# EFFECTIVE DESCENT MORPHISMS OF FILTERED PREORDERS

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ABSTRACT. We characterize effective descent morphisms of what we call filtered preorders, and apply these results to slightly improve a known result, due to the first author and F. Lucatelli Nunes, on the effective descent morphisms in lax comma categories of preorders. A filtered preorder, over a fixed preorder X, is defined as a preorder A equipped with a profunctor  $X \to A$  and, equivalently, as a set A equipped with a family  $(A_x)_{x \in X}$  of upclosed subsets of A with  $x' \leq x \Rightarrow A_x \subseteq A_{x'}$ .

### 1. INTRODUCTION

Throughout this paper, by a *preorder* we mean a set equipped with a reflexive transitive relation, and we write  $X = (X, \leq)$  for a fixed preorder. Such preorders were simply called ordered sets, e.g. in [1], and keeping the notation of [1], we will write:

- Ord for the category of preorders;
- $\operatorname{Ord}//X$  for the lax comma category of  $\operatorname{Ord}$  over X, and we recall that its objects are the same as in  $\operatorname{Ord}/X = (\operatorname{Ord} \downarrow X)$ , while a morphism

$$f: (A, \alpha) \to (B, \beta)$$

in  $\operatorname{Ord}//X$  is a morphism  $f: A \to B$  in  $\operatorname{Ord}$  with  $\alpha \leq \beta f$ , that is, with

$$\alpha(a)\leqslant\beta f(a)$$

for all  $a \in A$ .

Considering one of the problems formulated in [1], namely the problem of characterizing effective descent morphisms in  $\operatorname{Ord}//X$ , we found a 'nicer' category  $\operatorname{Ord}_X$ , of what we called 'filtered preorders', where this problem turned out to be easy enough to solve it completely. It also turned out that dealing with  $\operatorname{Ord}_X$  helps to improve what was done in [1] with  $\operatorname{Ord}//X$ . Here, a filtered preorder, over a fixed preorder X, is defined as a preorder A equipped with a profunctor  $X \to A$  and, equivalently, as a set A equipped with a family  $(A_x)_{x \in X}$  of upclosed subsets of A with  $x' \leq x \Rightarrow A_x \subseteq A_{x'}$ .

The paper is organized as follows:

Section 2 is devoted to general categorical remarks, notably involving the notion of map of adjunctions which was briefly mentioned in [6]. These remarks should be reformulated in the style of [7] one day, but we did not go further than what we needed for our purposes. We should particularly mention the contrast between 'strict' and 'pseudo-', which appears as the contrast between Corollary 9.6 in [5] and our Corollary 2.6(b). We are satisfied with 'strict' since we are using it for functors that preserve relevant chosen limits and colimits (specifically, pullbacks and coequalizers).

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Section 3 introduces filtered preorders and shows how their category extends the category Ord//X, while the purpose of Section 4 is to characterize effective descent morphisms in  $Ord_X$ , and Section 5 gives a new class of effective descent morphisms in Ord//X slightly improving Theorem 3.9 of [1].

## 2. General remarks on monadicity and descent

In this section we first deal with a fixed map  $(P, P'): \Phi \to X$  of adjunctions in the sense of [6] (see Section 7 of Chapter IV therein) displayed as

$$\begin{array}{cccc}
\mathcal{U} & \stackrel{\Phi}{\longrightarrow} & \mathcal{U}' \\
 P & & & & & \\
 P & & & & & \\
 W & \xrightarrow{} & \mathcal{W}' \\
 & & & \mathcal{W}'
\end{array}$$

In this diagram P and P' are functors while

$$\Phi = (\Phi_!, \Phi^*, \eta^{\Phi}, \varepsilon^{\Phi}) \colon \mathcal{U} \to \mathcal{U}' \text{ and } X = (X_!, X^*, \eta^X, \varepsilon^X) \colon \mathcal{W} \to \mathcal{W}'$$

are adjunctions with

$$X_!P = P'\Phi_!, \ P\Phi^* = X^*P', \ P\eta^{\Phi} = \eta^X P, \ P'\varepsilon^{\Phi} = \varepsilon^X P'.$$

