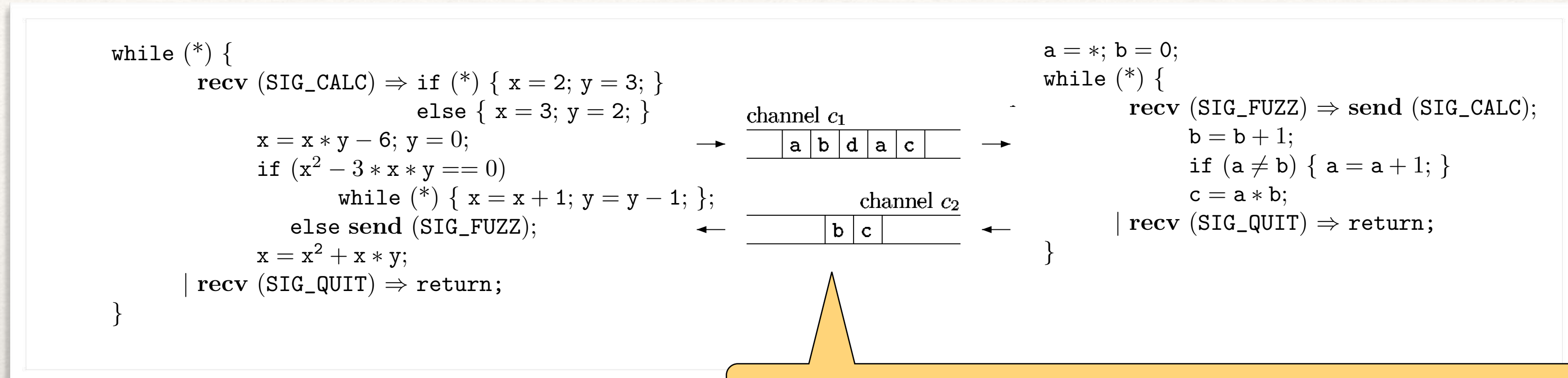
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- ❖ Finite networks of polynomial programs P_1, \dots, P_m communicating through **lossy** communication queues on a finite alphabet Σ

- ❖ State space = finite **product** of
 - spectra of polynomial rings $\mathbb{Z}[X_1, \dots, X_n]$, one for each P_i
 - Σ^* , with **word topology**, one for each communication queue

This is **Noetherian**, because:

- ❖ **Proposition.** Any finite product of Noetherian spaces is Noetherian.



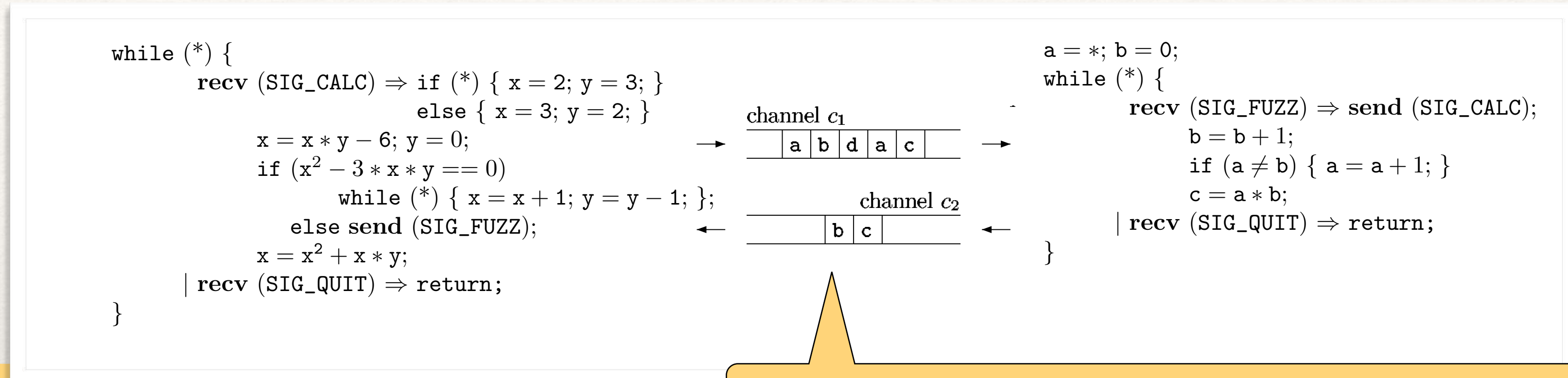
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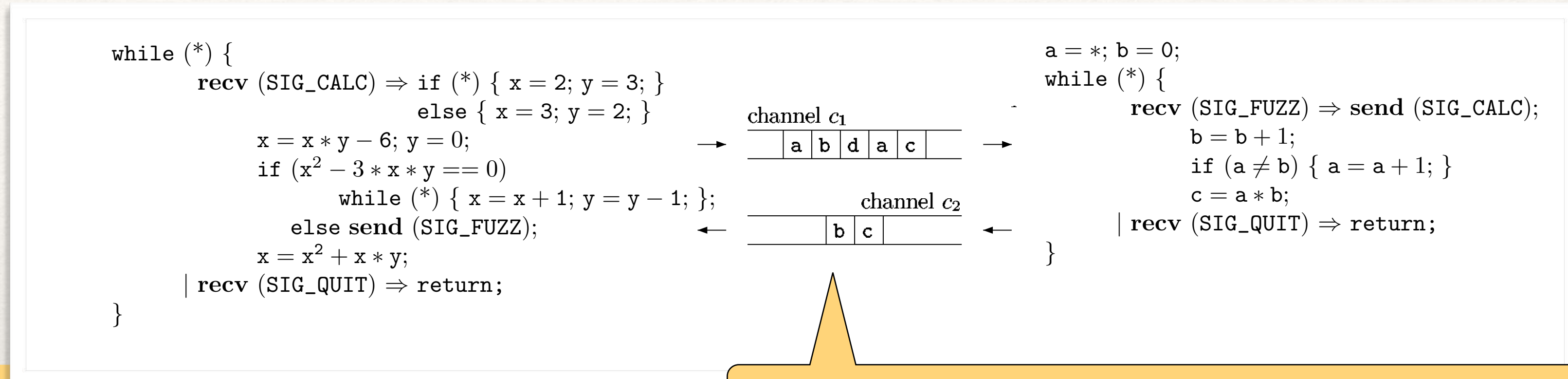


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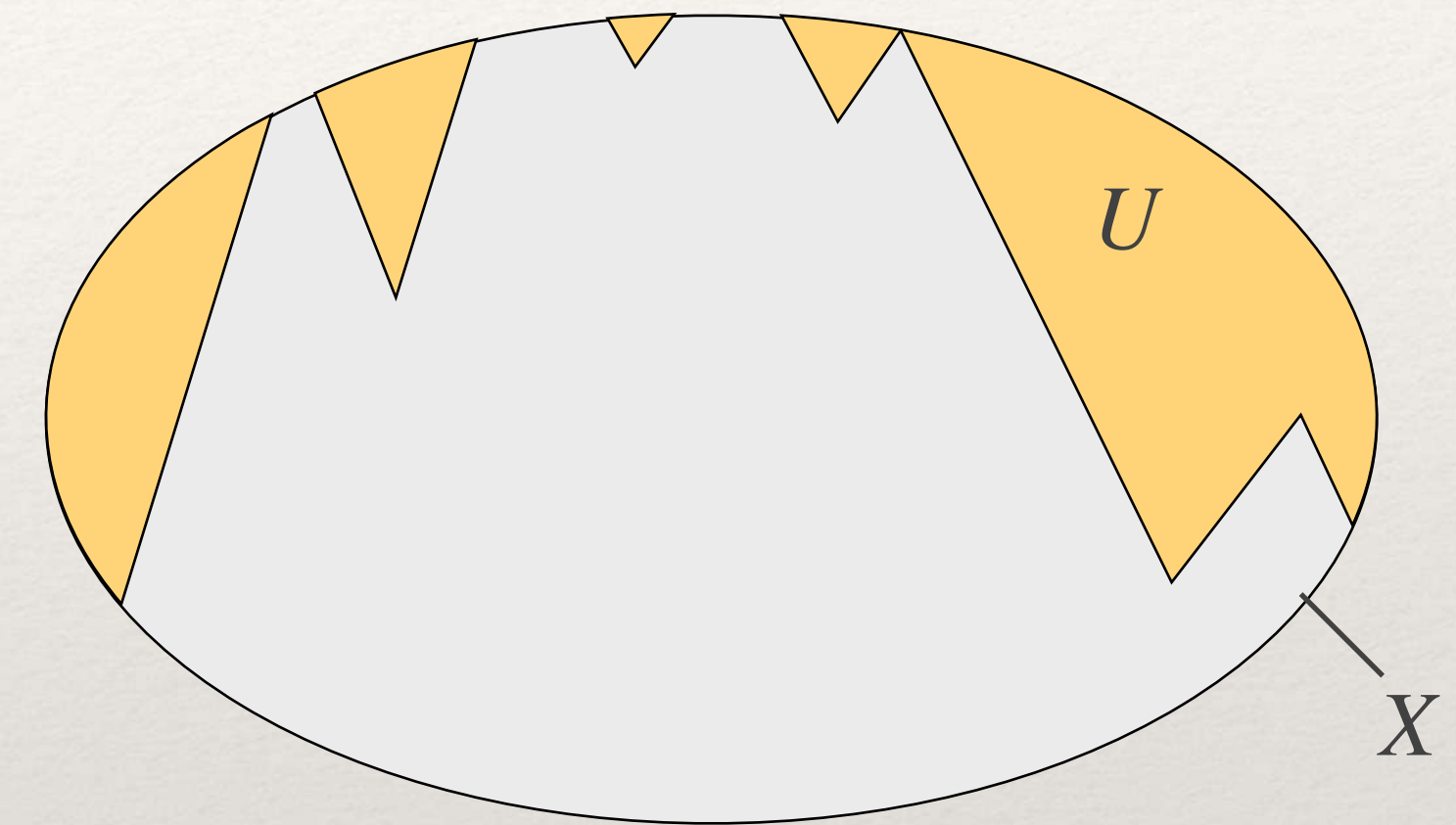
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- ❖ You still have to prove effectivity.
For that, you need to find a representation for open sets.
But open sets are **no longer** of the form $\uparrow \{x_1, \dots, x_n\}$

