

ON GENERALIZED EQUILOGICAL SPACES

WILLIAN RIBEIRO

ABSTRACT: In this paper we carry the construction of equiological spaces into an arbitrary category \mathbf{X} topological over \mathbf{Set} , introducing the category $\mathbf{X}\text{-Equ}$ of equiological objects. Similar to what is done for the category \mathbf{Top} of topological spaces and continuous functions, we study some features of the new category as (co)completeness and regular (co-)well-poweredness, as well as the fact that, under some conditions, it is a quasitopos. We achieve these various properties of the category $\mathbf{X}\text{-Equ}$ by representing it as a category of partial equiological objects, as a reflective subcategory of the exact completion \mathbf{X}_{ex} , and as the regular completion \mathbf{X}_{reg} . We finish with examples in the particular cases, amongst others, of ordered, metric, and approach spaces, which can all be described using the $(\mathbb{T}, \mathbf{V})\text{-Cat}$ setting.

KEYWORDS: equiological space, topological category, exact completion, regular completion, quasitopos, $(\mathbb{T}, \mathbf{V})\text{-category}$, modest set.

MATH. SUBJECT CLASSIFICATION (2010): 54A05, 54B30, 18D15, 54D80, 18B35.

Introduction

As a solution to remedy the problem of non-existence of general exponentials in \mathbf{Top} , Scott presents first in [Sco96], and later with his co-authors Bauer and Birkedal in [BBS04], the category \mathbf{Equ} of equiological spaces. Formed by equipping topological T_0 -spaces with arbitrary equivalence relations, \mathbf{Equ} contains \mathbf{Top}_0 (T_0 -spaces and continuous functions) as a full subcategory and it is cartesian closed. This fact is directly proven by showing an equivalence with the category \mathbf{PEqu} of partial equiological spaces, which is formed by equipping algebraic lattices with partial (not necessarily reflexive) equivalence relations. Also in [BBS04], equiological spaces are presented as modest sets of assemblies over algebraic lattices, offering a model for dependent type theory.

Contributing to the study of \mathbf{Equ} , a more general categorical framework, explaining why such (sub)categories are (locally) cartesian closed, was presented in [BCRS98, CR00, Ros99]. It turned out that \mathbf{Equ} is related to the free exact completion $(\mathbf{Top}_0)_{\text{ex}}$ of \mathbf{Top}_0 [Car95, CM82, CV98]. By the same

Received November 19, 2018.

Research supported by Centro de Matemática da Universidade de Coimbra – UID/MAT/00324/2013 and by the FCT PhD grant PD/BD/128059/2016, funded by the Portuguese Government through FCT/MCTES and co-funded by the European Regional Development Fund through the Partnership Agreement PT2020.

token, suppressing the T_0 -separation condition on the topological spaces, the category **Equ** is a full reflective subcategory of the exact completion \mathbf{Top}_{ex} of **Top**. More, the reflector preserves products and special pullbacks, from where it is concluded in [BCRS98] that **Equ** is locally cartesian closed, since \mathbf{Top}_{ex} is so [BCRS98, Theorem 4.1]. It is shown in [Ros98] that **Equ** can be presented as the free regular completion of **Top** [Car95, CV98], and [Men00] provides conditions for such regular completions to be quasitoposes.

In this paper we start with a category \mathbf{X} and a topological functor $|-|: \mathbf{X} \rightarrow \mathbf{Set}$, and, equipping each object X of \mathbf{X} with an equivalence relation on its underlying set $|X|$, we define the category $\mathbf{X}\text{-Equ}$ of equiological objects and their morphisms. Recovering the results for the particular case of **Top**, $\mathbf{X}\text{-Equ}$ is (co)complete and regular (co-)well-powered. Under the hypothesis of pre-order-enrichment, we explore the concepts of separated and injective objects of \mathbf{X} , leading us to the definition of a category $\mathbf{X}\text{-PEqu}$ of partial equiological objects. In the presence of a separation condition, we proceed presenting $\mathbf{X}\text{-Equ}$ as modest sets of assemblies over injective objects; from that, we verify its properties of cartesian closedness and regularity. This is the subject of our first section.

In Section 2, analogously to the case of **Top**, we get similar results when considering the exact completion \mathbf{X}_{ex} and the regular completion \mathbf{X}_{reg} of \mathbf{X} , culminating in the fact that $\mathbf{X}\text{-Equ}$ is a quasitopos, by the results of [Men00]. To do so, we use a general approach to study weak cartesian closedness of topological categories (see [CHR18]).

We finish with Section 3, where we briefly recall the $(\mathbb{T}, \mathbb{V})\text{-Cat}$ setting, which was introduced in [CT03] and further investigated in other papers [CH03, Hof07], and study the case when $\mathbf{X} = (\mathbb{T}, \mathbb{V})\text{-Cat}$, for a suitable monad \mathbb{T} and quantale \mathbb{V} , satisfying all conditions needed throughout the paper. Examples of such categories are **Ord** of preordered sets, **Met** of Lawvere generalized metric spaces [Law02] and **App** of Lowen approach spaces [Low97], amongst others. Based on full embeddings among those categories, we place full embeddings among their categories of equiological objects.

1. The category of (partial) equiological objects

Let \mathbf{X} be a category and $|-|: \mathbf{X} \rightarrow \mathbf{Set}$ be a topological functor. In particular \mathbf{X} is complete, cocomplete, and $|-|$ preserves both limits and colimits.

Definition 1.1. The category $\mathbf{X}\text{-Equ}$ is defined as follows.

- The objects are structures $\mathcal{X} = \langle X, \equiv_{|X|} \rangle$, where $X \in \mathbf{X}$ and $\equiv_{|X|}$ is an

equivalence relation on the set $|X|$; they are called *equiological objects of \mathbf{X}* .

• A morphism from $\mathcal{X} = \langle X, \equiv_{|X|} \rangle$ to $\mathcal{Y} = \langle Y, \equiv_{|Y|} \rangle$ is the equivalence class of a morphism $f: X \rightarrow Y$ in \mathbf{X} such that $|f|$ is an *equivariant* map, i.e. $x \equiv_{|X|} x'$ implies $|f|(x) \equiv_{|Y|} |f|(x')$, for all $x, x' \in |X|$, with the equivalence relation on morphisms defined by

$$f \equiv_{\mathcal{X} \rightarrow \mathcal{Y}} g \iff \forall x, x' \in |X|, (x \equiv_{|X|} x' \implies |f|(x) \equiv_{|Y|} |g|(x')).$$

One can see that $\equiv_{\mathcal{X} \rightarrow \mathcal{Y}}$ is indeed an equivalence relation; reflexivity follows from the fact that the underlying maps are equivariant, symmetry and transitivity follow from the same properties for $\equiv_{|X|}$ and $\equiv_{|Y|}$.

Identity of \mathcal{X} is given by $[1_X]$ and composition of classes $[f]: \mathcal{X} \rightarrow \mathcal{Y}$ and $[g]: \mathcal{Y} \rightarrow \mathcal{Z}$ is given by $[g] \cdot [f] = [g \cdot f]$, which is well-defined.

Theorem 1.1. *$\mathbf{X}\text{-Equ}$ is complete, cocomplete, regular well-powered and regular co-well-powered.*

The proof of Theorem 1.1 goes along the general lines of the proof of [BBS04, Theorem 3.10]. Limits and colimits are computed in \mathbf{X} and their underlying sets are endowed with appropriate equivalence relations. The properties of regular well- and regular co-well-poweredness follow from the description of equalizers and coequalizers in $\mathbf{X}\text{-Equ}$.

In general $\mathbf{X}\text{-Equ}$ is neither well-powered nor co-well-powered, as observed in [BBS04] for topological spaces. A morphism $[m]: \mathcal{X} \rightarrow \mathcal{Y}$ is a monomorphism in $\mathbf{X}\text{-Equ}$ if, and only if,

$$x \equiv_{|X|} x' \iff |m|(x) \equiv_{|Y|} |m|(x'), \forall x, x' \in |X|.$$

A morphism $[f]: \mathcal{X} \rightarrow \mathcal{Y}$ is an epimorphism in $\mathbf{X}\text{-Equ}$ if, and only if,

$$y \equiv_{|Y|} y' \iff \exists x, x' \in |X|; x \equiv_{|X|} x' \ \& \ y \equiv_{|Y|} |f|(x) \equiv_{|Y|} |f|(x') \equiv_{|Y|} y'.$$

Having the *embedding* and the *extension* theorems configured for powersets [BBS04, Theorems 3.6, 3.7], according to the authors, Scott has pointed out that those results in fact hold more generally to continuous lattices. Powersets can be generalized to algebraic lattices, and it is explained that “*The reason for considering algebraic lattices is that the lattice of continuous functions between powerset spaces is not usually a powerset space, but it is an algebraic lattice. And this extends to all algebraic lattices.*”, culminating in the well known fact that the category \mathbf{ALat} of algebraic lattices and Scott-continuous functions is cartesian closed [GHK+80, Chapter II, Theorem 2.10].