Equivalently, we can present it as an adjunction

$$((\Phi_!, X_!), (\Phi^*, X^*), (\eta^{\Phi}, \eta^X), (\varepsilon^{\Phi}, \varepsilon^X)) \colon (\mathfrak{U}, P, \mathfrak{W}) \to (\mathfrak{U}', P', \mathfrak{W}'),$$

in the 2-category ArrCAT (=  $Cat^2$  in the notation of, e.g., [4]) of arrows of the category of categories. In order to make this clear, let us just recall that:

- the objects of ArrCAT are all functors  $F: \mathcal{A} \to \mathcal{B}$ , written as triples  $(\mathcal{A}, F, \mathcal{B})$ ;
- a morphism  $(\mathcal{A}, F, \mathcal{B}) \to (\mathcal{A}', F', \mathcal{B}')$  is a pair of functors (K, L) making the diagram

$$\begin{array}{c|c} \mathcal{A} & \xrightarrow{K} \mathcal{A}' \\ F & & \downarrow F' \\ \mathcal{B} & \xrightarrow{L} \mathcal{B}' \end{array}$$

commute;

• for morphisms  $(K, L), (M, N) : (\mathcal{A}, F, \mathcal{B}) \to (\mathcal{A}', F', \mathcal{B}')$ , a 2-cell from (K, L) to (M, N) is a pair of natural transformations  $(\sigma, \tau) =$ 

$$\begin{array}{c|c} \mathcal{A} & \xrightarrow{K} & \mathcal{A}' \\ \mathcal{A} & \xrightarrow{\downarrow\sigma} & \mathcal{A}' \\ \mathcal{F} & & & \downarrow \\ \mathcal{F} & & & \downarrow \\ \mathcal{B} & \xrightarrow{\downarrow\tau} & \mathcal{B}' \\ \mathcal{B} & \xrightarrow{N} & \mathcal{B}' \end{array}$$

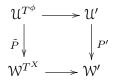
with  $F'\sigma = \tau F$ .

Our first remark is: the equations  $P\eta^{\Phi} = \eta^X P$  and  $P' \varepsilon^{\Phi} = \varepsilon^X P'$  immediately imply:

# Lemma 2.1. For the map of adjunctions above:

- (1) if  $\eta^{\Phi}$  is an isomorphism and P is surjective on objects, then  $\eta^X$  is an isomorphism;
- (2) if  $\varepsilon^{\Phi}$  is an isomorphism and P' is surjective on objects, then  $\varepsilon^X$  is an isomorphism;
- (3) in particular, if  $\Phi$  is a category equivalence and both P and P' are surjective on objects, then X is a category equivalence.

Next, our map of adjunctions above determines, whenever the categories  $\mathcal{U}'$  and  $\mathcal{W}'$  have chosen coequalizers preserved by the functor P', its *derived map of adjunctions* 



where:  $T^{\phi}$  and  $T^X$  are the monads determined by the adjunctions  $\Phi$  and X, respectively;  $\tilde{P}$  is induced by (P, P'); and the horizontal arrows are the comparison adjunctions. If (P, P') is a split epimorphism of adjunctions, then  $\tilde{P}$  also is a split epimorphism and, using Lemma 2.1, one can easily show that monadicity of  $\Phi$  implies monadicity of X. The precise statement is:

**Lemma 2.2.** Let  $(P, P'): \Phi \to X$  and  $(I, I'): X \to \Phi$  be maps of adjunctions with  $PI = 1_W$ and  $P'I' = 1_{W'}$ . If the categories  $\mathcal{U}'$  and  $\mathcal{W}'$  have chosen coequalizers preserved by the functors P' and I', then premonadicity of  $\Phi$  implies premonadicity of X and monadicity of  $\Phi$  implies monadicity of X.