Representations, sobifications

Representing open sets: the trick

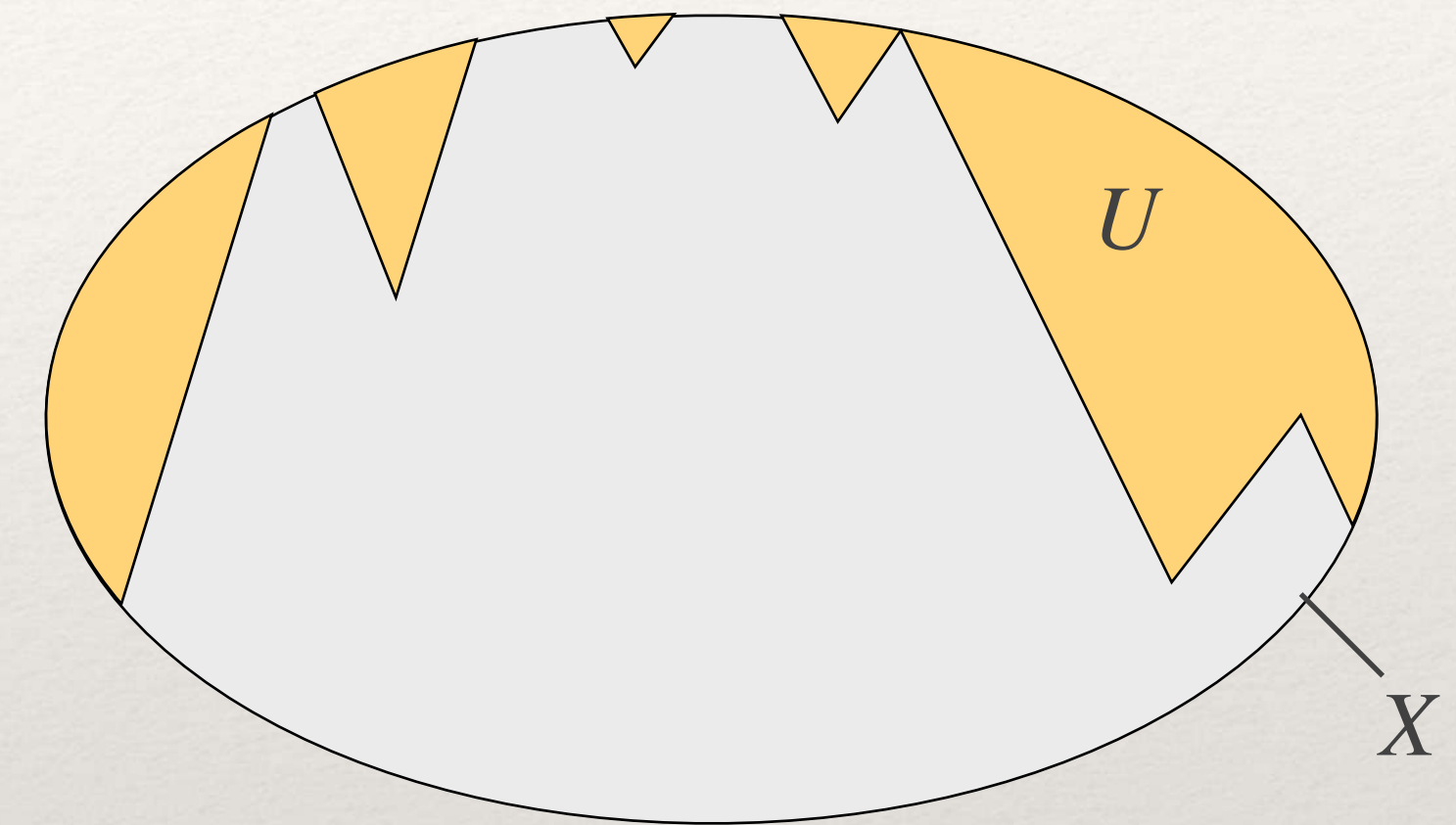
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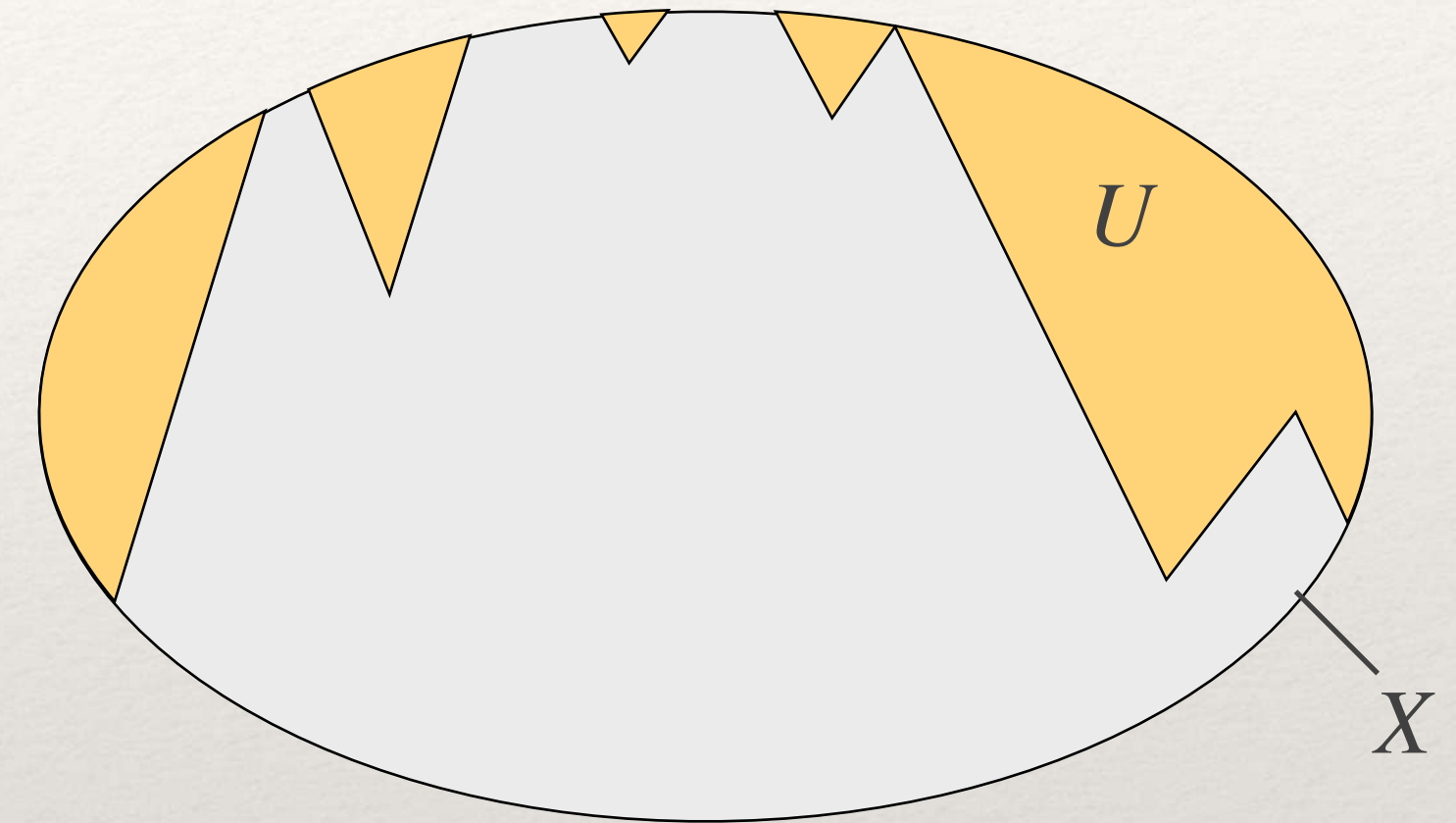
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Oops, I have not said what that was,
have I?