Algebraic lattices are in particular continuous lattices, therefore injective objects in \mathbf{Top}_0 . We show below that these – injectivity and separation – are the crucial properties in order to extend the arguments of [BBS04]. Next we assume that

(a) \mathbf{X} is a pre-order enriched category.

Definition 1.2. An object X of \mathbf{X} is said to be *separated* if, for each morphisms $f, g: Y \rightarrow X$ in \mathbf{X} , whenever $f \simeq g$ ($f \leq g$ and $g \leq f$), then $f = g$.

Hence an object X is separated if, for each object Y , the pre-ordered set of morphisms $\mathbf{X}(Y, X)$ is anti-symmetric. One can check that this is equivalent to the pre-ordered set $\mathbf{X}(\mathbf{1}, X)$ to be anti-symmetric, where $\mathbf{1} = L\{*\}$, with $L: \mathbf{Set} \rightarrow \mathbf{X}$ the left adjoint of $|-|$.

The full subcategory \mathbf{X}_{sep} of separated objects is replete and closed under mono-sources. Since $(\text{RegEpi}, \mathbb{M})$ is a factorization system for sources in the topological category \mathbf{X} , where \mathbb{M} stands for the class of mono-sources [AHS90, Proposition 21.14], closure under mono-sources then implies that \mathbf{X}_{sep} is regular epi-reflective in \mathbf{X} [HST14, II-Proposition 5.10.1].

We will consider (pseudo-)injective objects of \mathbf{X} with respect to $|-|$ -initial morphisms. Hence, denoting by \mathbf{X}_{inj} the full subcategory on the injectives, $Z \in \mathbf{X}_{\text{inj}}$ if, and only if, for each $|-|$ -initial morphism $y: X \rightarrow Y$ and morphism $f: X \rightarrow Z$, there exists a morphism $\hat{f}: Y \rightarrow Z$ such that $\hat{f} \cdot y \simeq f$;

$$\begin{array}{ccc} X & \xrightarrow{y} & Y \\ & \searrow f & \swarrow \hat{f} \\ & & Z \end{array} \quad \begin{array}{c} \simeq \\ \simeq \end{array}$$

\hat{f} is called an *extension* of f along y ; if Z is separated, then $\hat{f} \cdot y = f$. Moreover, assume that

(b) for each $X, Y \in \mathbf{X}$, $x, x' \in \mathbf{X}(\mathbf{1}, X)$ and $|-|$ -initial $f \in \mathbf{X}(X, Y)$,

$$f \cdot x \simeq f \cdot x' \implies x \simeq x'.$$

Thus, for X separated, if f is $|-|$ -initial, then f is an embedding (regular monomorphism, which in our case is equivalent to $|-|$ -initial with $|f|$ an injective map); hence restricting ourselves to the separated objects, injectivity with respect to $|-|$ -initial morphisms coincides with injectivity with respect to embeddings.

Definition 1.3. The category $\mathbf{X}\text{-PEqu}$ of *partial equiological objects of \mathbf{X}* consists of:

- objects are structures $\mathcal{X} = \langle X, \equiv_{|X|} \rangle$, where $X \in \mathbf{X}_{\text{inj}}$ and $\equiv_{|X|}$ is a partial (not necessarily reflexive) equivalence relation on the set $|X|$;
- a morphism from $\langle X, \equiv_{|X|} \rangle$ to $\langle Y, \equiv_{|Y|} \rangle$ is the equivalence class of an \mathbf{X} -morphism $f: X \rightarrow Y$ such that $|f|$ is an equivariant map, with the equivalence relation on morphisms as in Definition 1.1.

In order to verify an equivalence between the categories of equiological and partial equiological objects, we will restrict ourselves to the separated objects, so we consider now that the objects in the structures of Definitions 1.1 and 1.3 are all separated, and denote the resulting categories by $\mathbf{X}\text{-Equ}_{\text{sep}}$ and $\mathbf{X}\text{-PEqu}_{\text{sep}}$, respectively. We also assume that

(c) \mathbf{X} has enough injectives, meaning that, for each $X \in \mathbf{X}$, there exists an $|-|$ -initial morphism $y_X: X \rightarrow \hat{X}$, with $\hat{X} \in \mathbf{X}_{\text{inj}}$, and, if X is separated, so is \hat{X} .

When X is separated, as we have seen before, y_X is an embedding.

Theorem 1.2. $\mathbf{X}\text{-Equ}_{\text{sep}}$ and $\mathbf{X}\text{-PEqu}_{\text{sep}}$ are equivalent.

Proof. As in the proof of [BBS04, Theorem 3.12], a functor $R: \mathbf{X}\text{-PEqu}_{\text{sep}} \rightarrow \mathbf{X}\text{-Equ}_{\text{sep}}$ is defined taking each separated partial equiological object \mathcal{X} to $R\mathcal{X} = \langle RX, \equiv_{|RX|} \rangle$, where $|RX| = \{x \in |X| \mid x \equiv_{|X|} x\}$ and $\equiv_{|RX|}$ is the restriction of $\equiv_{|X|}$. For a morphism $[f]: \mathcal{X} \rightarrow \mathcal{Y}$, $|f|(|RX|) \subseteq |RY|$, so we take the (co)restriction $\overline{|f|}: |RX| \rightarrow |RY|$, lift to an \mathbf{X} -morphism $\overline{f}: RX \rightarrow RY$ and set $R[f] = \overline{[f]}$.

To prove that R is faithful one only needs to observe that, for elements $x, x' \in X$, if $x \equiv_{|X|} x'$, then $x' \equiv_{|X|} x$, and consequently $x \equiv_{|X|} x$ and $x' \equiv_{|X|} x'$, whence $x, x' \in |RX|$; and to prove that R is full one uses the injectivity of Y , providing an extension $\hat{f}: X \rightarrow Y$ of $i_{RY} \cdot f$ along i_{RX} .

Finally, for essential surjectivity let $\mathcal{X} = \langle X, \equiv_{|X|} \rangle \in \mathbf{X}\text{-Equ}_{\text{sep}}$ and consider the embedding $y_X: X \rightarrow \hat{X}$, $\hat{X} \in \mathbf{X}_{\text{sep, inj}}$. Endow $|\hat{X}|$ with the following partial equivalence relation

$$\varphi \equiv_{|\hat{X}|} \psi \iff \exists x, x' \in |X|; \varphi = |y_X|(x), \psi = |y_X|(x') \ \& \ x \equiv_{|X|} x',$$

that is, two elements of $|\hat{X}|$ are related if, and only if, they are the images by $|y_x|$ of elements that are related in $|X|$. The sets $|R\hat{X}|$ and $|X|$ are in bijection; using the $|-|$ -initiality of y_x and $i_{R\hat{X}}$, this bijection proves to be an isomorphism in \mathbf{X} , and consequently in $\mathbf{X}\text{-Equ}_{\text{sep}}$, by the definition of $\equiv_{|\hat{X}|}$, so $R\langle\hat{X}, \equiv_{|\hat{X}|}\rangle \cong \mathcal{X}$. \blacksquare

For our next result we must assume that

(d) *every injective object of \mathbf{X} is exponentiable.*

Binary products and exponentials of injective objects are again injective, so \mathbf{X}_{inj} is a cartesian closed subcategory of \mathbf{X} . We also assume that

(e) *the reflector from \mathbf{X} to \mathbf{X}_{sep} preserves finite products;*

whence the exponential of separated objects, when it exists, is again separated [Day72, Sch84].

