

THE MAXIMAL LENGTH OF A CHAIN IN THE BRUHAT ORDER FOR A CLASS OF BINARY MATRICES

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ABSTRACT: We answer to a question by Brualdi and Deaett about the maximal length of a chain in the Bruhat order for an interesting combinatorial class of binary matrices.

KEYWORDS: Bruhat order, row and column sum vectors, $(0,1)$ -matrices, interchange, minimal matrix, maximal matrix.

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1. Overview and definitions

Let m and n be two positive integers and let $R = (r_1, \dots, r_m)$ and $S = (s_1, \dots, s_n)$ be positive integral vectors. The class of all m -by- n $(0,1)$ -matrices with row sum vector R and column sum vector S is denoted by $\mathcal{A}(R, S)$. Combinatorial properties of such class of matrices have been thoroughly explored over the years (cf. e.g. [2, 3, 4, 5, 6, 7, 8, 10, 11, 15, 17] and references therein). One of the most remarkable results for $\mathcal{A}(R, S)$ known as the Gale–Ryser Theorem, independently due to David Gale [12] and Herbert J. Ryser [17], states that a necessary and sufficient condition for $\mathcal{A}(R, S)$ to be nonempty is $S \trianglelefteq R^*$, where R^* is the conjugate of R and \trianglelefteq is the standard majorization order. We are assuming that R and S are both non-increasing vectors. An important case in which nonemptiness is assured occurs when $m = n$, k is a positive integer number such that $0 \leq k \leq n$, and $R = S = (k, \dots, k)$ is the constant vector having each component equal to k . In this case we write $\mathcal{A}(n, k)$ instead of $\mathcal{A}(R, S)$.

In [8] a Bruhat partial order \preceq on a nonempty class $\mathcal{A}(R, S)$ was defined using a characterization of the Bruhat order on S_n , the symmetric group of n elements, seen as the set of permutation matrices $\mathcal{A}(n, 1)$. Precisely, for an $m \times n$ matrix $A = (a_{ij})$, let $\Sigma_A = (\sigma_{ij}(A))$ be the $m \times n$ matrix defined

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by

$$\sigma_{ij}(A) = \sum_{k=1}^i \sum_{\ell=1}^j a_{k\ell}, \quad \text{for } 1 \leq i \leq m, 1 \leq j \leq n.$$

If $A_1, A_2 \in \mathcal{A}(R, S)$, then $A_1 \preceq A_2$ if and only if $\Sigma_{A_1} \geq \Sigma_{A_2}$ in the entrywise order, i.e., $\sigma_{ij}(A_1) \geq \sigma_{ij}(A_2)$, for all $1 \leq i \leq m$ and $1 \leq j \leq n$.

In [7], Brualdi and Deaett characterize all families of the class $\mathcal{A}(n, k)$ with a unique minimal element.

Theorem 1.1 ([7, Theorem 5.1]). *Let $n \geq 1$ and k be integers with $0 \leq k \leq n$. Then the class $\mathcal{A}(n, k)$ has a unique minimal element in the Bruhat order if and only if $k \in \{0, 1, n-1, n\}$ or $n = 2k$. Moreover the minimal matrix in $\mathcal{A}(2k, k)$ is*

$$P_k = J_k \oplus J_k = \begin{pmatrix} J_k & O_k \\ O_k & J_k \end{pmatrix},$$

where J_k is the matrix of all 1's of order k and O_k is the zero matrix also of order k .

Since $\mathcal{A}(n, k) \simeq \mathcal{A}(n, n-k)$ (the map $A \mapsto J_n - A$ does the job), $|\mathcal{A}(n, 0)| = 1$, and $\mathcal{A}(n, 1) \simeq S_n$, the most interesting case is in fact $\mathcal{A}(2k, k)$.

We remark that $|\mathcal{A}(2k, k)|$ is the sequence A058527 in the *The On-Line Encyclopedia of Integer Sequences* [18]. For instances of values of such sequence the reader is referred to [19]. We observe also that computing a closed formula for such sequence is an open problem which looks quite hard (see e.g. [1, 9, 13, 14, 16] and the references therein for some asymptotic results).

In [7, Section 6] an example is provided to show that Bruhat order \preceq is not graded, and it is wondered if the maximal length of a chain in the Bruhat order in the class $\mathcal{A}(2k, k)$ from the minimal element P_k to the maximal element

$$Q_k = \begin{pmatrix} O_k & J_k \\ J_k & O_k \end{pmatrix}$$

is $4k^2$.

In this paper we prove that the above conjecture is false and that the correct answer is a much larger number: namely k^4 .

2. The main result

For any $A, B \in \mathcal{A}(R, S)$ such that $A \preceq B$, as an immediate consequence of the definition of Bruhat order, an upper bound for the length of any admissible chain between A and B is clearly given by

$$\varphi(A, B) = \sum_{i=1}^m \sum_{j=1}^n [\sigma_{ij}(A) - \sigma_{ij}(B)] .$$

Since the poset $(\mathcal{A}(2k, k), \preceq)$ admits a minimum P_k and a maximum Q_k , any chain between two elements pairwise comparable can be extended to a chain between P_k and Q_k .