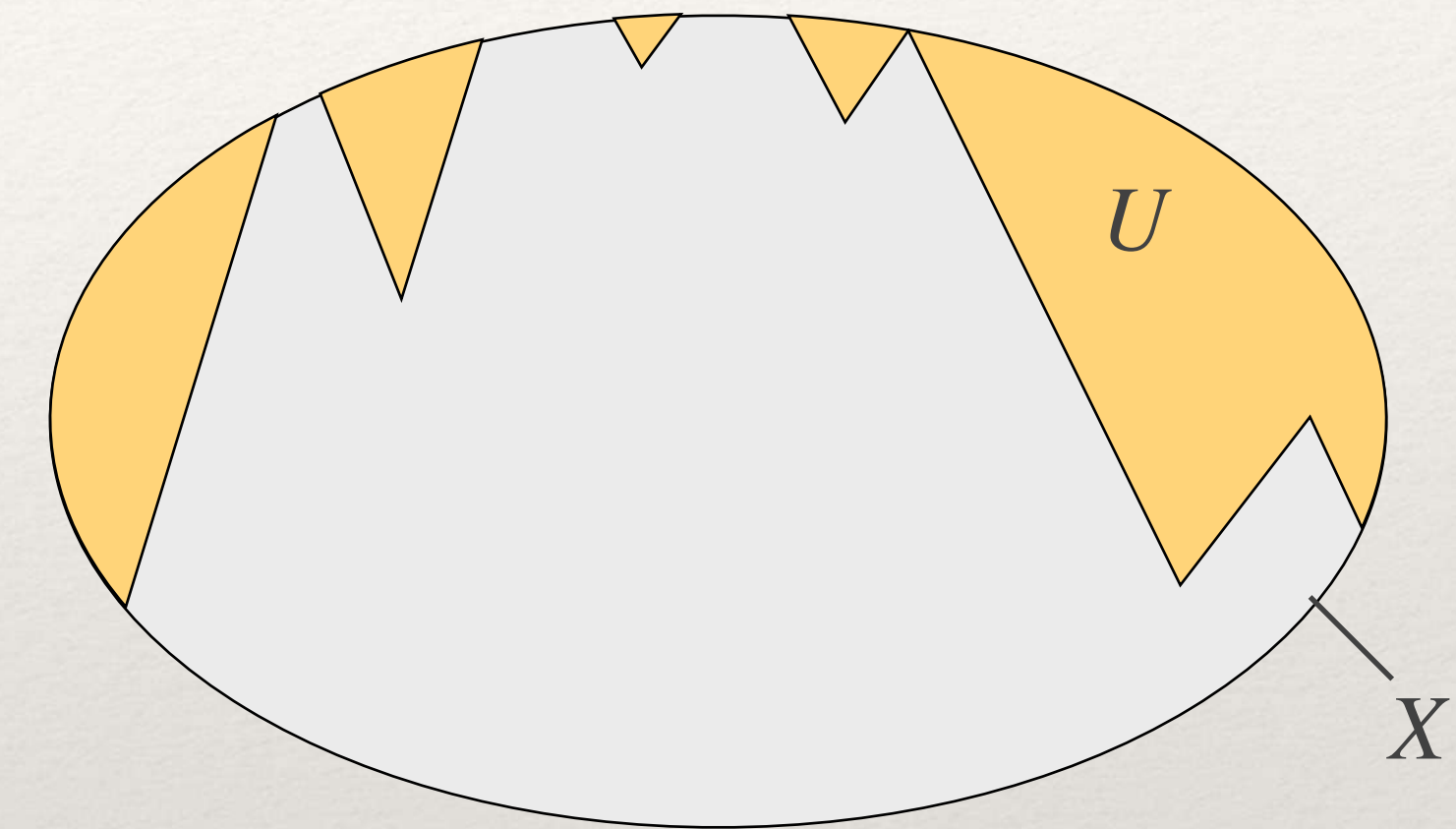
Sober spaces and sobrifications

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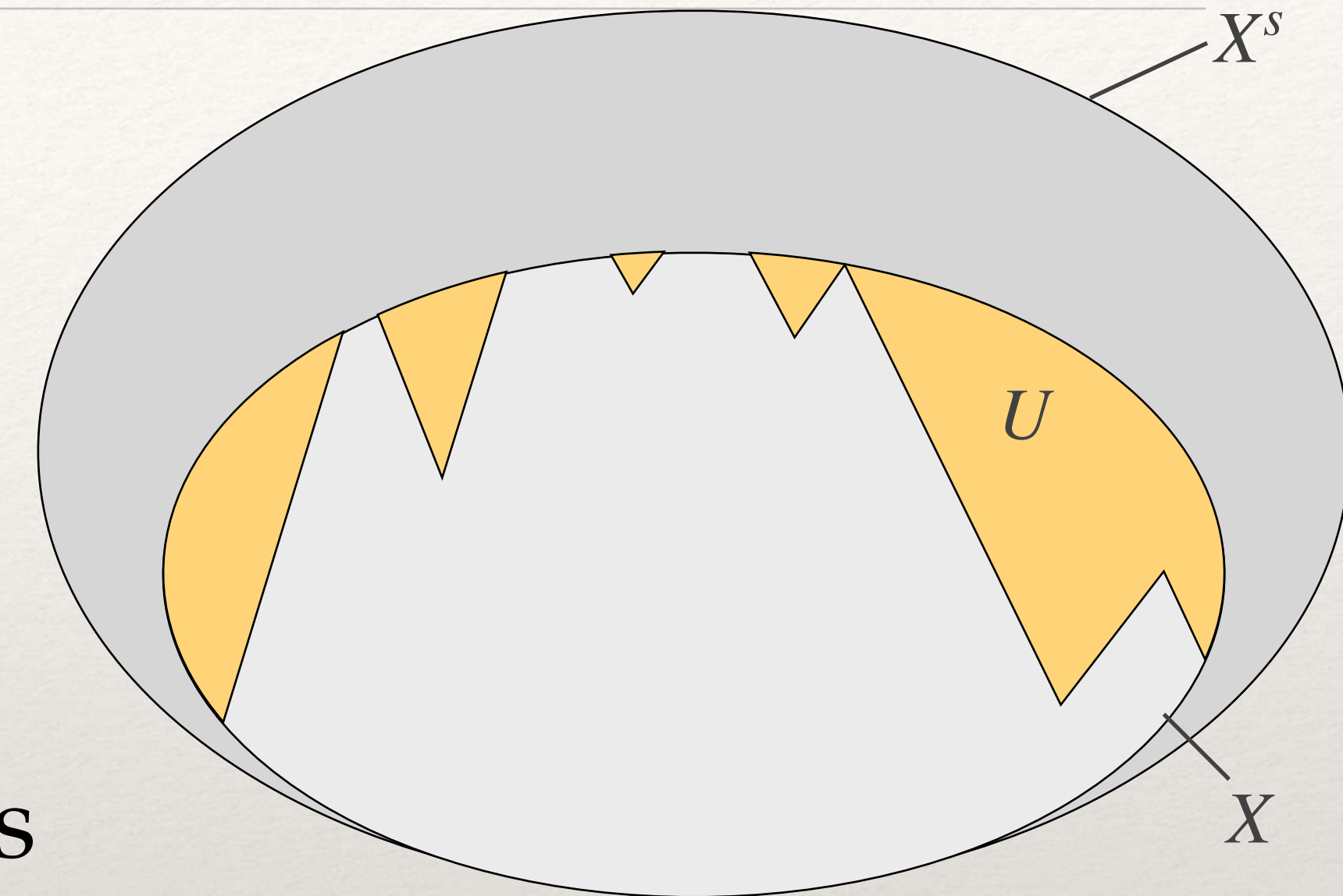
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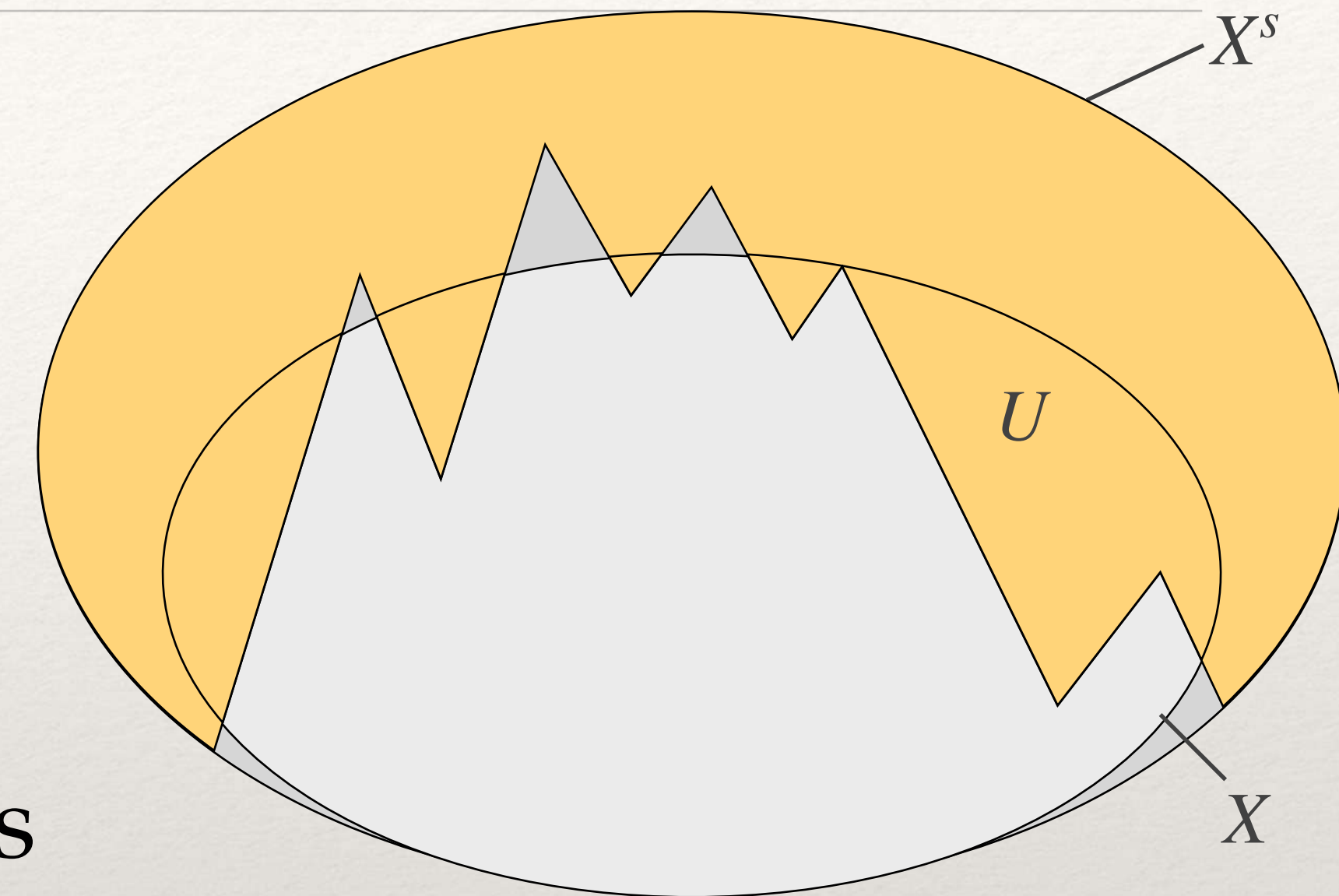
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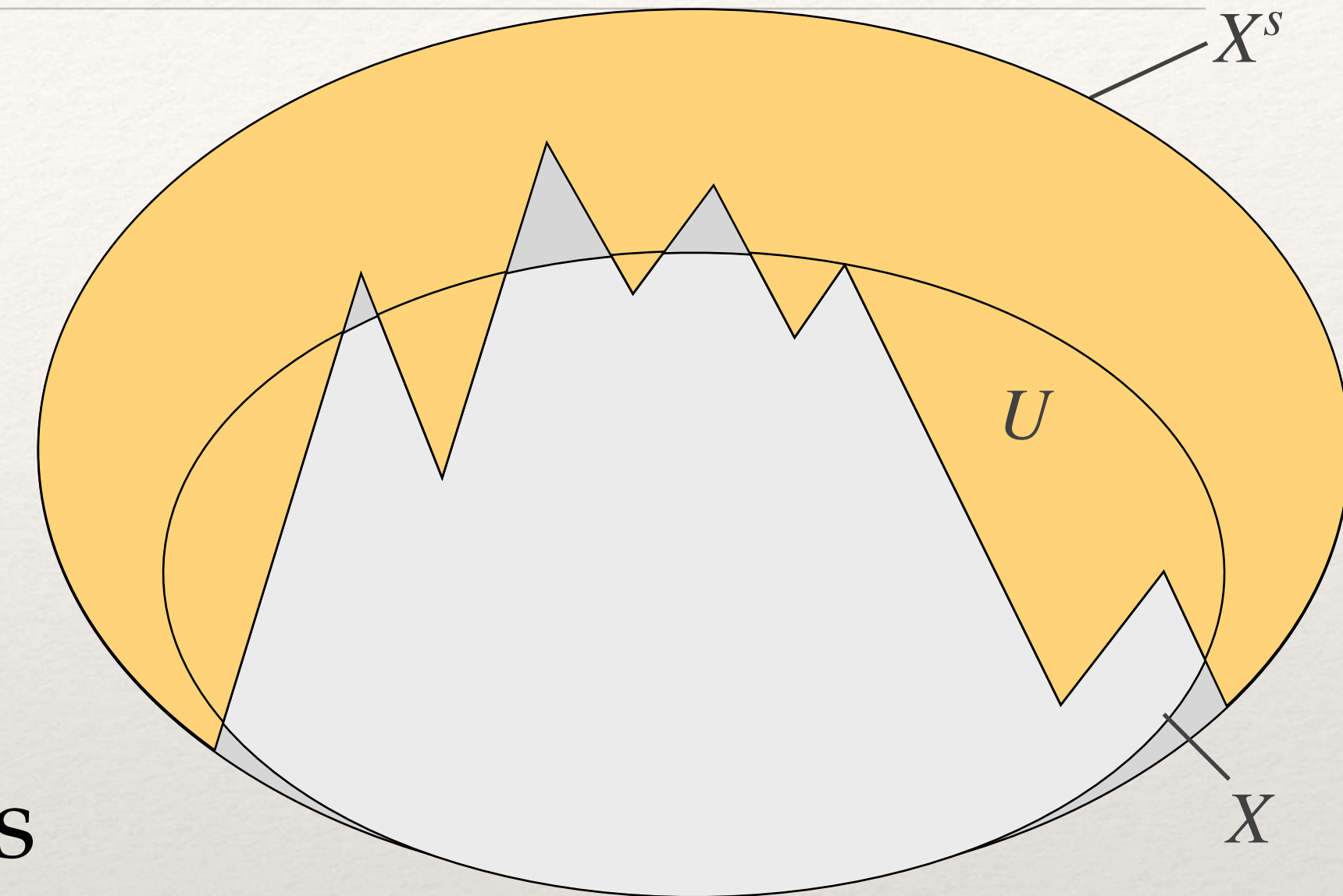
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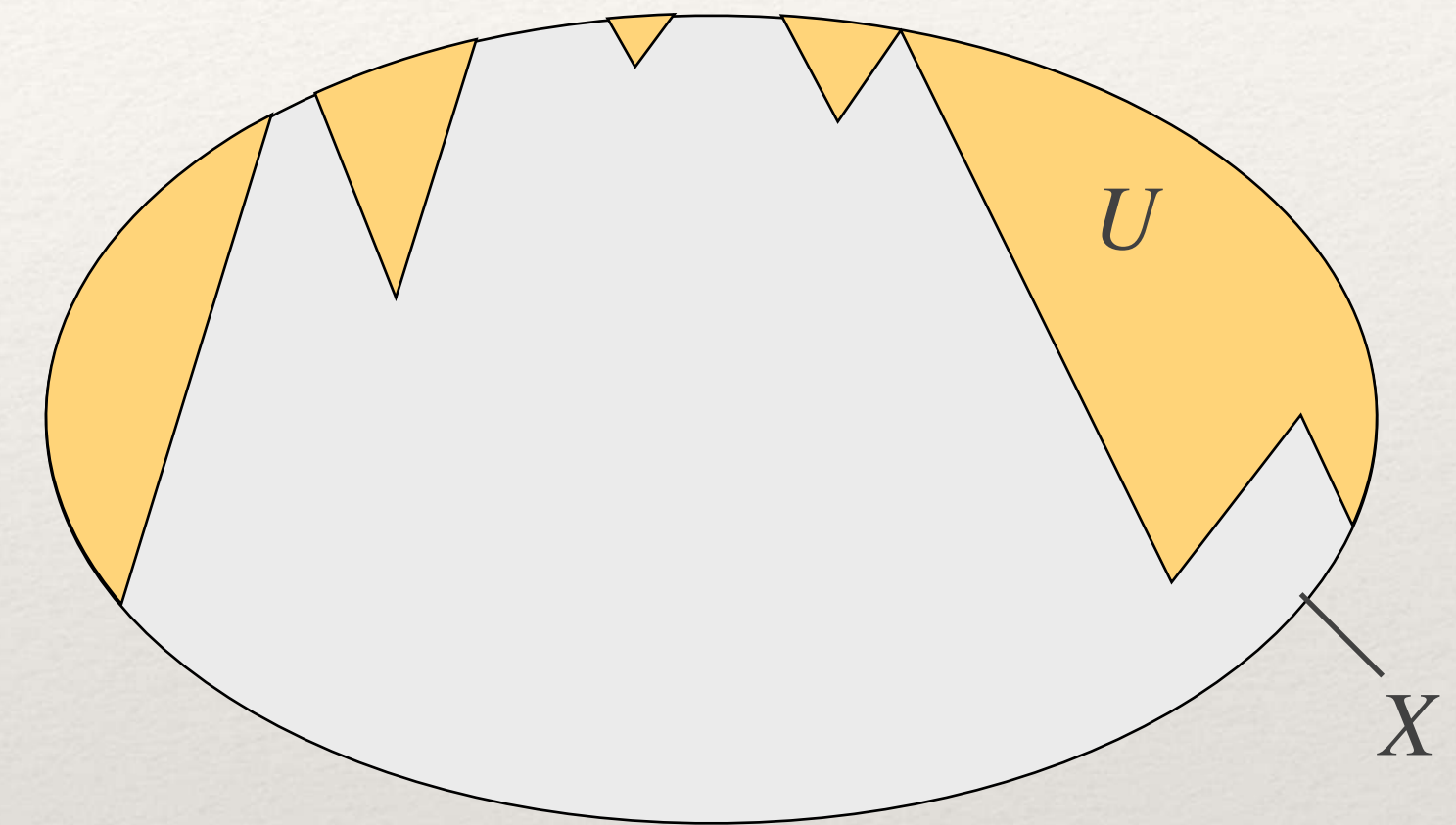