Theorem 1.3. $\mathbf{X}\text{-PEqu}_{\text{sep}}$ *is cartesian closed.*

Proof: Let $\mathcal{X} = \langle X, \equiv_{|X|} \rangle$ and $\mathcal{Y} = \langle Y, \equiv_{|Y|} \rangle$ be partial equiological separated objects. We build the exponential Y^X in $\mathbf{X}_{\text{sep, inj}}$ and endow $|Y^X|$ with the partial equivalence relation: $\alpha \equiv_{|Y^X|} \beta$ if, and only if, for all $x, x' \in X$,

$$x \equiv_{|X|} x' \implies \alpha(x) = |\text{ev}|(\alpha, x) \equiv_{|Y|} |\text{ev}|(\beta, x') = \beta(x'),$$

for each $\alpha, \beta \in |Y^X|$, where $\text{ev}: Y^X \times X \rightarrow Y$ is the evaluation morphism in \mathbf{X} . Then $\mathcal{Y}^{\mathcal{X}} = \langle Y^X, \equiv_{|Y^X|} \rangle \in \mathbf{X}\text{-PEqu}_{\text{sep}}$ and $|\text{ev}|: |Y^X| \times |X| \rightarrow |Y|$ is equivariant, so $[\text{ev}]: \mathcal{Y}^{\mathcal{X}} \times \mathcal{X} \rightarrow \mathcal{Y}$ is a valid morphism in $\mathbf{X}\text{-PEqu}_{\text{sep}}$. More, $[\text{ev}]$ satisfies the universal property: for each morphism $[f]: \mathcal{Z} \times \mathcal{X} \rightarrow \mathcal{Y}$, $\mathcal{Z} = \langle Z, \equiv_{|Z|} \rangle \in \mathbf{X}\text{-PEqu}_{\text{sep}}$, there exists a unique $[\bar{f}]: \mathcal{Z} \rightarrow \mathcal{Y}^{\mathcal{X}}$ commuting the diagram below.

$$\begin{array}{ccc} \mathcal{Y}^{\mathcal{X}} \times \mathcal{X} & \xrightarrow{[\text{ev}]} & \mathcal{Y} \\ \uparrow [\bar{f}] \times 1_{\mathcal{X}} & & \nearrow [f] \\ \mathcal{Z} \times \mathcal{X} & & \end{array}$$

The morphism $\bar{f}: Z \rightarrow Y^X$ is the transpose of $f: Z \times X \rightarrow Y$, so that

$$\text{ev} \cdot (\bar{f} \times 1_X) = f,$$

and $[\bar{f}]$ is indeed unique, for if $[f']: \mathcal{Z} \rightarrow \mathcal{Y}^{\mathcal{X}}$ is such that $[\text{ev} \cdot (f' \times 1_X)] = [f]$, then for each $z \equiv_{|Z|} z'$ in Z and $x \equiv_{|X|} x'$ in X ,

$$\begin{aligned} \bar{f}(z)(x) &= \text{ev} \cdot (\bar{f} \times 1_X)(z, x) = f(z, x) \equiv_{|Y|} f(z', x') \\ &\equiv_{|Y|} \text{ev} \cdot (f' \times 1_X)(z', x') = f'(z')(x'), \end{aligned}$$

hence $\bar{f}(z) \equiv_{|Y^X|} f'(z')$, i.e. $[\bar{f}] = [f']$. ■

Therefore, by Theorem 1.2, $\mathbf{X}\text{-Equ}_{\text{sep}}$ is cartesian closed. We remark that the proof of Theorem 1.3 also applies to $\mathbf{X}\text{-PEqu}$ without separation.

To finish this section we discuss the presentation of equiological objects as modest sets of assemblies, following what is done in [BBS04, Section 4].

Definition 1.4. The category of *assemblies* $\text{Assm}(\mathbf{X}_{\text{inj}})$ over injective objects of \mathbf{X} consists of the following data: objects are triples (A, X, E_A) , where A is a set, $X \in \mathbf{X}_{\text{inj}}$, and $E_A: A \rightarrow \mathcal{P}|X|$ is a function such that $E_A(a) \neq \emptyset$, for each $a \in A$, with $\mathcal{P}|X|$ the powerset of $|X|$. The elements of $E_A(a)$ are called *realizers* for a . A morphism between assemblies (A, X, E_A) and (B, Y, E_B) is a map $f: A \rightarrow B$ for which there exists a morphism $g: X \rightarrow Y$ in \mathbf{X} such that $|g|(E_A(a)) \subseteq E_B(f(a))$; g is said to be a *realizer* for f , and we say that $|g|$ *tracks* f .

Definition 1.5. An object $(A, X, E_A) \in \text{Assm}(\mathbf{X}_{\text{inj}})$ is called a *modest set* if, for all $a, a' \in A$, $a \neq a'$ implies $E_A(a) \cap E_A(a') = \emptyset$. The full subcategory of the assemblies that are modest sets is denoted by $\text{Mdst}(\mathbf{X}_{\text{inj}})$.

With these definitions, we get the same properties as those for the particular case of topological spaces, which we highlight in the following items, omitting some of the proofs that follow directly from the ones in [BBS04].

(1) $\text{Mdst}(\mathbf{X}_{\text{inj}})$ and $\text{Assm}(\mathbf{X}_{\text{inj}})$ have finite limits and the inclusion from modest sets to assemblies preserves them.

(2) $\text{Mdst}(\mathbf{X}_{\text{inj}})$ and $\text{Assm}(\mathbf{X}_{\text{inj}})$ are cartesian closed and $\text{Mdst}(\mathbf{X}_{\text{inj}}) \rightarrow \text{Assm}(\mathbf{X}_{\text{inj}})$ preserves exponentials. For (A, X, E_A) and (B, Y, E_B) in $\text{Assm}(\mathbf{X}_{\text{inj}})$, the exponential is (C, Y^X, E_C) , where $C = \{f: A \rightarrow B \mid \exists g: X \rightarrow Y \in \mathbf{X} \text{ realizer for } f\}$, and $E_C(f) = \{\alpha \in |Y^X| \mid \alpha \text{ tracks } f\}$; here, for simplicity, we denote also by α the map from $|X|$ to $|Y|$, given by $x \mapsto |\text{ev}|(\alpha, x)$,

for each $x \in |X|$, where $\text{ev}: Y^X \times X \rightarrow Y$ is the evaluation map. If (B, Y, E_B) is a modest set, then so is (C, Y^X, E_C) , for if $f, f': A \rightarrow B$ are tracked by $\alpha \in |Y^X|$, then for each $a \in A$, take $x \in E_A(a) \neq \emptyset$, then $\alpha(x) \in E_B(f(a)) \cap E_B(f'(a)) \neq \emptyset$, whence $f(a) = f'(a)$ and then $f = f'$.

(3) $\text{Mdst}(\mathbf{X}_{\text{inj}})$ is a reflective subcategory of $\text{Assm}(\mathbf{X}_{\text{inj}})$.

(4) The regular subobjects of (A, X, E_A) in $\text{Assm}(\mathbf{X}_{\text{inj}})$, or in $\text{Mdst}(\mathbf{X}_{\text{inj}})$, are in bijective correspondence with the powerset of A .

(5) $\text{Mdst}(\mathbf{X}_{\text{inj}})$ and $\text{Assm}(\mathbf{X}_{\text{inj}})$ are regular categories.

Theorem 1.4. $\mathbf{X}\text{-PEqu}$ and $\text{Mdst}(\mathbf{X}_{\text{inj}})$ are equivalent.

Proof. Define the functor $F: \text{Mdst}(\mathbf{X}_{\text{inj}}) \rightarrow \mathbf{X}\text{-PEqu}$ assigning to (A, X, E_A) the object $\langle X, \equiv_{|X|} \rangle$, where

$$x \equiv_{|X|} x' \iff \exists a \in A; x, x' \in E_A(a),$$

and on morphisms F assigns to each $f: (A, X, E_A) \rightarrow (B, Y, E_B)$ the class of a realizer $g: X \rightarrow Y$ for f ; $\equiv_{|X|}$ is indeed an equivalence relation and two realizers for f are in the same equivalence class, so F is well-defined.