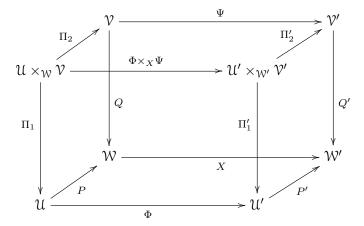
Let  $\mathcal{C}_0$  and  $\mathcal{C}_1$  be categories with chosen pullbacks, and  $S_1: \mathcal{C}_1 \to \mathcal{C}_0$  a functor that preserves them. Then any morphism  $p_1: E_1 \to B_1$  in  $\mathcal{C}_1$  determines a map

$$\begin{array}{c} (\mathfrak{C}_{1} \downarrow E_{1}) \xrightarrow{(p_{1!}, p_{1}^{*}, \eta^{p_{1}}, \varepsilon^{p_{1}})} \\ (\mathfrak{C}_{1} \downarrow E_{1}) \xrightarrow{(p_{0!}, p_{1}^{*}, \eta^{p_{1}}, \varepsilon^{p_{1}})} \\ (\mathfrak{C}_{0} \downarrow E_{0}) \xrightarrow{(p_{0!}, p_{0}^{*}, \eta^{p_{0}}, \varepsilon^{p_{0}})} \\ \end{array}$$

of adjunctions, in which  $(p_0: E_0 \to B_0) = S_1(p_1: E_1 \to B_1)$ , P and P' are induced by  $S_1$ , and the horizontal arrows are the suitable change-of-base adjunctions. From Lemma 2.2, having in mind the monadic approach to descent theory (see e.g. [3]) we obtain:

**Corollary 2.3.** Let  $C_0$  and  $C_1$  be categories with chosen pullbacks and chosen coequalizers, and  $S_1: C_1 \to C_0$  and  $J_1: C_0 \to C_1$  functors that preserves them and have  $S_1J_1 = 1_{C_0}$ . Then the functor  $S_1$  sends descent morphisms in  $C_1$  to descent morphisms in  $C_0$  and effective descent morphisms in  $C_1$  to effective descent morphisms in  $C_0$ .

Now consider a cube diagram



whose left-hand and the right-hand faces are pullbacks in the category of categories, while all other faces are maps of adjunctions. Here we assume that  $\mathcal{U} \times_{\mathcal{W}} \mathcal{V}$  and  $\mathcal{U}' \times_{\mathcal{W}'} \mathcal{V}'$  are constructed 'as usually', that is, as suitable categories of pairs, and the adjunction  $\Phi \times_X \Psi$  is defined by

$$(\Phi \times_X \Psi)_!(U,V) = (\Phi_!(U), \Psi_!(V)), \ (\Phi \times_X \Psi)^*(U',V') = (\Phi^*(U'), \Psi^*(V')),$$

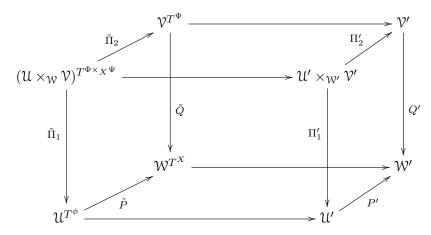
$$\eta_{(U,V)}^{\Phi \times_X \Psi} = (\eta_U^{\Phi}, \eta_V^{\Psi}), \ \varepsilon_{(U',V')}^{\Phi \times_X \Psi} = (\varepsilon_{U'}^{\Phi}, \varepsilon_{V'}^{\Psi}).$$

From the construction of  $\eta_{(U,V)}^{\Phi \times_X \Psi}$  and  $\varepsilon_{(U',V')}^{\Phi \times_X \Psi}$ , we obtain:

Lemma 2.4. Under the assumptions above:

- (1) if  $\eta^{\Phi}$  and  $\eta^{\Psi}$  are isomorphisms, then so is  $\eta^{\Phi \times_X \Psi}$ ;
- (2) if  $\varepsilon^{\Phi}$  and  $\varepsilon^{\Psi}$  are isomorphisms, then so is  $\varepsilon^{\Phi \times_X \Psi}$ ;
- (3) if  $\Phi$  and  $\Psi$  are category equivalences, then so is  $\Phi \times_X \Psi$ .