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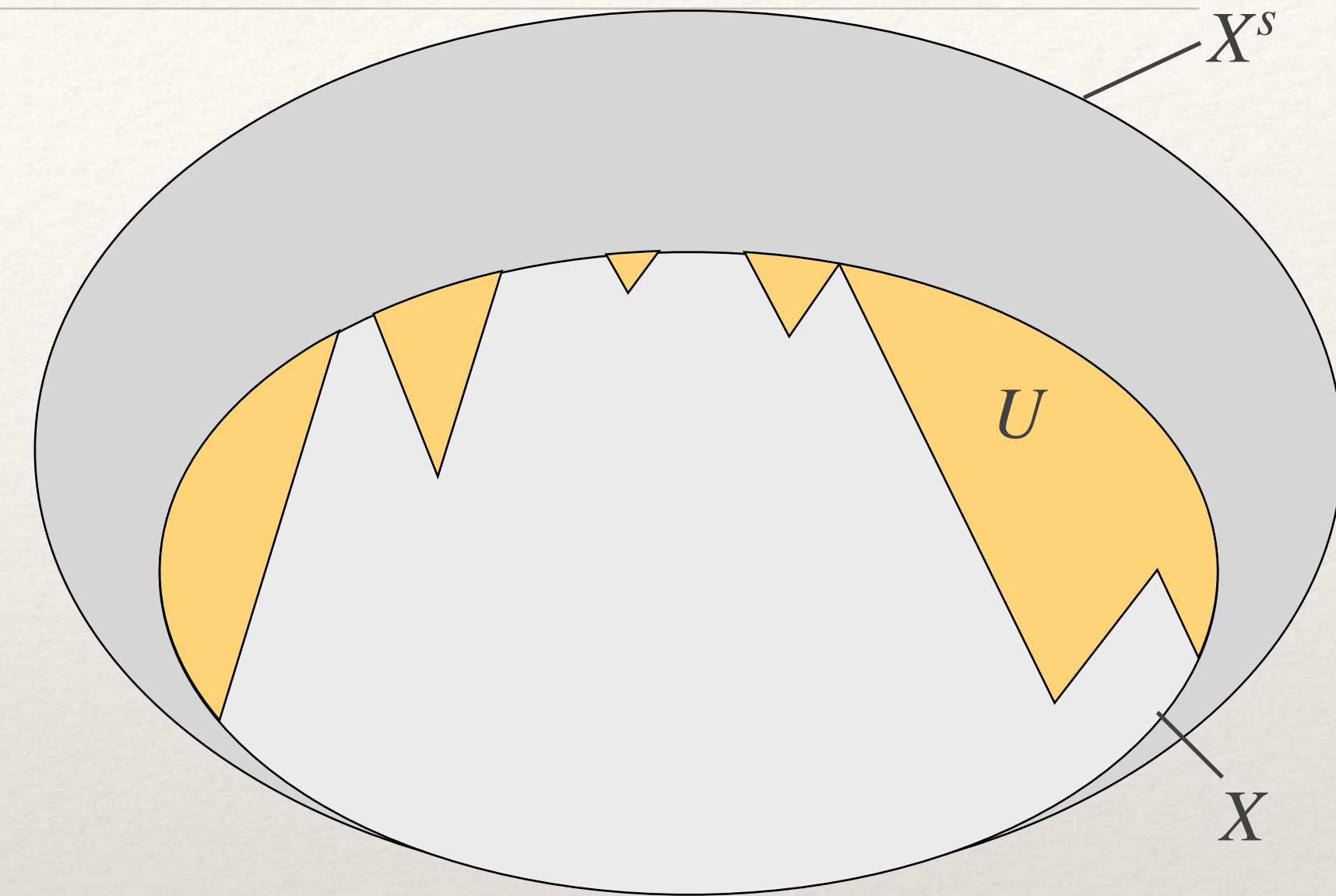
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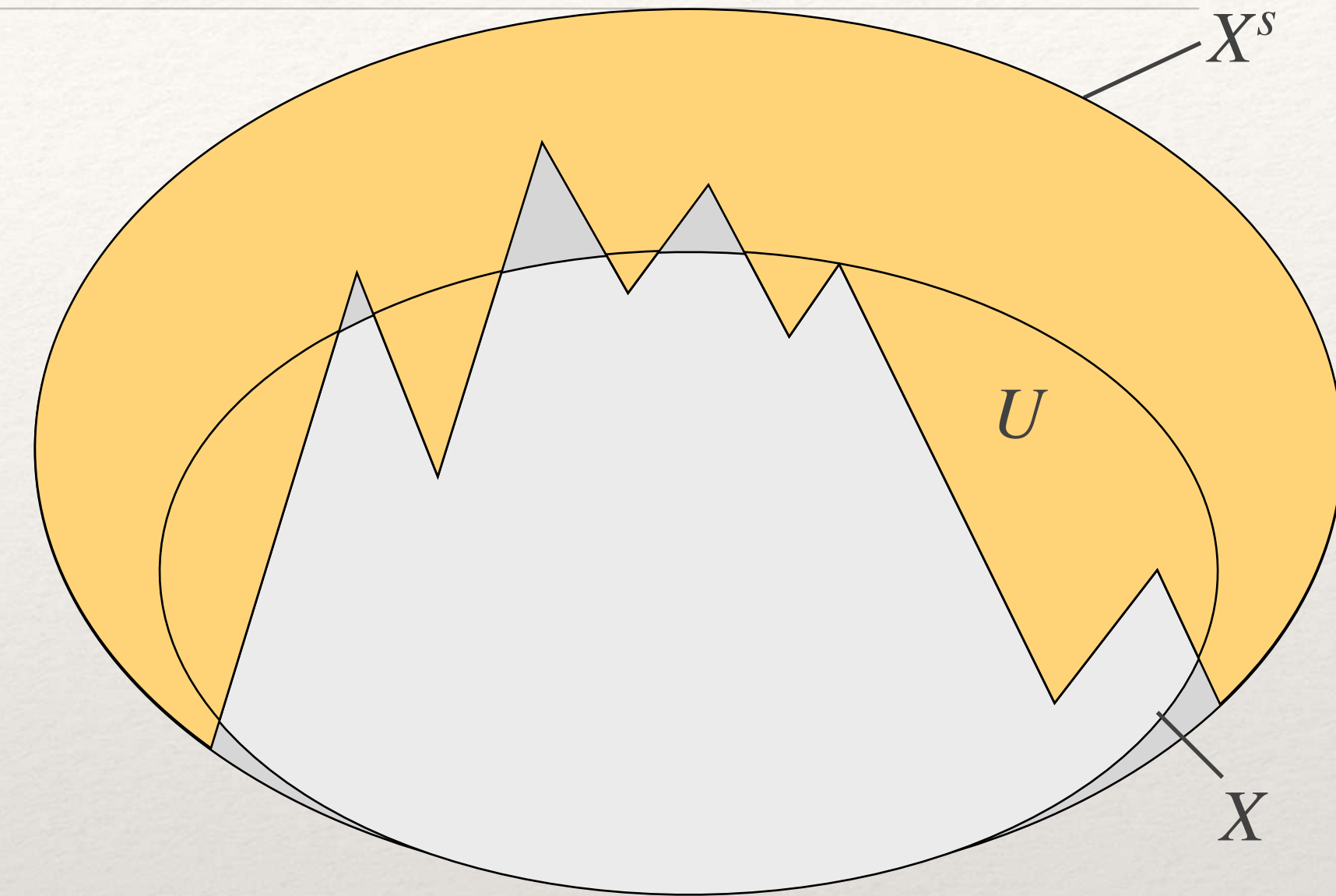
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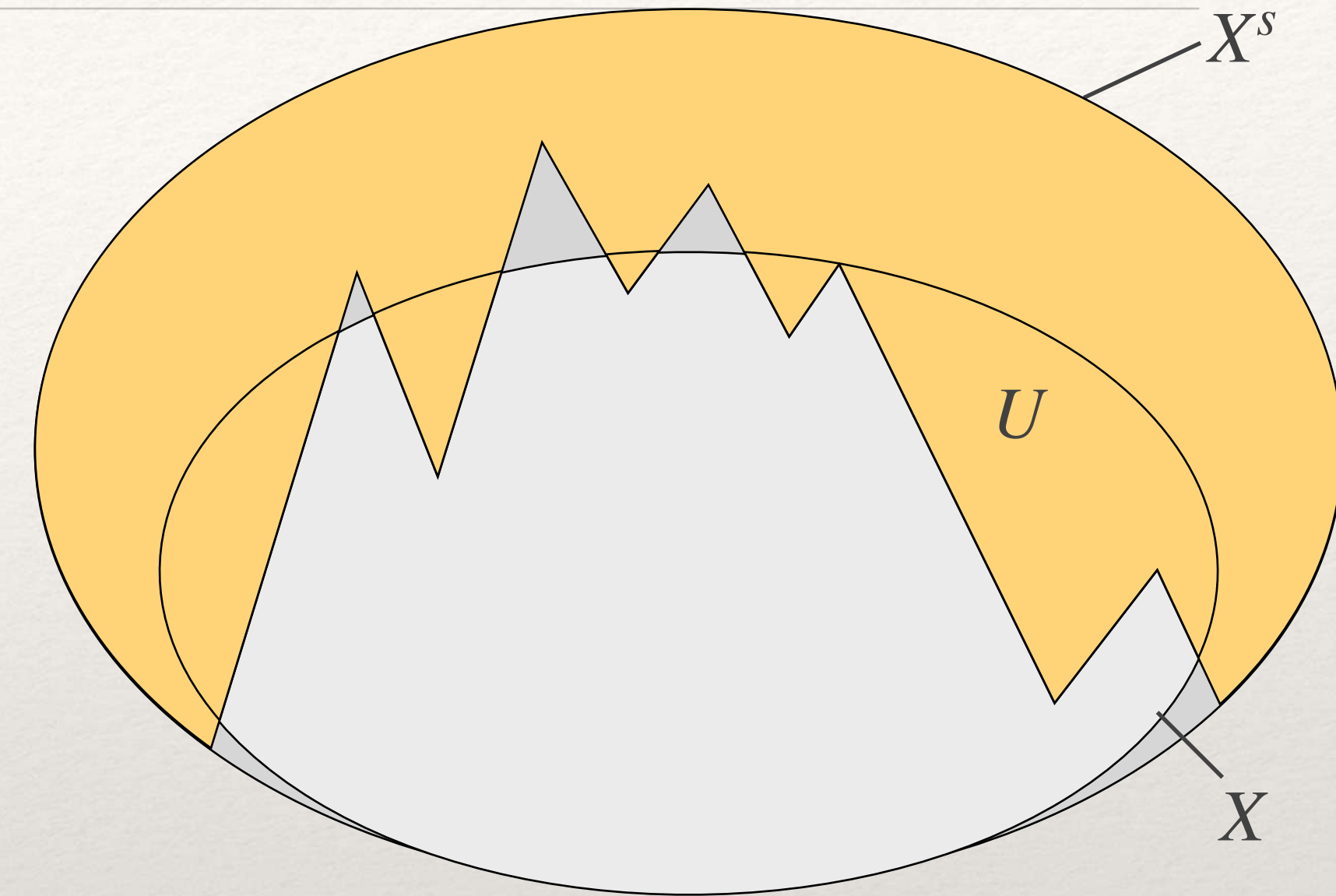
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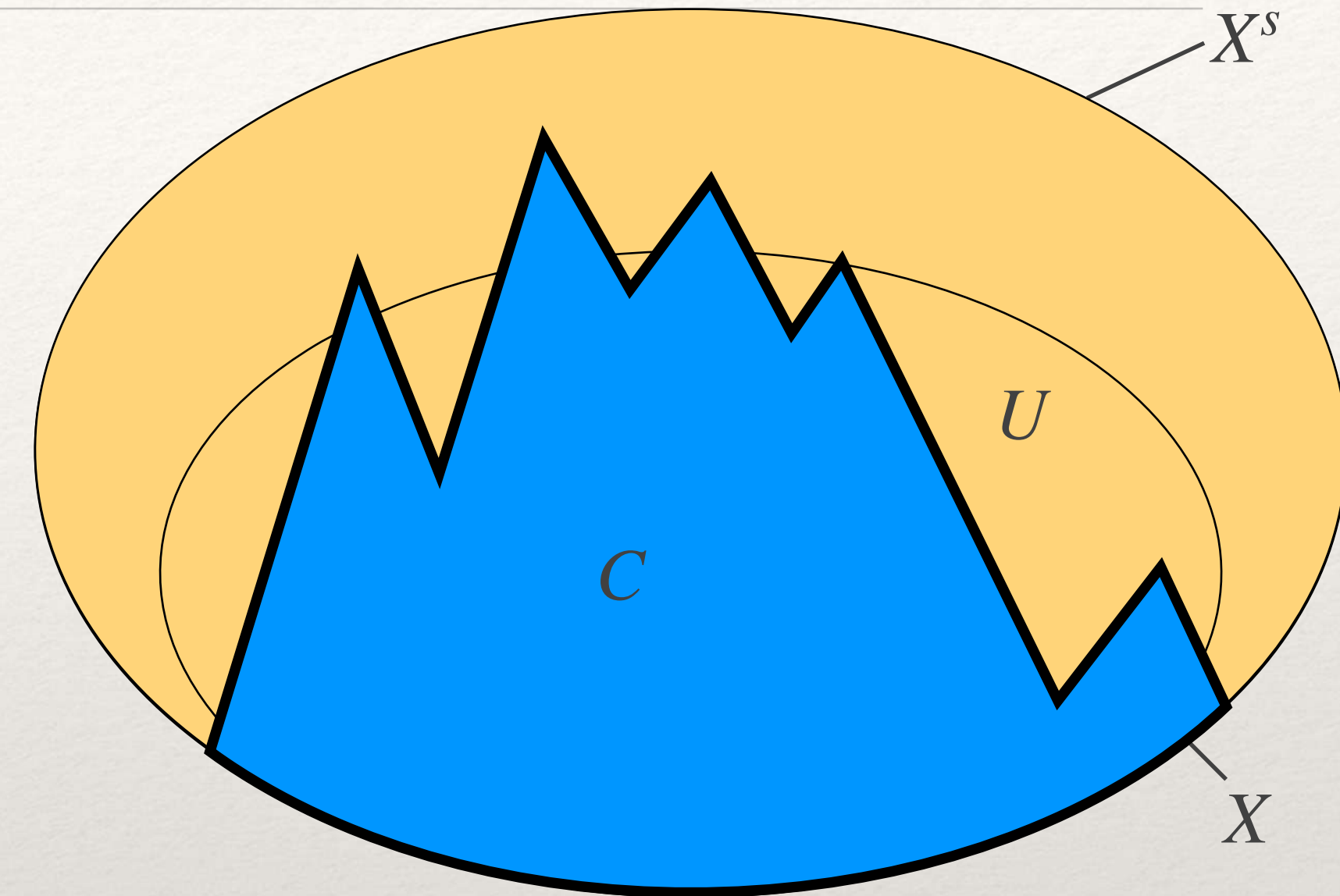
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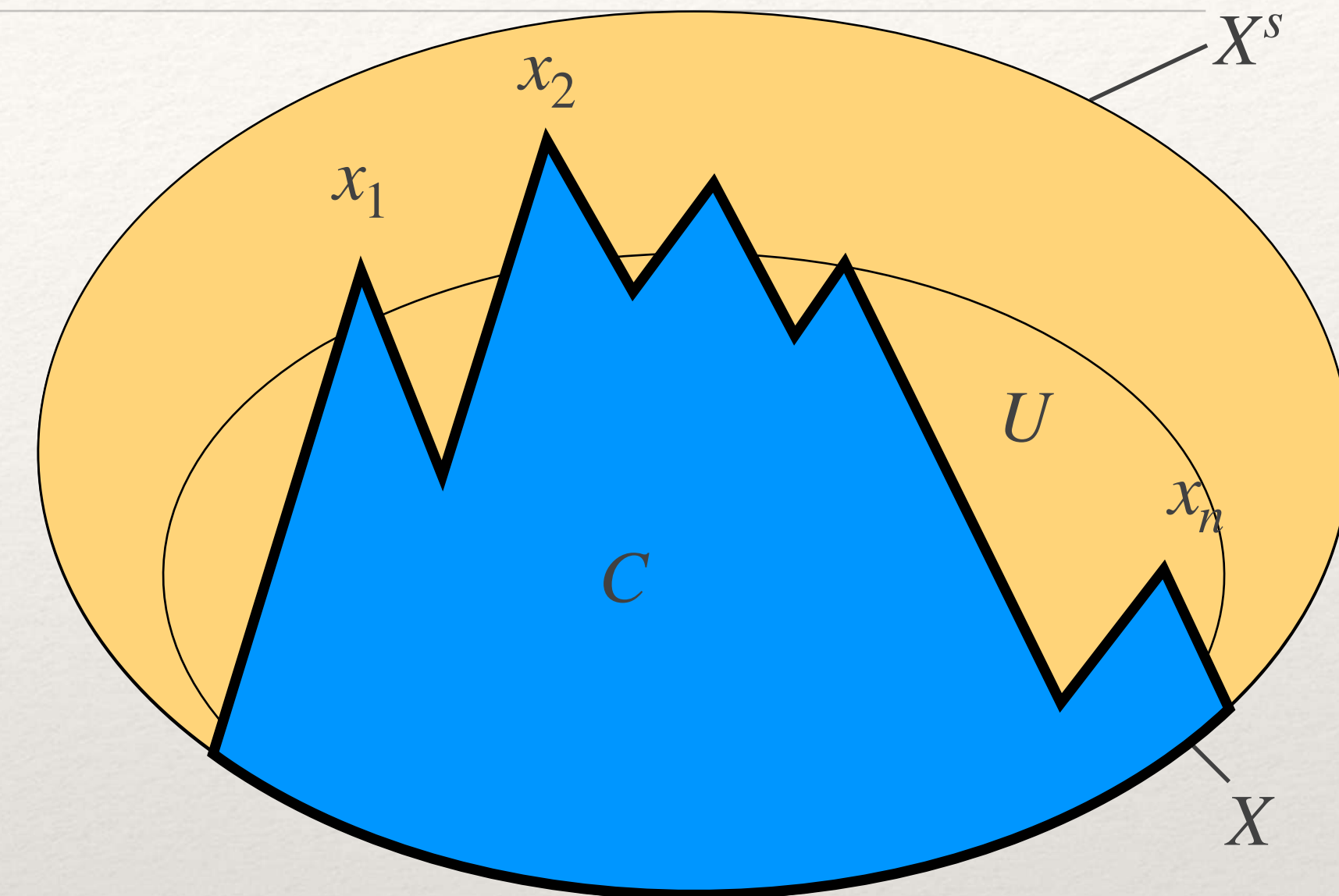