Faithfulness of F follows from the observation in item (2) above: two maps tracked by the same realizer must be equal. To see that F is full, take a morphism $[g]: F(A, X, E_A) \rightarrow F(B, Y, E_B)$ in $\mathbf{X}\text{-PEqu}$. For each $a \in A$, let $x \in E_A(a) \neq \emptyset$; then $x \equiv_{|X|} x$ and so $|g|(x) \equiv_{|Y|} |g|(x)$, that is, there exists $b \in B$ such that $|g|(x) \in E_B(b)$, whence we set $f(a) = b$; this b is uniquely determined since (B, Y, E_B) is a modest set, therefore we have a map $f: A \rightarrow B$, which, by definition, has g as a realizer.

Now let $\langle X, \equiv_{|X|} \rangle$ be a partial equiological object and define (A, X, E_A) by $A = \{x \in |X| \mid x \equiv_{|X|} x\} / \equiv_{|X|}$ and $E_A([x]) = [x] \subseteq \mathcal{P}|X|$. Hence $F(A, X, E_A) = \langle X, \equiv_{|X|} \rangle$ and F is essentially surjective. \blacksquare

The same argument can be repeated replacing \mathbf{X}_{inj} with $\mathbf{X}_{\text{sep, inj}}$, so we obtain $\text{Mdst}(\mathbf{X}_{\text{sep, inj}}) \cong \mathbf{X}\text{-PEqu}_{\text{sep}} \cong \mathbf{X}\text{-Equ}_{\text{sep}}$. Properties from items (1) to (5) remain valid, so they also hold for $\mathbf{X}\text{-Equ}_{\text{sep}}$.

2. Equiological objects and exact completion

The category Equ of equiological spaces can also be obtained as a full reflective subcategory of the exact completion [BCRS98, Car95, CM82] Top_{ex} of

the category of topological spaces [Ros99], and this is a particular instance of a general process to obtain such categories [BCRS98].

We can describe the exact completion \mathbf{X}_{ex} of \mathbf{X} as: objects are *pseudo-equivalence relations* on \mathbf{X} , that is, parallel pairs of morphisms $X_1 \begin{array}{c} \xrightarrow{r_1} \\ \xrightarrow{r_2} \end{array} X_0$ of \mathbf{X} satisfying

(i) *reflexivity*: there exists a morphism $r: X_0 \rightarrow X_1$ such that $r_1 \cdot r = 1_{X_0} = r_2 \cdot r$;

$$\begin{array}{ccccc} & & X_1 & & \\ & r_1 \swarrow & \uparrow r & \searrow r_2 & \\ X_0 & \xleftarrow{1_{X_0}} & X_0 & \xrightarrow{1_{X_0}} & X_0 \end{array}$$

(ii) *symmetry*: there exists a morphism $s: X_1 \rightarrow X_1$ such that $r_1 \cdot s = r_2$ and $r_2 \cdot s = r_1$;

$$\begin{array}{ccccc} & & X_1 & & \\ & r_1 \swarrow & \uparrow s & \searrow r_2 & \\ X_0 & \xleftarrow{r_2} & X_1 & \xrightarrow{r_1} & X_0 \end{array}$$

(iii) *transitivity*: for $r_3, r_4: X_2 \rightarrow X_1$ a pullback of r_1, r_2 , there exists a morphism $t: X_2 \rightarrow X_1$ commuting the following diagram

$$\begin{array}{ccccc} & & X_1 & & \\ & & \uparrow t & & \\ & & X_2 & & \\ & r_1 \swarrow & \searrow r_3 & \swarrow r_4 & \searrow r_2 \\ & X_1 & \searrow r_2 & X_1 & \\ & \swarrow r_1 & \searrow r_2 & \swarrow r_1 & \searrow r_2 \\ X_0 & & X_0 & & X_0 \end{array}$$

A morphism from $X_1 \begin{array}{c} \xrightarrow{r_1} \\ \xrightarrow{r_2} \end{array} X_0$ to $Y_1 \begin{array}{c} \xrightarrow{s_1} \\ \xrightarrow{s_2} \end{array} Y_0$ is an equivalence classe $[f]$ of an \mathbf{X} -morphism $f: X_0 \rightarrow Y_0$ such that there exists $g: X_1 \rightarrow Y_1$ in \mathbf{X} satisfying

$$f \cdot r_i = s_i \cdot g, \quad i = 1, 2.$$

$$\begin{array}{ccc} X_1 & \xrightarrow{g} & Y_1 \\ r_1 \downarrow & & \downarrow s_1 \\ & & r_2 \downarrow & & \downarrow s_2 \\ X_0 & \xrightarrow{f} & Y_0. \end{array}$$

Here two morphisms $f_1, f_2: X_0 \rightarrow Y_0$ are related if, and only if, there exists a morphism $h: X_0 \rightarrow Y_1$ such that $f_i = s_i \cdot h$, $i = 1, 2$.

$$\begin{array}{ccc} X_1 & & Y_1 \\ \downarrow r_1 & \nearrow h & \downarrow s_1 \\ & & r_2 \downarrow & & \downarrow s_2 \\ X_0 & \xrightarrow{f_1} & Y_0 \\ & \xrightarrow{f_2} & \end{array}$$

Since it is topological over \mathbf{Set} , \mathbf{X} has a stable factorization system for morphisms given by (Epi, RegMono) [AHS90, Remark 15.2(3), Proposition 21.14] (see also [CHR18]). Let $\text{PER}(\mathbf{X}, \text{RegMono})$ denote the full subcategory of \mathbf{X}_{ex} of the pseudo-equivalence relations $X_1 \xrightarrow[r_2]{r_1} X_0$ such that $\langle r_1, r_2 \rangle: X_1 \rightarrow X_0 \times X_0$ is a regular monomorphism.

Lemma 2.1. *\mathbf{X} -Equ and $\text{PER}(\mathbf{X}, \text{RegMono})$ are equivalent.*

Proof. For each equilogical object $\langle X, \equiv_{|X|} \rangle$, consider $E_X = \{(x, x') \in |X| \times |X| \mid x \equiv_{|X|} x'\}$ and the source $(\pi_i^X: E_X \rightarrow |X|)_{i=1,2}$ of the projections from E_X onto $|X|$. Take its $|-|$ -initial lifting, which by abuse of notation we denote

by $(\pi_i^X: E_X \rightarrow X)_{i=1,2}$. Hence $E_X \xrightarrow[\pi_2^X]{\pi_1^X} X$ belongs to $\text{PER}(\mathbf{X}, \text{RegMono})$ and

each morphism $[f]: \mathcal{X} \rightarrow \mathcal{Y}$ in $\mathbf{X}\text{-Equ}$ is a valid morphism

$$[f]: (E_X, X, \pi_1^X, \pi_2^X) \rightarrow (E_Y, Y, \pi_1^Y, \pi_2^Y)$$

in $\text{PER}(\mathbf{X}, \text{RegMono})$.

That correspondence defines a functor which is fully faithful and, for a pseudo-equivalence relation $X_1 \xrightarrow[r_2]{r_1} X_0$ in $\text{PER}(\mathbf{X}, \text{RegMono})$, define the

equiological object $\langle X_0, \equiv_{|X_0|} \rangle$ by

$$x_0 \equiv_{|X_0|} x'_0 \iff (\exists \text{ (unique) } x_1 \in X_1) |r_1|(x_1) = x_0 \ \& \ |r_2|(x_1) = x'_0,$$

for each $x_0, x'_0 \in X_0$. Then $E_{X_0} \begin{array}{c} \xrightarrow{\pi_1^{X_0}} \\ \xrightarrow{\pi_2^{X_0}} \end{array} X_0$ is isomorphic to $X_1 \begin{array}{c} \xrightarrow{r_1} \\ \xrightarrow{r_2} \end{array} X_0$ in

$\text{PER}(\mathbf{X}, \text{RegMono})$ and the functor is essentially surjective. \blacksquare

Hence [BCRS98, Theorem 4.3] states that

Theorem 2.1. $\mathbf{X}\text{-Equ} \cong \text{PER}(\mathbf{X}, \text{RegMono})$ is a full reflective subcategory of \mathbf{X}_{ex} ; the reflector preserves finite products and commutes with change of base in the codomain.

Next we wish to prove that $\mathbf{X}\text{-Equ}$ is cartesian closed, so by Theorem 2.1 and [Sch84, Theorem 1.2], it suffices to show that \mathbf{X}_{ex} is cartesian closed. To do so, we will apply the following result derived from [Ros99, Theorem 1, Lemma 4] (see also [CHR18, Theorem 1.1]).