After some lengthy but rather straightforward computations we get

$$\sigma_{ij}(P_k) = \begin{cases} ij & \text{if } i, j \leq k \\ ik & \text{if } i \leq k \leq j \\ jk & \text{if } i \geq k \geq j \\ ij - k(i + j - 2k) & \text{if } i, j \geq k \end{cases} ,$$

$$\sigma_{ij}(Q_k) = \begin{cases} 0 & \text{if } i, j \leq k \\ i(j - k) & \text{if } i \leq k \leq j \\ j(i - k) & \text{if } i \geq k \geq j \\ k(i + j - 2k) & \text{if } i, j \geq k \end{cases} ,$$

and $\varphi(P_k, Q_k) = k^4$.

Hence it suffices to present an instance of a chain between P_k and Q_k having exactly such length. We do that in an algorithmic way, presenting a procedure to generate an order preserving path in the Hasse diagram of $\mathcal{A}(2k, k)$.

Procedure (Switch(t, r)). $1 \leq t, r \leq 2k - 1$.

Input: $A = (a_{ij}) \in \mathcal{A}(2k, k)$ such that the submatrix

$$\begin{pmatrix} a_{t,r} & a_{t,r+1} \\ a_{t+1,r} & a_{t+1,r+1} \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} .$$

Output: $B = (b_{ij}) \in \mathcal{A}(2k, k)$ such that $b_{ij} = a_{ij}$ if $1 \leq i, j \leq 2k$ and $(i, j) \notin \{(t, r), (t, r + 1), (t + 1, r), (t + 1, r + 1)\}$, and

$$\begin{pmatrix} b_{t,r} & b_{t,r+1} \\ b_{t+1,r} & b_{t+1,r+1} \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} .$$

It is easy to see that executing procedure $\text{Switch}(t, r)$ the output covers the input in the Bruhat order for any choice of parameters. Our chain will be made by repeated applications of the procedure $\text{Switch}(t, r)$.

Procedure ($\text{Switch-rows}(t)$). $1 \leq t \leq 2k - 1$.

Input: $A = (a_{ij}) \in \mathcal{A}(2k, k)$ such that rows $t, t + 1$ equal

$$\begin{pmatrix} 1, \dots, 1, 0, \dots, 0 \\ 0, \dots, 0, 1, \dots, 1 \end{pmatrix}.$$

For $\alpha = k$ down to 1 do

 Begin

 For $\beta = \alpha$ to $\alpha + k - 1$ do $\text{Switch}(t, \beta)$.

 End.

Output: $B = (b_{ij}) \in \mathcal{A}(2k, k)$ such that $b_{ij} = a_{ij}$ for any $1 \leq i, j \leq 2k$ such that $i \neq t, t + 1$, and rows $t, t + 1$ equal

$$\begin{pmatrix} 0, \dots, 0, 1, \dots, 1 \\ 1, \dots, 1, 0, \dots, 0 \end{pmatrix}.$$

Algorithm 2.1 ($\text{Chain}(k)$). $0 \neq k \in \mathbb{N}$.

Input: P_k .

For $\alpha = k$ down to 1 do

 Begin

 For $\beta = \alpha$ to $\alpha + k - 1$ do $\text{Switch-rows}(\beta)$.

 End.

Output: Q_k .

We can see that, for any choice of parameters, the procedure “Switch” is invoked k^2 times by procedure “Switch-rows”, and that algorithm “Chain” recalls procedure “Switch-rows” k^2 times as well, so there are k^4 application of procedure “Switch”. Since, as already remarked, each time that procedure “Switch” is recalled we are moving up (by one cover relation) in the Hasse diagram of the poset $(\mathcal{A}(2k, k), \preceq)$, all the constructed elements are pairwise distinct members of the desired chain. As a consequence we obtain our result.

Theorem 2.2. *For any positive integer k , the maximal length of a chain in the Bruhat order in $\mathcal{A}(2k, k)$ equals k^4 .*

For the sake of clarity, we present in detail our construction of the chain for the case $k = 2$.

Example 2.1. $k = 2$.

$$\begin{array}{c}
 \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 \end{pmatrix} \mapsto \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \end{pmatrix} \mapsto \begin{pmatrix} 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{pmatrix} \\
 \mapsto \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{pmatrix} \mapsto \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \end{pmatrix} \mapsto \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \end{pmatrix} \\
 \mapsto \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{pmatrix} \mapsto \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{pmatrix} \mapsto \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix} \\
 \mapsto \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix} \mapsto \begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix} \mapsto \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix} \\
 \mapsto \begin{pmatrix} 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix} \mapsto \begin{pmatrix} 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 \end{pmatrix} \mapsto \begin{pmatrix} 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 \end{pmatrix} \\
 \mapsto \begin{pmatrix} 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 \end{pmatrix} \mapsto \begin{pmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \end{pmatrix}.
 \end{array}$$

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