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Reminder

- ❖ **Proposition.** If:
 - X is well-founded
 - (Property T) X is finitary
 - (Property W) For all $x, y \in X$, $\downarrow x \cap \downarrow y$ is finitarythen X is **Noetherian** in the upper topology and the **closed** sets are exactly the **finitary** subsets.

This turns out to be the general form of all **sober** Noetherian spaces.

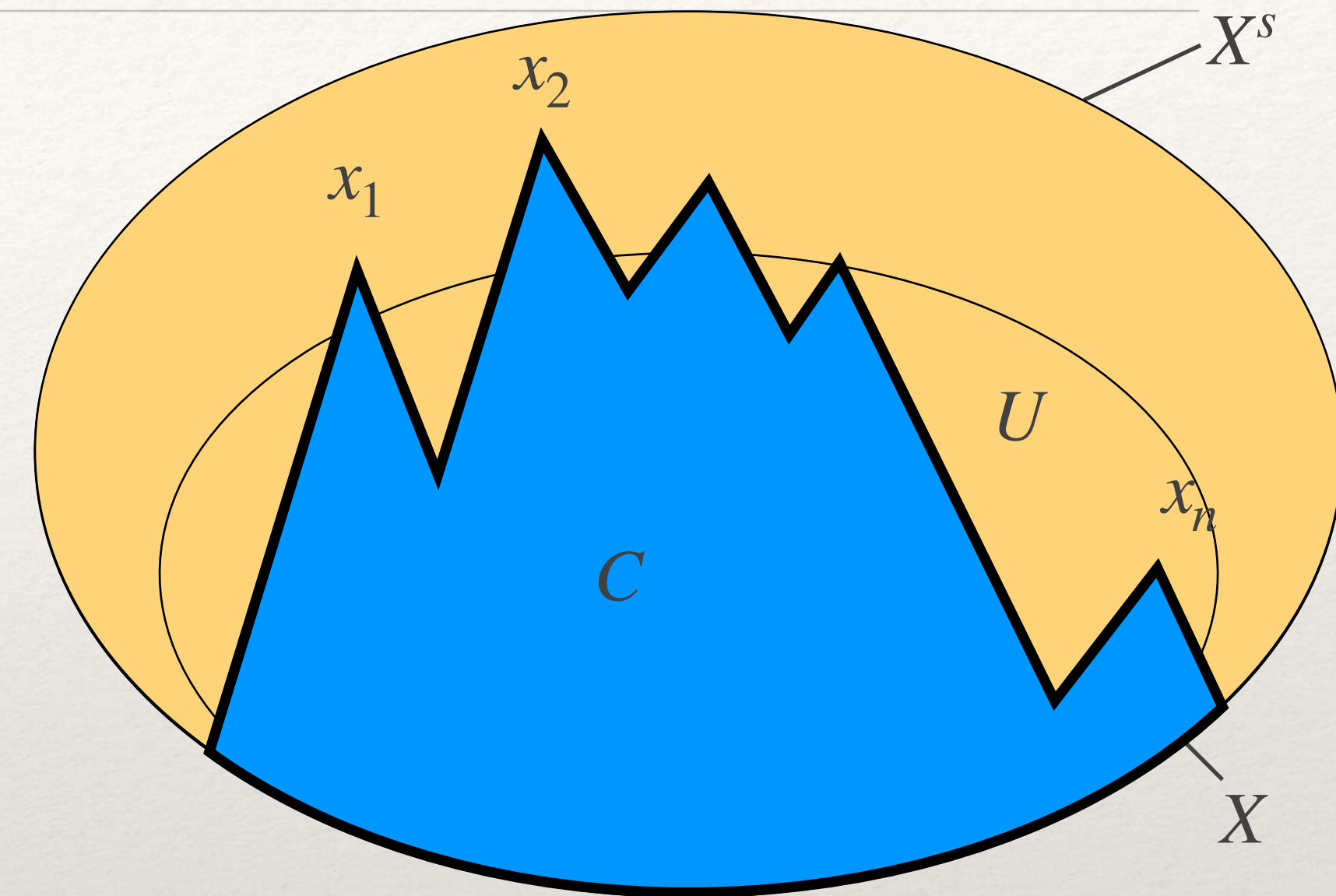
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In a sober Noetherian space, every closed set C is a **finitary** subset $\downarrow \{x_1, \dots, x_n\}$.

- ❖ Hence we can represent U by
(the complement of the downward closure in X^s)
of **finitely many** points... in X^s



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Representing points in sobrifications

- ❖ For a finite set Σ , with the discrete topology,

$$\Sigma^s = \Sigma$$

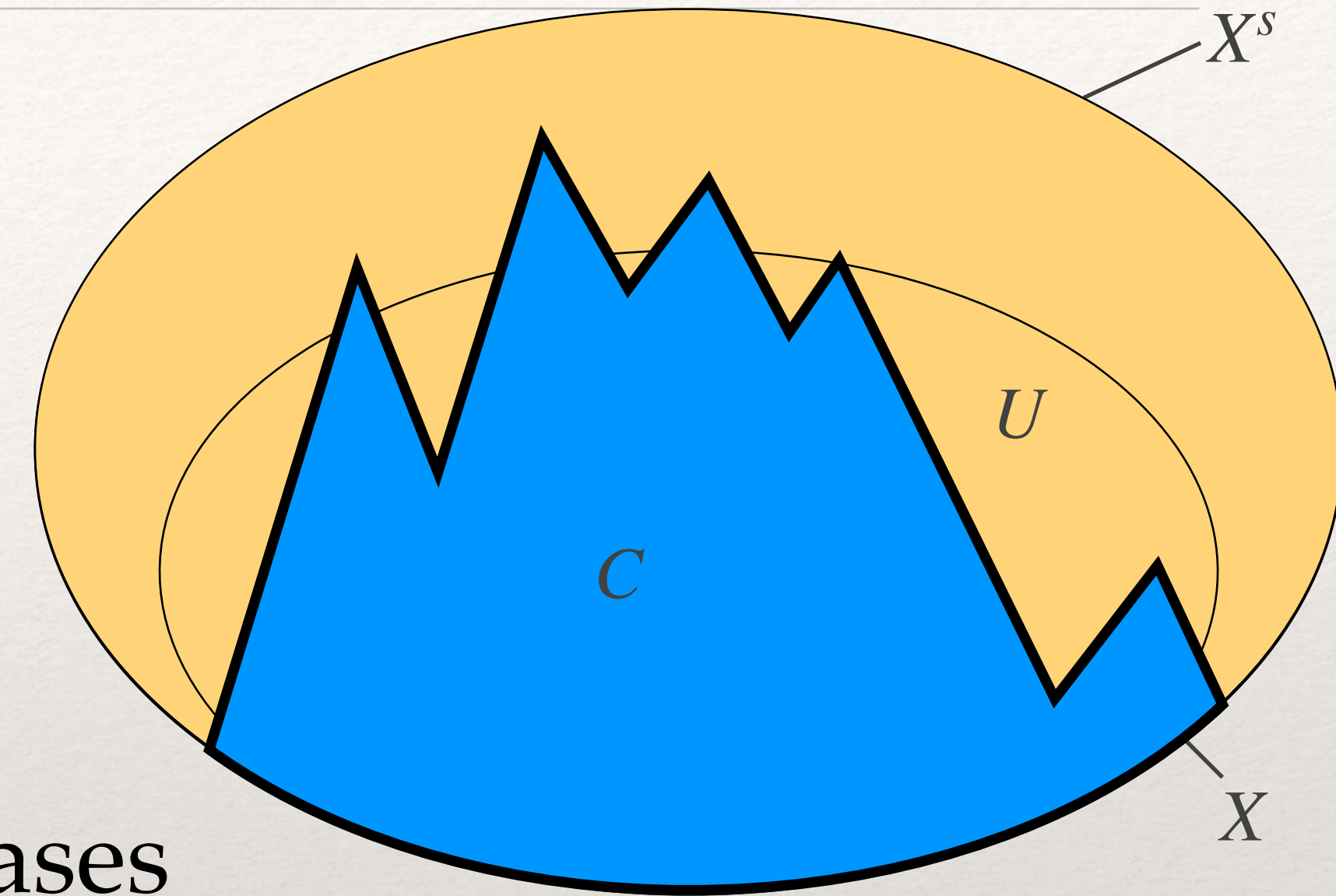
- ❖ Products: $(X \times Y)^s = X^s \times Y^s$