Theorem 2.2. Let \mathbf{X} be a complete, infinitely extensive and well-powered category with $(\text{RegEpi}, \text{Mono})$ -factorizations such that $f \times 1$ is an epimorphism whenever f is a regular epimorphism. Then \mathbf{X}_{ex} is cartesian closed provided \mathbf{X} is weakly cartesian closed.

Since \mathbf{X} is topological over Set , in order to use the latter theorem, we will assume that

(f) \mathbf{X} is infinitely extensive;

more, assuming also conditions (a) to (e) from the previous section, following the same steps of the proofs of [CHR18, Theorems 5.3 and 5.5], we deduce the following result.

Proposition 2.1. \mathbf{X} is weakly cartesian closed.

Furthermore, we can verify that \mathbf{X}_{ex} is actually locally cartesian closed. Consider the restriction functor $|-|: \mathbf{X}_{\text{inj}} \rightarrow \text{Pfn}$, where Pfn is the category of sets and partial functions. The category $\mathcal{F}(\mathbf{X}_{\text{inj}}, |-|)$, or simply $\mathcal{F}(\mathbf{X}_{\text{inj}})$, is described in [CR00] as follows: objects are triples $(X, A, \sigma: A \rightarrow |X|)$, where X is an injective object of \mathbf{X} , A is a set and σ is a function; a morphism

$f: (X, A, \sigma: A \rightarrow |X|) \rightarrow (Y, B, \delta: B \rightarrow |Y|)$ is a map $f: A \rightarrow B$ such that there exists $g: X \rightarrow Y$ in \mathbf{X} commuting the diagram

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ \sigma \downarrow & & \downarrow \delta \\ |X| & \xrightarrow{|g|} & |Y|. \end{array}$$

Proposition 2.2. *The categories \mathbf{X} and $\mathcal{F}(\mathbf{X}_{\text{inj}})$ are equivalent.*

Proof. Define the functor $G: \mathbf{X} \rightarrow \mathcal{F}(\mathbf{X}_{\text{inj}})$ by

$$GX = (\hat{X}, |X|, \sigma_x = |y_x|: |X| \rightarrow |\hat{X}|),$$

where y_x is the $|-|$ -initial morphism assured by condition **(c)** in the previous section; for each morphism $f: X \rightarrow Y$, injectivity of \hat{Y} implies the existence of a morphism $g: \hat{X} \rightarrow \hat{Y}$ extending $y_y \cdot f$ along y_x

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ y_x \downarrow & & \downarrow y_y \\ \hat{X} & \xrightarrow{g} & \hat{Y}, \end{array}$$

hence we set $Gf = |f|$. G is faithful and to see it is full, let $f: |X| \rightarrow |Y|$ be a map commuting the diagram

$$\begin{array}{ccc} |X| & \xrightarrow{f} & |Y| \\ |y_x| \downarrow & & \downarrow |y_y| \\ |\hat{X}| & \xrightarrow{|g|} & |\hat{Y}|, \end{array}$$

for some $g: \hat{X} \rightarrow \hat{Y}$ in \mathbf{X} , then $|-|$ -initiality of y_y implies the existence of a unique $\bar{f}: X \rightarrow Y$ such that $G\bar{f} = |\bar{f}| = f$.

For essential surjectivity, let $(X, A, \sigma: A \rightarrow |X|)$ in $\mathcal{F}(\mathbf{X}_{\text{inj}})$ and take the $|-|$ -initial lifting of σ , that we denote by $\sigma_{\text{ini}}: A_{\text{ini}} \rightarrow X$, so $|A_{\text{ini}}| = A$ and $|\sigma_{\text{ini}}| = \sigma$. Hence $GA_{\text{ini}} = (\hat{A}_{\text{ini}}, A, |y_{A_{\text{ini}}}|: A \rightarrow |\hat{A}_{\text{ini}}|)$ and we verify that the identity map $1_A: A \rightarrow A$ is a morphism from (X, A, σ) to $(\hat{A}_{\text{ini}}, A, |y_{A_{\text{ini}}}|)$,

and vice-versa. The latter fact comes readily from injectivity of \hat{A}_{ini} and $|-|$ -initiality of σ_{ini} :

$$\begin{array}{ccc} A & \xrightarrow{1_A} & A \\ \sigma \downarrow & & \downarrow |y_{A_{\text{ini}}}| \\ |X| & \xrightarrow{|g_1|} & |\hat{A}_{\text{ini}}|, \end{array}$$

for some morphism $g_1 : X \rightarrow \hat{A}_{\text{ini}}$, and by injectivity of X and $|-|$ -initiality of $y_{A_{\text{ini}}}$:

$$\begin{array}{ccc} A & \xrightarrow{1_A} & A \\ |y_{A_{\text{ini}}}| \downarrow & & \downarrow \sigma \\ |\hat{A}_{\text{ini}}| & \xrightarrow{|g_2|} & |X|, \end{array}$$

for some morphism $g_2 : \hat{A}_{\text{ini}} \rightarrow X$. Therefore, $GA_{\text{ini}} \cong (X, A, \sigma : A \rightarrow |X|)$ in $\mathcal{F}(\mathbf{X}_{\text{inj}})$. \blacksquare

Since \mathbf{X}_{inj} is (weakly) cartesian closed, as shown in [CR00], $\mathbf{X} \cong \mathcal{F}(\mathbf{X}_{\text{inj}})$ has all weak simple products (in particular it is weakly cartesian closed), and more, $\mathbf{X} \cong \mathcal{F}(\mathbf{X}_{\text{inj}})$ is weakly locally cartesian closed, i.e. it has weak dependent products, whence by [CR00, Theorem 3.3], $\mathbf{X}_{\text{ex}} \cong \mathcal{F}(\mathbf{X}_{\text{inj}})_{\text{ex}}$ is locally cartesian closed.

Therefore, by Theorem 2.1, we conclude that $\mathbf{X}\text{-Equ}$ is locally cartesian closed (see for instance [HST14, III-Corollary 4.6.2]), and, being complete and cocomplete, one may ask whether this category is actually a quasitopos.

As discussed in [Ros98], “.. *the full subcategory of \mathcal{C}_{ex} consisting of those equivalence spans which are kernel pairs in \mathcal{C} gives the free regular completion \mathcal{C}_{reg} of \mathcal{C} .*”, where in that context *equivalence span* means pseudo-equivalence relation. Hence, similar to what is observed in [Men00] for topological spaces, the category $\mathbf{X}\text{-Equ}$, presented by $\text{PER}(\mathbf{X}, \text{RegMono})$, is equivalent to the regular completion \mathbf{X}_{reg} of \mathbf{X} .

It is easy to depict the latter equivalence using the classical description of \mathbf{X}_{reg} [Car95]: objects are \mathbf{X} -morphisms $f : X \rightarrow Y$, and a morphism from $f : X \rightarrow Y$ to $g : Z \rightarrow W$ is an equivalence class $[l]$ of an \mathbf{X} -morphism

$l: X \rightarrow Z$ such that $g \cdot l \cdot f_0 = g \cdot l \cdot f_1$, where f_0, f_1 form the kernel pair of f .

$$\begin{array}{ccc} \text{Ker}(f) & \xrightarrow{f_1} & X \\ f_0 \downarrow & \lrcorner & \downarrow f \\ X & \xrightarrow{f} & Y \end{array}$$

Two such arrows l and m are equivalent if $g \cdot l = g \cdot m$.

$$\begin{array}{ccc} X & & Z \\ f \downarrow & \xrightarrow{[l]} & \downarrow g \\ Y & & W \end{array}$$

Lemma 2.2. \mathbf{X}_{reg} and $\text{PER}(\mathbf{X}, \text{RegMono})$ are equivalent.