- ❖ $\text{Spec}(\mathbb{Z}[X_1, \dots, X_n])$: already sober,
points = prime ideals, represented as Gröbner bases

- ❖ $(X^*)^s$ consists of **word products**

$$P ::= \epsilon \mid C^?P \mid F^*P$$

$$\text{with } C \in X^s, F = C_1 \cup \dots \cup C_n \ (C_i \in X^s)$$



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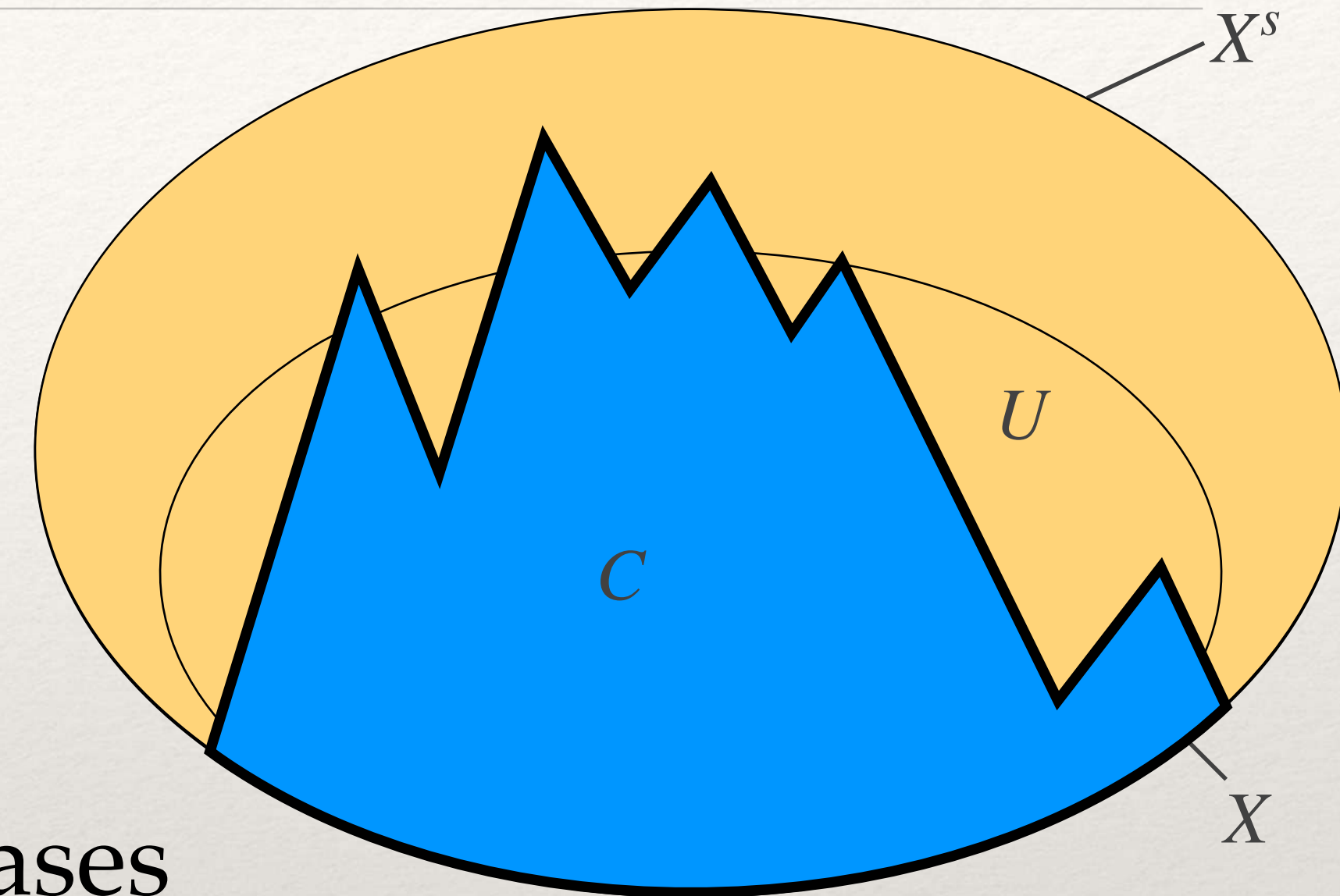
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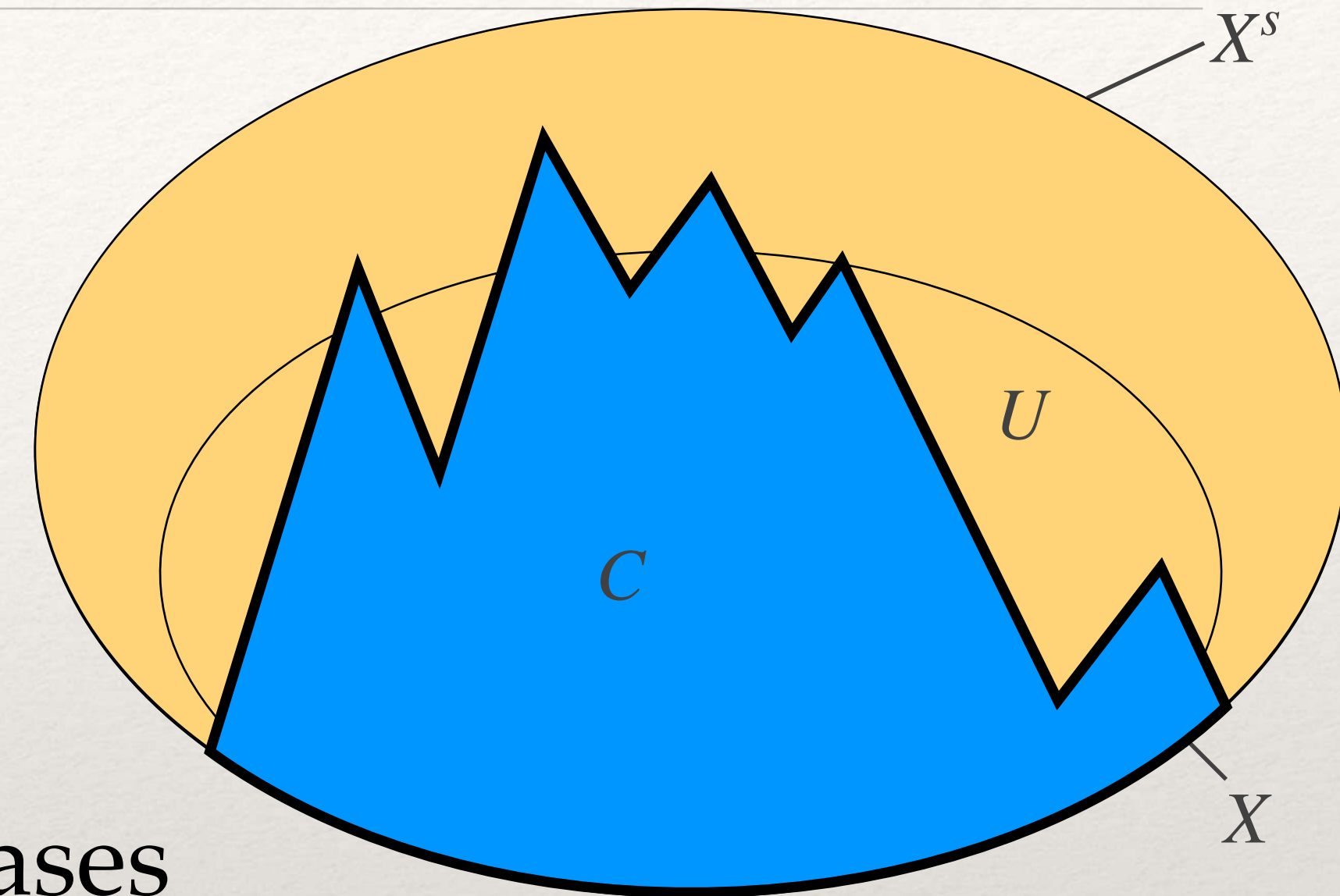
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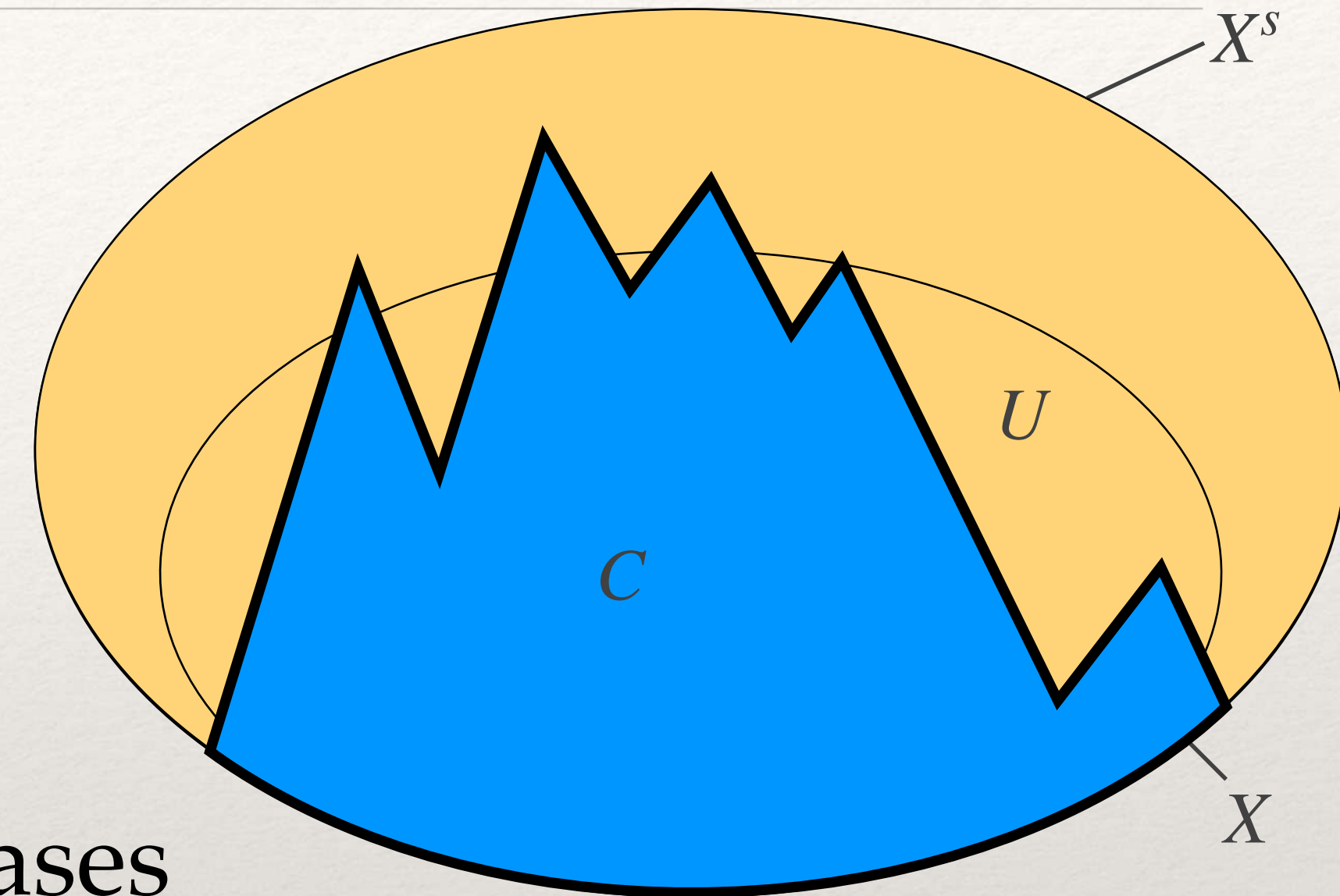
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e.g.,
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- All those are **representable** on a computer (Finkel, JGL 2009, 2021)

Structures of Noetherian spaces

-
-
- ❖ **Maximal order types** of well-partial-orderings
 - ❖ **Statures** of Noetherian spaces as generalization of maximal order types
 - ❖ ... we are not really changing the subject,
and we will use the **representations** of points in X^S again

Maximal order types

- ❖ A well-partial-ordering is a well-quasi-ordering that is antisymmetric
- ❖ **Theorem (Wolk 1967).** A wpo is a partial ordering whose linear extensions are all **well-founded**

Note: every linear well-founded ordering is isomorphic to a unique ordinal,
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❖ **Theorem (de Jongh, Parikh 1977).** Among those, one has **maximal** order type.