Proof. Define $F: \mathbf{X}_{\text{reg}} \rightarrow \text{PER}(\mathbf{X}, \text{RegMono})$ as in the diagram below,

$$\begin{array}{ccc} (f: X \rightarrow Y) & \longmapsto & (\text{Ker}(f), X, f_0, f_1) \\ [l] \downarrow & & \downarrow [l] \\ (g: Z \rightarrow W) & \longmapsto & (\text{Ker}(g), Z, g_0, g_1) \end{array}$$

so it is a well-defined functor, since $l: X \rightarrow Z$ satisfies $g \cdot l \cdot f_0 = g \cdot l \cdot f_1$ if, and only if, there exists a unique $\bar{l}: \text{Ker}(f) \rightarrow \text{Ker}(g)$ such that $g_0 \cdot \bar{l} = l \cdot f_0$ and $g_1 \cdot \bar{l} = l \cdot f_1$.

$$\begin{array}{ccccc} \text{Ker}(f) & & \xrightarrow{l \cdot f_1} & & Z \\ & \searrow \bar{l} & & \nearrow g_1 & \\ & & \text{Ker}(g) & \xrightarrow{g_1} & Z \\ & & g_0 \downarrow & \lrcorner & \downarrow g \\ & \searrow l \cdot f_0 & & \nearrow g & \\ & & Z & \xrightarrow{g} & W \end{array}$$

Then F is fully faithful, and it is essentially surjective because each pseudo-equivalence relation $X_1 \xrightarrow[r_2]{r_1} X_0$ with $\langle r_1, r_2 \rangle: X_1 \rightarrow X_0 \times X_0$ a regular monomorphism is seen to form the kernel pair of the $|-|$ -final lifting $\bar{p}: X_0 \rightarrow$

\tilde{X}_0 of the projection map $p: |X_0| \rightarrow |X_0|/\sim$, where \sim is the equivalence relation on $|X_0|$ defined in the proof of Lemma 2.1. \blacksquare

We now intend to apply [Men00, Corollary 8.4.2]; by condition **(f)** and Proposition 2.2, \mathbf{X} is an (infinitely) extensive weakly locally cartesian closed category, hence we are only missing the *chaotic situation* described right after [Men00, Lemma 7.3.3]. This comes from the observation that the topos \mathbf{Set} is a mono-localization of \mathbf{X} , since the topological functor $|-|: \mathbf{X} \rightarrow \mathbf{Set}$ is faithful, preserves finite limits and has a full embedding as a right adjoint [AHS90, Proposition 21.12]. Therefore, by Lemma 2.1, Lemma 2.2 and [Men00, Corollary 8.4.2], we conclude:

Theorem 2.3. *\mathbf{X} -Equ is a quasitopos.*

3. The case $\mathbf{X}=(\mathbb{T}, \mathbf{V})$ -Cat

We briefly introduce the (\mathbb{T}, \mathbf{V}) -Cat setting, and refer the reader to the reference [CT03] for details (see also [HST14]).

Although introduced in a more general setting, we are interested here in the case when

- $\mathbf{V} = (\mathbf{V}, \otimes, k)$ is a commutative unital quantale (see for instance [HST14, II-Section 1.10]) which is also a *Heyting algebra* (so that the operation infimum \wedge also has a right adjoint), and
- $\mathbb{T} = (T, m, e): \mathbf{Set} \rightarrow \mathbf{Set}$ is a monad satisfying the Beck-Chevalley condition (T preserves weak pullbacks and the naturality squares of m are weak pullbacks [CHJ14]) that is laxly extended to the pre-ordered category $\mathbf{V}\text{-Rel}$, which has as objects sets and as morphisms \mathbf{V} -relations $r: X \dashrightarrow Y$, i.e. \mathbf{V} -valued maps $r: X \times Y \rightarrow \mathbf{V}$.

Hence we assume that there exists a functor $T: \mathbf{V}\text{-Rel} \rightarrow \mathbf{V}\text{-Rel}$ extending T , by abuse of notation denoted by the same letter, that commutes with involution: $T(r^\circ) = (Tr)^\circ$, for each $r: X \dashrightarrow Y \in \mathbf{V}\text{-Rel}$, where $r^\circ(y, x) = r(x, y)$, for each $(x, y) \in X \times Y$. The functor T turns m and e into oplax transformations, meaning that the naturality diagrams become:

$$\begin{array}{ccccc} X & \xrightarrow{e_X} & TX & \xleftarrow{m_X} & T^2X \\ r \downarrow & \leq & \downarrow Tr & \geq & \downarrow T^2r \\ Y & \xrightarrow{e_Y} & TY & \xleftarrow{m_Y} & T^2Y, \end{array}$$

for each \mathbf{V} -relation $r: X \dashrightarrow Y$.

Hence we have a lax monad on $\mathbf{V}\text{-Rel}$ [CH04] and $(\mathbb{T}, \mathbf{V})\text{-Cat}$ is defined as the category of Eilenberg-Moore lax algebras for that lax monad: objects are pairs (X, a) , where X is a set and $a: TX \dashrightarrow X$ is a \mathbf{V} -relation, which is reflexive and transitive.

$$\begin{array}{ccccc} X & \xrightarrow{e_X} & TX & \xleftarrow{Ta} & T^2X \\ & \searrow & \downarrow a & \leq & \downarrow m_X \\ & & X & \xleftarrow{a} & TX \\ & \nearrow 1_X & & & \end{array}$$

Such pairs are called $(\mathbb{T}, \mathbf{V})\text{-categories}$; a morphism from (X, a) to (Y, b) is a map $f: X \rightarrow Y$ commuting the diagram below.

$$\begin{array}{ccc} TX & \xrightarrow{Tf} & TY \\ a \downarrow & \leq & \downarrow b \\ X & \xrightarrow{f} & Y \end{array}$$

Such a map is called $(\mathbb{T}, \mathbf{V})\text{-functor}$.

We are also going to restrict ourselves to the case that the extension T to $\mathbf{V}\text{-Rel}$ is determined by a \mathbb{T} -algebra structure map $\xi: TV \rightarrow V$, so we are in the setting of *topological theories* [Hof07] (see also [CT14]), hence \mathbf{V} has a $(\mathbb{T}, \mathbf{V})\text{-category}$ structure given by the composite

$$TV \xrightarrow{\xi} V \xrightarrow{\text{hom}} V,$$

where $\text{hom}: V \times V \rightarrow V$ is the left adjoint of \otimes , so

$$u \otimes v \leq w \iff u \leq \text{hom}(v, w),$$

for each u, v, w in the quantale \mathbf{V} .

The forgetful functor $|-|: (\mathbb{T}, \mathbf{V})\text{-Cat} \rightarrow \mathbf{Set}$ is topological [CH03, CT03], and before we provide examples of categories given by $(\mathbb{T}, \mathbf{V})\text{-Cat}$, we verify that, for suitable monad \mathbb{T} and quantale \mathbf{V} satisfying the conditions assumed so far in this section, $(\mathbb{T}, \mathbf{V})\text{-Cat}$ satisfies all conditions **(a)** to **(f)** from the two previous sections. In each item, we highlight the properties that are needed in order to achieve the respective condition, adding the references where that is proved.

(a) $(\mathbb{T}, \mathbf{V})\text{-Cat}$ is pre-ordered enriched. For $(\mathbb{T}, \mathbf{V})\text{-categories}$ (X, a) and (Y, b) , consider the following relation on the set of $(\mathbb{T}, \mathbf{V})\text{-functors}$ from (X, a)

to (Y, b) :

$$f \leq g \iff \forall x \in X, k \leq b(e_Y(f(x)), g(x)).$$

This determines a pre-order, first defined in [CT03], which is compatible with composition of (\mathbb{T}, \mathbf{V}) -functors. One can also check that a (\mathbb{T}, \mathbf{V}) -category (X, a) is separated if, and only if, the following pre-order on X is anti-symmetric:

$$x \leq x' \iff k \leq a(e_X(x), x')$$

(see [HST14, III-Proposition 3.3.1]).

(b) *|-|-initial (\mathbb{T}, \mathbf{V}) -functors reflect the order.* Let $(X, a), (Y, b)$ be (\mathbb{T}, \mathbf{V}) -categories, $x, x': \mathbf{1} \rightarrow (X, a)$ be (\mathbb{T}, \mathbf{V}) -functors, where $\mathbf{1} = (\{*\}, e_{\{*\}}^\circ)$ (the discrete structure on the singleton [HST14, III-Section 3.2]), and $f: (X, a) \rightarrow (Y, b)$ an $|-|-$ initial (\mathbb{T}, \mathbf{V}) -functor such that $f \cdot x \simeq f \cdot x'$; $|-|-$ initiality of f means that $a(\mathfrak{x}, x) = b(Tf(\mathfrak{x}), f(x))$, for each $\mathfrak{x} \in TX, x \in X$. We calculate:

$$\begin{aligned} k &\leq b(e_Y(f \cdot x(*)), f \cdot x'(*)) && \text{(definition of } f \cdot x \leq f \cdot x') \\ &\leq b(Tf \cdot e_X(x(*)), f \cdot x'(*)) && \text{(composition is associative, } e \text{ is natural)} \\ &\leq a(e_X(x(*)), x'(*)) && \text{(} f \text{ is } |-|- \text{initial),} \end{aligned}$$

so $x \leq x'$ and in the same fashion $x' \leq x$, thus $x \simeq x'$.

(c) *(\mathbb{T}, \mathbf{V}) -Cat has enough injectives.* The tensor product \otimes of \mathbf{V} induces a functor $\otimes: (\mathbb{T}, \mathbf{V})\text{-Cat} \times (\mathbb{T}, \mathbf{V})\text{-Cat} \rightarrow (\mathbb{T}, \mathbf{V})\text{-Cat}$, with

$$(X, a) \otimes (Y, b) = (X \times Y, c),$$

where, for each $\mathfrak{w} \in T(X \times Y), (x, y) \in X \times Y$,

$$c(\mathfrak{w}, (x, y)) = a(T\pi_X(\mathfrak{w}), x) \otimes b(T\pi_Y(\mathfrak{w}), y),$$

and π_X, π_Y are the projections from $X \times Y$ onto X and Y , respectively. The following facts can be found in [CH09, Hof11, CCH15]: for each (\mathbb{T}, \mathbf{V}) -category (X, a) , $a: TX \dashrightarrow X$ defines a (\mathbb{T}, \mathbf{V}) -functor

$$a: X^{\text{op}} \otimes X \rightarrow \mathbf{V},$$

where $X^{\text{op}} = (TX, m_X \cdot (Ta)^\circ \cdot m_X)$; the \otimes -exponential mate $y_X: X \rightarrow \mathbf{V}^{X^{\text{op}}}$ of a is fully faithful; the (\mathbb{T}, \mathbf{V}) -category $PX = \mathbf{V}^{X^{\text{op}}}$ is injective and if (X, a) is separated, so is PX .

(d) *Injectives are exponentiable.* Conditions under which injectivity implies exponentiability in $(\mathbb{T}, \mathbf{V})\text{-Cat}$ are studied in [CHR18]. We recall them next. Consider the maps

$$\mathbf{V} \otimes \mathbf{V} \xrightarrow{\otimes} \mathbf{V} \quad \text{and} \quad X \xrightarrow{(-, u)} X \otimes \mathbf{V}, \quad (1)$$

with $(\mathbf{V}, \text{hom}_\varepsilon), (X, a) \in (\mathbb{T}, \mathbf{V})\text{-Cat}$. Define also for a \mathbf{V} -relation $r: X \dashrightarrow Y$ and $u \in \mathbf{V}$, the \mathbf{V} -relation $r \otimes u: X \dashrightarrow Y$ given by

$$(r \otimes u)(x, y) = r(x, y) \otimes u, \quad (2)$$

for each $(x, y) \in X \times Y$. As a final condition, assume that, for all $u, v, w \in \mathbf{V}$,

$$w \wedge (u \otimes v) = \{u' \otimes v' \mid u' \leq u, v' \leq v, u' \otimes v' \leq w\}, \quad (3)$$

which is equivalent to exponentiability of injective \mathbf{V} -categories (see [HR13, Theorem 5.3]). Then [CHR18, Theorem 5.4] says the following:

Theorem 3.1. *Suppose that: the maps \otimes and $(-, u)$ in (1) are (\mathbb{T}, \mathbf{V}) -functors; for every injective (\mathbb{T}, \mathbf{V}) -category (X, a) and every $u \in \mathbf{V}$, $T(a \otimes u) = Ta \otimes u$, with those \mathbf{V} -relations defined as in (2); and (3) holds. Then every injective (\mathbb{T}, \mathbf{V}) -category is exponentiable in $(\mathbb{T}, \mathbf{V})\text{-Cat}$.*

(e) *The reflector from $(\mathbb{T}, \mathbf{V})\text{-Cat}$ to $(\mathbb{T}, \mathbf{V})\text{-Cat}_{\text{sep}}$ preserves finite products.* This is proved in [CHR18, Proposition 5.4].

(f) *$(\mathbb{T}, \mathbf{V})\text{-Cat}$ is infinitely extensive.* This is proved in [MST06] under the condition that T is a taut functor [Man02], what comes for free from the assumption that T preserves weak pullbacks.

To give examples of categories satisfying all the conditions above, we consider:

- the identity monad $\mathbb{I} = (\text{Id}, 1, 1)$ on Set laxly extended to the identity lax monad on $\mathbf{V}\text{-Rel}$;
- the ultrafilter monad \mathbb{U} with the Barr extension to $\mathbf{V}\text{-Rel}$ [HST14, IV-Corollary 2.4.5], with \mathbf{V} integral and *completely distributive* (see, for instance, [HST14, II-Section 1.11]);
- the list monad (or free monoid monad) $\mathbb{L} = (L, m, e)$ (see [HST14, II-Examples 3.1.1(2)]), with the extension $L: \mathbf{V}\text{-Rel} \rightarrow \mathbf{V}\text{-Rel}$ that sends each

$r: X \dashrightarrow Y$ to the \mathbf{V} -relation $Lr: LX \dashrightarrow LY$ given by

$$Lr((x_1, \dots, x_n), (y_1, \dots, y_m)) = \begin{cases} r(x_1, y_1) \otimes \dots \otimes r(x_n, y_n), & \text{if } n = m \\ \perp, & \text{if } n \neq m; \end{cases}$$

• the monad $\mathbb{M} = (- \times M, m, e)$, for a monoid $(M, \cdot, 1_M)$, with $m_x: X \times M \times M \rightarrow X \times M$ given by $m_x(x, a, b) = (x, a \cdot b)$ and $e_x: X \rightarrow X \times M$ given by $e_x(x) = (x, 1_M)$ (see [HST14, V-Section 1.4]). The extension $- \times M: \mathbf{V}\text{-Rel} \rightarrow \mathbf{V}\text{-Rel}$ sends the \mathbf{V} -relation $r: X \dashrightarrow Y$ to the \mathbf{V} -relation $r \times M: X \times M \dashrightarrow Y \times M$ with

$$r \times M((x, a), (y, b)) = \begin{cases} r(x, y), & \text{if } a = b, \\ \perp, & \text{if } a \neq b. \end{cases}$$

As well as the quantales: $\mathbf{2} = (\{\perp, \top\}, \wedge, \top)$, $\mathbf{P}_+ = ([0, \infty]^{\text{op}}, +, 0)$, $\mathbf{P}_{\max} = ([0, \infty]^{\text{op}}, \max, 0)$, $\mathbf{2}^2 = (\{\perp, u, v, \top\}, \wedge, \top)$ (the *diamond lattice* [HST14, II-Exercise 1.H]) and Δ (the quantale of distribution functions [HR13]). We assemble the table:

$\mathbf{V} \backslash \mathbb{T}$	\mathbb{I}	\mathbb{U}	\mathbb{L}	\mathbb{M}
$\mathbf{2}$	Ord	Top	MultiOrd	$(\mathbb{M}, \mathbf{2})\text{-Cat}$
\mathbf{P}_+	Met	App		
\mathbf{P}_{\max}	UltMet	NA-App		
$\mathbf{2}^2$	BiRel	BiTop		
Δ	ProbMet			

(4)

- **Ord** is the category of pre-ordered spaces,
- **Met** is the category of Lawvere generalized metric spaces [Law02],
- **UltMet** is the full subcategory of **Met** of ultra-metric spaces [HST14, III-Exercise 2.B],
- **BiRel** is the one of sets and birelations [HST14, III-Examples 1.1.1(3)],
- **ProbMet** is the category of probabilistic metric spaces [HR13],
- **Top** is the usual category of topological spaces and continuous functions,
- **App** is that of Lowen's approach spaces [Low97], and
- **NA-App** is the full subcategory of **App** of non-Archimedean approach spaces studied in details in [CVO17], and denoted in [Hof14] by **UApp**,
- **BiTop** is the category of bitopological spaces and bicontinuous maps [HST14, III-Exercise 2.D],

- **MultiOrd** is the category of multi-ordered sets [HST14, V-Section 1.4], and
- **(M, 2)-Cat** can be interpreted as the category of M -labelled ordered sets [HST14, V-Section 1.4].