❖ Any meaningful equivalent of that notion for Noetherian spaces?

But first, why should we bother about maximal order types anyway?

Why bother about maximal order types?

- ❖ First studied by de Jongh and Parikh (1977)
then Schmidt (1979)
- ❖ Many applications in proof theory (reverse mathematics):
Simpson (1985), after Friedman
van den Meeren, Rathjen, Weiermann (2014, 2015)
etc.
- ❖ Ordinal complexity of the size-change principle for proving
the termination of programs and rewrite systems
Blass and Gurevich (2008)
- ❖ and...

Why bother about maximal order types?

- ❖ Figueira, Figueira, Schmitz and Schnoebelen (2011),
Schmitz and Schnoebelen (2011) (and others)

obtain **complexity upper bounds**
for algorithms whose termination
is based upon wqo arguments
(e.g., coverability)

length function
(complexity upper bound)

Theorem 5.3 (Main Theorem).

Let g be a smooth control function

eventually bounded by a function in \mathcal{F}_γ

and let A be an exponential nwqo

with **maximal order type** $< \omega^{\beta+1}$.

Then $L_{A,g}$ is bounded by a function in:

- ❖ \mathcal{F}_β if $\gamma < \omega$ (e.g., if g is primitive recursive) and $\beta \geq \omega$
- ❖ $\mathcal{F}_{\gamma+\beta}$ if $\gamma \geq 2$ and $\beta < \omega$.

class of functions
elementary recursive in F_β
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- ❖ E.g., coverability
in lossy channel systems
is F_{ω^ω} -complete.
(way larger than Ackermann)

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Going topological

- ❖ Let us return to the question of finding a **Noetherian analogue** of maximal order types

A wrong idea: minimal T_0 topologies

- ❖ Partial ordering $\sim T_0$ topology
Extension \sim coarser T_0 topology
Linear extension = maximal extension \sim minimal T_0 topology
- ❖ Studied by Larson (1969).
A minimal T_0 topology is necessarily the **upper** topology of a **linear** ordering.
- ❖ Unfortunately, minimal T_0 topologies do not exist in general:
Fact. \mathbb{R}_{cof} is Noetherian, but has no coarser minimal T_0 topology.
(Its uncountably many proper closed subsets would all have to be finite, and linearly ordered.)

Statures of wpos

❖ **Theorem** (Kříž 1997, Blass and Gurevich 2008).

Maximal order type of a wpo (X, \leq)

= **ordinal rank** $||X||$ of the top element X

in the poset $(\mathcal{D}X, \subseteq)$ of downwards-closed subsets of X

The **stature** of X

❖ Ordinal rank inductively defined by:

$$||F|| = \sup\{ ||F'|| + 1 \mid F' \in \mathcal{D}X, F' \subsetneq F \}$$

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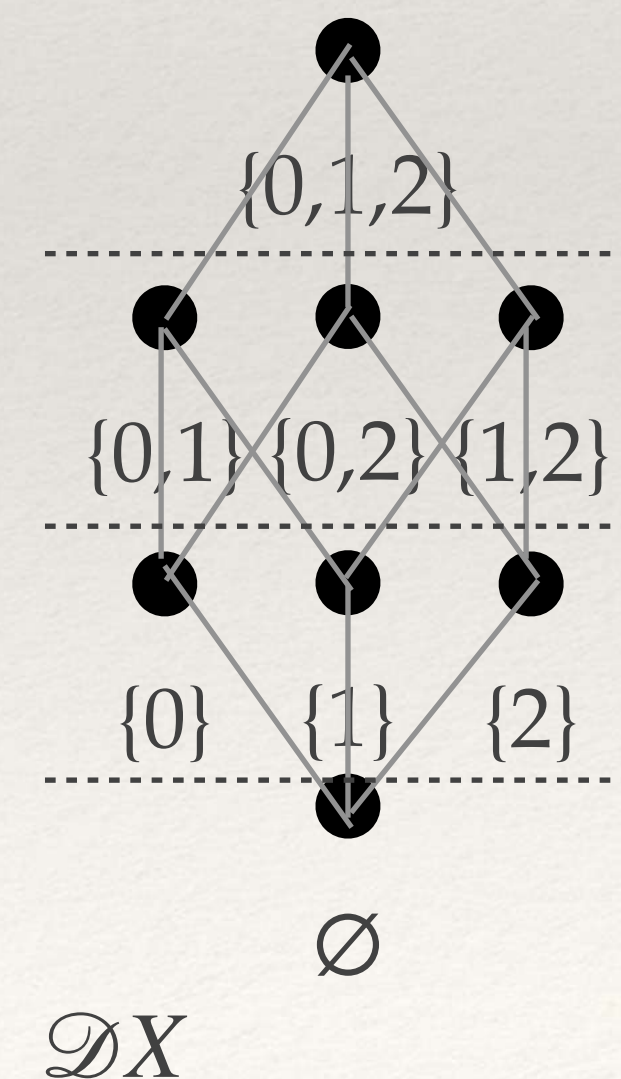
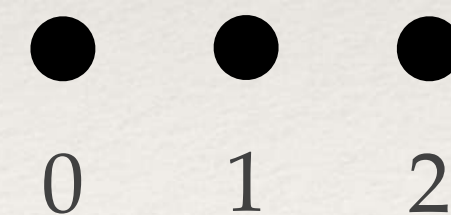
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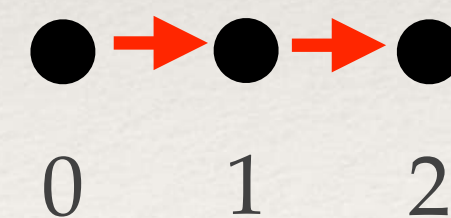
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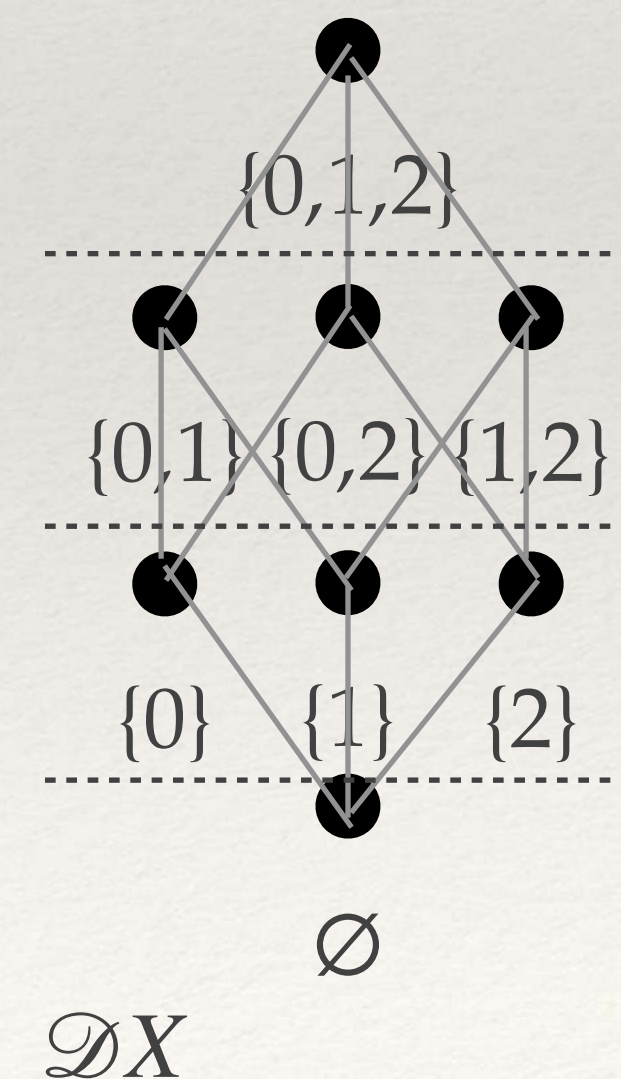
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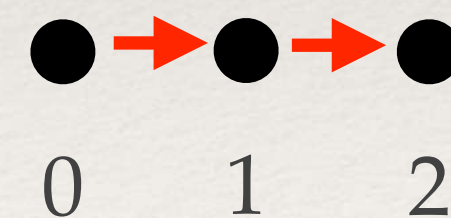
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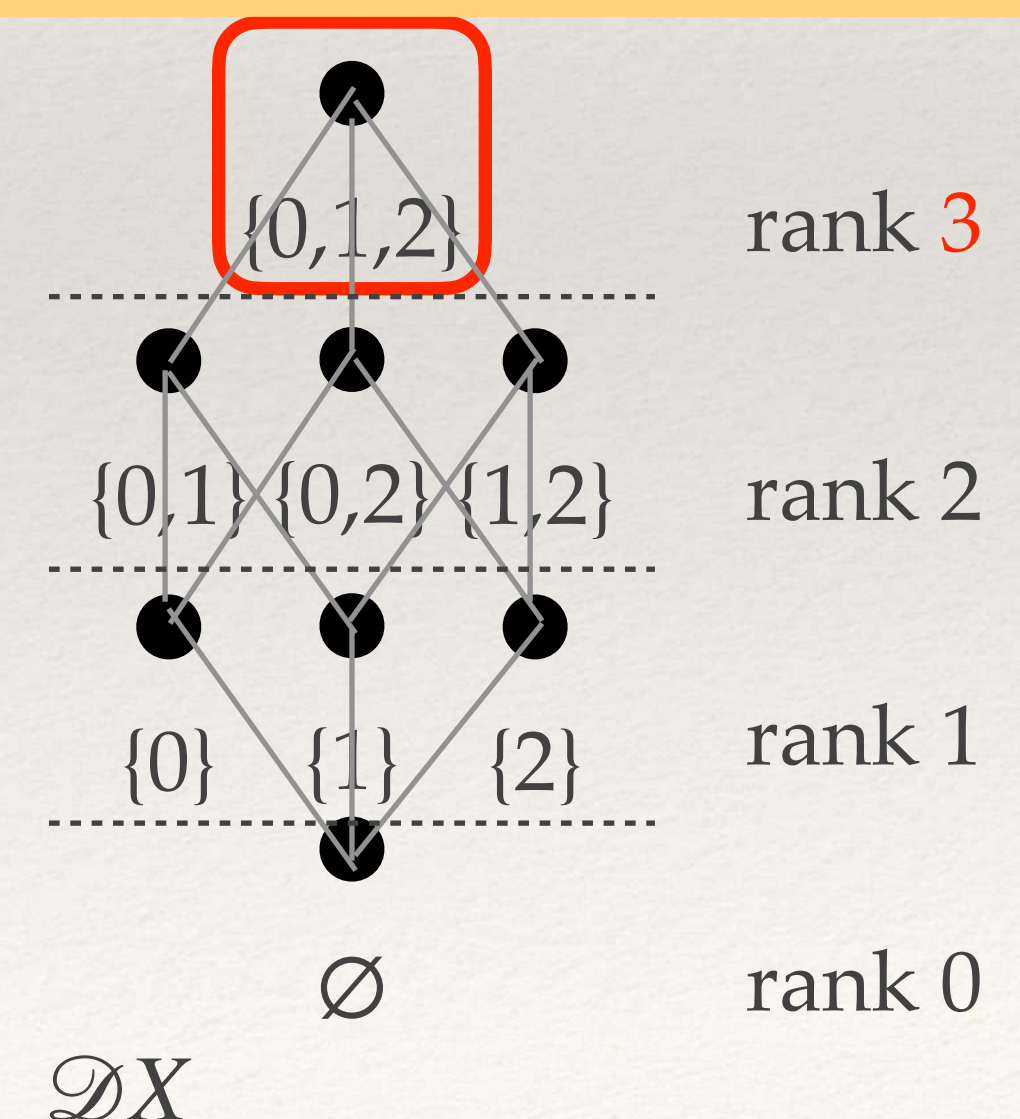
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Statures of Noetherian spaces

❖ **Definition.** The **stature** of a Noetherian space X is the **ordinal rank** $||X||$ of the top element X in the poset $(\mathcal{H}X, \subseteq)$ of **closed** subsets of X

$$❖ \quad ||F|| = \sup\{ ||F'|| + 1 \mid F' \in \mathcal{H}X, F' \subsetneq F \}$$

❖ Matches previous definition:
for a wqo in its Alexandroff topology,
closed = downwards-closed

$$\mathcal{H}X = \mathcal{D}X$$

X is Noetherian iff:

- (6) Every antitonic chain $F_1 \supseteq F_2 \supseteq \dots \supseteq F_n \supseteq \dots$ of closed subsets is **stationary**
- (7) $\mathcal{H}X$ is well-founded.