For instance, an object of **Ord-Equ** is a pre-ordered set (X, \leq) together with an equivalence relation \equiv_X on X ; separatedness of (X, \leq) means that \leq is anti-symmetric, so the objects of **Ord-Equ_{sep}** are partially ordered sets equipped with equivalence relations on their underlying sets. Further, a partial equilogical separated object in **Ord-PEqu_{sep}** is a complete lattice (injective ordered set) together with an equivalence relation on the underlying set. In the same fashion, the objects of the category **Mdst(Ord_{sep, inj})** are triples $(A, (X, \leq), E_A)$, with A a set, $E_A: A \rightarrow \mathcal{P}X$ a function, and (X, \leq) a complete lattice.

Furthermore, from Section 2 we conclude that, together with **Top**, all the other categories in Table 4 are weakly locally cartesian closed and their exact completions are locally cartesian closed categories; moreover, their categories of equilogical objects, which are equivalent to their regular completions, are quasitoposes that fully embed the original categories.

Concerning four of those categories, we also have adjunctions

$$\begin{array}{ccc}
 \mathbf{Top} & \overset{\dashleftarrow}{\perp} & \mathbf{App} \\
 \uparrow & & \uparrow \\
 \mathbf{Ord} & \overset{\dashleftarrow}{\perp} & \mathbf{Met},
 \end{array}$$

(Note: The diagram shows solid arrows for the adjunctions $\mathbf{Top} \dashv \mathbf{App}$ and $\mathbf{Ord} \dashv \mathbf{Met}$, and dotted arrows for the embeddings $\mathbf{Ord} \hookrightarrow \mathbf{Top}$ and $\mathbf{Met} \hookrightarrow \mathbf{App}$. The hook-arrows are full embeddings.)

where both solid and dotted diagrams commute, the hook-arrows are full embeddings and the two full embeddings $\mathbf{Ord} \hookrightarrow \mathbf{App}$ coincide (see [HST14, III-Section 3.6]). One can see that those adjunctions extend to the respective categories of equilogical objects,

$$\begin{array}{ccc}
 \mathbf{Equ} & \overset{\dashleftarrow}{\perp} & \mathbf{App-Equ} \\
 \uparrow & & \uparrow \\
 \mathbf{Ord-Equ} & \overset{\dashleftarrow}{\perp} & \mathbf{Met-Equ}
 \end{array}$$

(Note: The diagram shows solid arrows for the adjunctions $\mathbf{Equ} \dashv \mathbf{App-Equ}$ and $\mathbf{Ord-Equ} \dashv \mathbf{Met-Equ}$, and dotted arrows for the embeddings $\mathbf{Ord-Equ} \hookrightarrow \mathbf{Equ}$ and $\mathbf{Met-Equ} \hookrightarrow \mathbf{App-Equ}$. The hook-arrows are full embeddings.)

and we describe them now.

(1) **Ord-Equ to Met-Equ.** Each ordered equiological object $\langle (X, \leq), \equiv_x \rangle$ is taken to $\langle (X, d_\leq), \equiv_x \rangle$, where the metric d_\leq is given by

$$d_\leq(x, x') = \begin{cases} 0, & \text{if } x \leq x' \\ \infty, & \text{otherwise,} \end{cases}$$

for each $x, x' \in X$. The left adjoint of this inclusion assigns $\langle (X, \leq_d), \equiv_x \rangle$ to $\langle (X, d), \equiv_x \rangle$, with $x \leq_d x'$ if and only if $d(x, x') < \infty$, for each $x, x' \in X$. Hence the category **Ord-Equ** is fully embedded in **Met-Equ** as the metric equiological objects $\langle (X, d), \equiv_x \rangle$ for which there exists an order \leq on X such that $d = d_\leq$.

(2) **Ord-Equ to Equ.** Each $\langle (X, \leq), \equiv_x \rangle$ is taken to $\langle (X, \tau_\leq), \equiv_x \rangle$, where τ_\leq is the Alexandroff topology: open sets are generated by the down-sets $\downarrow x$, $x \in X$. For its right adjoint, to an equiological space $\langle (X, \tau), \equiv_x \rangle$ is assigned $\langle (X, \leq_\tau), \equiv_x \rangle$, where \leq_τ is the specialization order of (X, τ) : for each $x, x' \in X$, $x \leq_\tau x'$ if and only if $\dot{x} \rightarrow x'$, where \dot{x} denotes the principal ultrafilter on x , and \rightarrow denotes the convergence relation between ultrafilters and points of X determined by τ ; observe that this is the induced order described in item **(a)** above. Hence the category **Ord-Equ** is fully embedded in **Equ** as the equiological spaces $\langle (X, \tau), \equiv_x \rangle$ for which there exists an order \leq on X such that $\tau = \tau_\leq$, and those are exactly the Alexandroff spaces: arbitrary intersections of open sets are open (see [HST14, II-Example 5.10.5, III-Example 3.4.3(1)]).

(3) **Met-Equ to App-Equ.** A metric equiological object $\langle (X, d), \equiv_x \rangle$ becomes an approach equiological one $\langle (X, \delta_d), \equiv_x \rangle$, where the approach distance is given by $\delta_d(x', A) = \inf\{d(x, x') \mid x \in A\}$, for each $x' \in X$, $A \in \mathcal{P}X$ [HST14, III-Examples 2.4.1(1)]. The right adjoint of this embedding assigns $\langle (X, d_\delta), \equiv_x \rangle$ to $\langle (X, \delta), \equiv_x \rangle$, where $d_\delta(x, x') = \sup\{\delta(x', A) \mid x \in A \in \mathcal{P}X\}$, for each $x, x' \in X$. Hence **Met-Equ** is identified within **App-Equ** as the approach equiological objects $\langle (X, \delta), \equiv_x \rangle$ such that $\delta = \delta_d$, for some metric d on X , that is, (X, δ) is a metric approach space [Low97, Chapter 3].

(4) **Equ to App-Equ.** An equiological space $\langle (X, \tau), \equiv_x \rangle$ becomes an approach equiological one $\langle (X, \delta_\tau), \equiv_x \rangle$, where the approach distance is given by

$$\delta_\tau(x', A) = \begin{cases} 0, & \text{if } A \in \mathfrak{x}, \text{ for some } \mathfrak{x} \in UX \text{ with } \mathfrak{x} \rightarrow x' \\ \infty, & \text{otherwise,} \end{cases}$$

for each $x' \in X$, $A \in \mathcal{P}X$, where UX denotes the set of ultrafilters on X [HST14, III-Examples 2.4.1(2)]. The left adjoint of this embedding is slightly more elaborate: for an approach equilogical object $\langle (X, \delta), \equiv_x \rangle$, consider the convergence relation \rightarrow between ultrafilters in UX and points of X given by

$$\mathfrak{x} \rightarrow x \iff \sup\{\delta(x, A) \mid A \in \mathfrak{x}\} < \infty;$$

this convergence defines a pseudo-topological space [Cho48], to which we apply the reflector described in [HST14, III-Exercise 3.D], obtaining an equilogical space $\langle (X, \tau_\delta), \equiv_x \rangle$. Hence **Equ** is identified within **App-Equ** as the approach equilogical objects $\langle (X, \delta), \equiv_x \rangle$ such that $\delta = \delta_\tau$, for some topology τ on X , that is, (X, δ) is a topological approach space [Low97, Chapter 2].

Open question. The conditions **(a)** to **(f)** of Sections 1 and 2 were derived from the successful attempt of generalizing the structures/constructions to $(\mathbb{T}, \mathbb{V})\text{-Cat}$, for suitable \mathbb{T} and \mathbb{V} . Requiring those conditions on an arbitrary category with a topological functor over **Set** produced the same desired results. However, we do not know an example of a category satisfying those conditions which cannot be described as $(\mathbb{T}, \mathbb{V})\text{-Cat}$.

Acknowledgments

This work was done during the preparation of the author's PhD thesis, under the supervision of Maria Manuel Clementino, whom the author thanks for proposing the investigation and advising the whole study. I also thank Fernando Lucatelli Nunes for fruitful discussions.

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WILLIAN RIBEIRO

CMUC, DEPARTMENT OF MATHEMATICS, UNIVERSITY OF COIMBRA, 3001-501 COIMBRA, PORTUGAL

E-mail address: willian.ribeiro.vs@gmail.com