Some statures of Noetherian spaces

- ❖ We have already obtained **statures** of quite a few Noetherian constructions (JGL, Laboureix 2022)
- ❖ Let me focus on X^*

X	$\text{sob } X$		$\ X\ $	
Finite T_0	$\leq \text{card } X$		$\text{card } X$	Lem. 6.1
Ordinal α (Alex.)	$\alpha / \alpha + 1$	Lem. 6.2	α	Lem. 6.2
Ordinal α (Scott)	$\alpha / \alpha + 1$	Lem. 6.2	$\alpha / \alpha - 1$	Lem. 6.2
Cofinite topology	$1 / 2$	Thm. 7.1	$\min(\text{card } X, \omega)$	Thm. 7.2
$X + Y$	$\max(\text{sob } X, \text{sob } Y)$	Prop. 8.4	$\ X\ \oplus \ Y\ $	Prop. 8.2
$X +_{\text{lex}} Y$	$\text{sob } X + \text{sob } Y$	Prop. 9.4	$\ X\ + \ Y\ $	Prop. 9.2
X_{\perp}	$1 + \text{sob } X$	Prop. 9.6	$1 + \ X\ $	Prop. 9.6
$X \times Y$	$(\text{sob } X \oplus \text{sob } Y) - 1$	Prop. 10.1	$\ X\ \otimes \ Y\ $	Thm. 10.9
$\mathcal{H}_{\text{ov}} X, \mathcal{H}_{\text{fin}} X, \mathbb{P}X, \mathbb{P}_{\text{fin}} X$	$\ X\ + 1$	Thm. 11.1	$\geq 1 + \ X\ ,$ $\leq \omega^{\ X\ }$	Prop. 11.2
X^*	$\omega^{\ X\ ^{\circ}} + 1$ $(\alpha^{\circ} \stackrel{\text{def}}{=} \alpha + 1 \text{ if } \alpha = \epsilon + n, \epsilon \text{ critical, } n \in \mathbb{N},$ $\alpha \text{ otherwise})$	Thm. 12.13	$\omega^{\omega^{\ X\ ^{\circ}}}$ $(\alpha'^{\circ} \stackrel{\text{def}}{=} \alpha - 1 \text{ if } \alpha \text{ finite,}$ $\alpha^{\circ} \text{ otherwise})$	Thm. 12.22
$\bigtriangleright_{n=1}^{+\infty} X_n$	$\bigoplus_{n=1}^{+\infty} \text{rsob } X_n + 1 /$ $\bigoplus_{n=1}^k \text{rsob } X_n + \omega + 1$	Thm. 13.4	$\bigotimes_{n=1}^{+\infty} \ X_n\ /$ $\bigotimes_{m=1}^k \ X_m\ \times \omega$	Thm. 13.8
X^{\triangleright}	$\omega^{\alpha_1+1} + 1$ where $\text{sob } X - 1 =_{\text{CNF}} \omega^{\alpha_1} + \dots$	Cor. 13.7	$\omega^{\omega^{\beta_1+1}} / \omega$ where $\ X\ =_{\text{CNF}} \omega^{\alpha_1} + \dots,$ $\alpha_1 =_{\text{CNF}} \omega^{\beta_1} + \dots$	Cor. 13.9
X^{\circledast}	$\geq (\omega \times \ X\) + 1,$ $\leq (\ X\ \otimes \omega) + 1$	Prop. 14.8, Prop. 14.9	$\omega^{\hat{\alpha}}$ $(\hat{\alpha} \stackrel{\text{def}}{=} \omega^{\alpha_1^{\circ}} + \dots + \omega^{\alpha_m^{\circ}}$ if $\alpha =_{\text{CNF}} \omega^{\alpha_1} + \dots + \omega^{\alpha_m}$)	Thm. 14.20

The stature of X^*

- ❖ **Theorem** (JGL, Laboureix 2022). If $X \neq \emptyset$ is Noetherian and $\alpha = ||X||$, then $||X^*|| = \omega^{\omega^{\alpha \pm 1}}$
(+1 if $\alpha = \epsilon_\beta + n$, -1 if α finite)
- ❖ Not very surprising: already known when X wqo (Schmidt 1979)

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❖ The proof is very different, and is **localic**.

Explicitly, we do not reason on points (words),

but on **closed sets** = finite unions of word products

$(X^*)^s$ consists of **word products**

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with $C \in X^s$, $F = C_1 \cup \dots \cup C_n$ ($C_i \in X^s$)

An excerpt from the proof of $||X^*|| \geq \omega^{\omega^{\alpha \pm 1}}$

- ❖ Let $F \subsetneq F \cup C$, $\mathbf{C}_0 = \emptyset$, $\mathbf{C}_{n+1} = (F^*C^?)^n F^*$, $\mathcal{A}_n = \{\mathbf{A} \in \mathcal{H}X \mid \mathbf{C}_n \subseteq \mathbf{A} \subsetneq \mathbf{C}_{n+1}\}$
- ❖ Map $(\mathbf{B} \subsetneq \mathbf{B}^+) \in \text{Step}(\mathcal{H}(F^*))$, $\mathbf{A} \in \mathcal{A}_n$ to $(F^*C^?)^{n+1} \mathbf{B} \cup \mathbf{A} C^? \mathbf{B}^+ \cup \mathbf{C}_{n+1}$

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- ❖ This is strictly monotonic : $\text{Step}(\mathcal{H}(F^*)) \times_{\text{lex}} \mathcal{A}_n \rightarrow \mathcal{A}_{n+1}$
- ❖ If $||F^*|| \geq \omega^{\omega^\beta}$ then $||\mathbf{C}_{n+1}|| \geq \omega^{\omega^\beta \times (n+1)}$,
so $||(F \cup C)^*|| \geq \omega^{\omega^{\beta+1}}$, by taking suprema over $n \in \mathbb{N}$
- ❖ This is the key step in a well-founded induction on $F \in \mathcal{H}X$
showing $||F^*|| \geq \omega^{\omega^{||F|| \pm 1}}$
- ❖ Finally, let $F = X$; by definition, $||X|| = \alpha$. \square

A finite union
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The stature of $\mathbb{Z}[X_1, \dots, X_n]$

- ❖ The ordinal height of the lattice of ideals of $\mathbb{Z}[X_1, \dots, X_n]$ is $\omega^n + 1$ (Aschenbrenner, Pong 2004)
- ❖ Hence $||\text{Spec}(\mathbb{Z}[X_1, \dots, X_n])|| = \omega^n$ (argument not quite written out yet, probably well-known)

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- ❖ Together with $||X \times Y|| = ||X|| \otimes ||Y||$ (JGL, Laboureix 2022) extending the same formula on wqos (de Jongh, Parikh 1977), we obtain the **stature** of the state space of **concurrent polynomial programs...**

The stature of the state space of concurrent polynomial programs

- ❖ m programs, each on n variables
- p queues, on $k \geq 1$ letters

- ❖ Stature of state space =

$$(\omega^n)^m \otimes (\omega^{\omega^{k-1}})^p$$

$$= \omega^{nm \oplus \omega^{k-1} \cdot p}$$

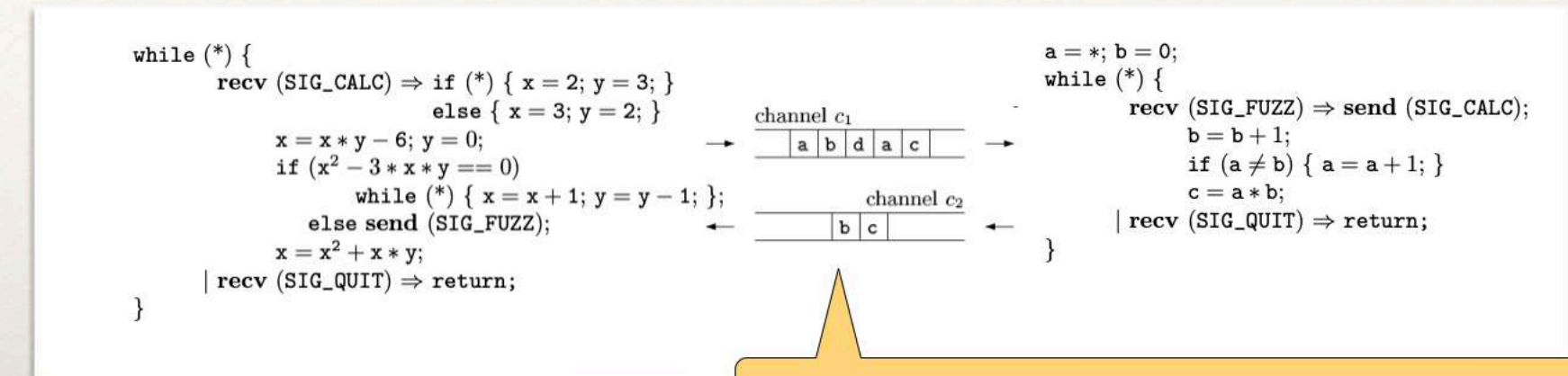
- ❖ Note that the contribution of the polynomial programs (nm) is **much lower** than the contribution of the queues ($\omega^{k-1} \cdot p$)

Concurrent polynomial programs

⊞

(JGL 2011)

- * Finite networks of polynomial programs P_1, \dots, P_m communicating through **lossy** communication queues on a finite alphabet Σ



- * State space = finite **product** of
 - **spectra** of polynomial rings $\mathbb{Z}[X_1, \dots, X_n]$, one for each P_i
 - Σ^* , with **word topology**, one for each communication queue

Our findings on statures so far

- ❖ We have already obtained **statures** of quite a few Noetherian constructions

statures

Our findings on statures so far

- ❖ We have already obtained **statures** of quite a few Noetherian constructions
- ❖ We retrieve the known formulae from wqo theory, which **extend** properly

X	$ X $
finite T_0	card X
ordinal α (Alex.)	α
$Y+Z$	$ Y \oplus Z $
$Y+_{\text{lex}}Z$	$ Y + Z $
Y_{\perp}	$1 + Y $
$Y \times Z$	$ Y \otimes Z $
fin. words Y^*	$\omega^{\{\omega^{ Y +1}\}}$
multisets Y^{\odot}	$\omega^{\tilde{\alpha}} [Y = \alpha]$

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statures

Our findings on statures so far

- ❖ We have already obtained **statures** of quite a few Noetherian constructions
- ❖ We retrieve the known formulae from wqo theory, which **extend** properly
- ❖ and **new formulae** for non-wqo Noetherian spaces

X	$ X $
finite T_0	card X
ordinal α (Alex.)	α
$Y+Z$	$ Y \oplus Z $
$Y+_{\text{lex}}Z$	$ Y + Z $
Y_{\perp}	$1 + Y $
$Y \times Z$	$ Y \otimes Z $
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cofinite topology	$\min(\text{card } Y, \omega)$
$\mathcal{H}Y, \mathcal{P}Y$	$1 + Y \dots \omega^{ Y }$
words, prefix top.	$\omega^{\{\omega^{\beta+1}\}}$ $[Y = \omega^{\{\omega^{\beta+\dots}\}} + \dots]$
$Y < \alpha$	$\leq \omega^{\{\omega^{(Y +\alpha)+1}\}}$

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- ❖ Application to actual **complexity** upper bounds?

Conclusion

Conclusion, research directions

- ❖ A rich theory extending **wqos** into the topological: **Noetherian** spaces
- ❖ Old results extend, new results pop up (powersets, spectra, infinite words)
- ❖ Ordinal analysis: the **stature** $||X||$ = ordinal rank of top element of $\mathcal{H}X$
as an analogue of maximal order types
- ❖ Still in its infancy