If the categories  $\mathcal{U}', \mathcal{V}'$ , and  $\mathcal{W}'$  have chosen coequalizers preserved by the functors P' and Q', then we have the *derived cube diagram* 



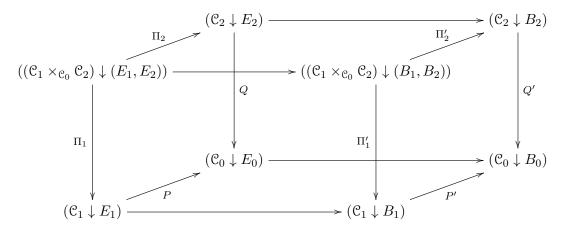
obtained in an obvious way using derived maps of adjunctions. Note, in particular, that  $(\mathcal{U} \times_{\mathcal{W}} \mathcal{V})^{T^{\Phi \times_X \Psi}}$  can be identified with  $\mathcal{U}^{T^{\phi}} \times_{\mathcal{W}^{TX}} \mathcal{V}^{T^{\Psi}}$ . And, from Lemma 2.4, we obtain

**Lemma 2.5.** For maps  $(P, P'): \Phi \to X$  and  $(Q, Q'): \Psi \to X$  of adjunctions, premonadicity of  $\Phi$  and  $\Psi$  implies premonadicity of  $\Phi \times_X \Psi$  and monadicity of  $\Phi$  and  $\Psi$  implies premonadicity of  $\Phi \times_X \Psi$ , provided the categories  $\mathfrak{U}', \mathfrak{V}'$ , and  $\mathfrak{W}'$  have chosen coequalizers preserved by the functors P' and Q'.

Let

$$\mathcal{C}_1 \xrightarrow{S_1} \mathcal{C}_0 \xleftarrow{S_2} \mathcal{C}_2$$

be a cospan of categories having chosen pullbacks preserved by the functors  $S_1$  and  $S_2$ . Given a morphism  $(p_1, p_2): (E_1, E_2) \to (B_1, B_2)$  in  $\mathcal{C}_1 \times_{\mathcal{C}_0} \mathcal{C}_2$ , we can take the originally considered cube diagram to be



where (P, P') is induced by  $S_1$ , (Q, Q') is induced by  $S_2$ , and horizontal arrows are the suitable change-of-base adjunctions. From Lemma 2.5, since the change-of-base adjunction along a morphism is premonadic/monadic if and only if it is a descent/effective descent morphism (see e.g. [3]), we obtain:

**Corollary 2.6.** If the categories  $C_i$  (i = 0, 1, 2) have chosen pullbacks and chosen coequalizers preserved by the functors  $S_1$  and  $S_2$ , then:

- (1) if  $p_1$  and  $p_2$  are descent morphisms, then so is  $(p_1, p_2)$ ;
- (2) if  $p_1$  and  $p_2$  are effective descent morphisms, then so is  $(p_1, p_2)$ .

*Example 2.7.* Let us choose the data above as follows:

- $\mathcal{C}_0$  to be the category of sets.
- $\bullet \ {\mathfrak C}_1 = {\sf Ord}.$
- $\mathcal{C}_2$  to be the category of pairs  $(A, (A_x)_{x \in X})$ , where A is a set and  $(A_x)_{x \in X}$  is a family of subsets of A with  $x' \leq x \Rightarrow A_x \subseteq A_{x'}$ . A morphism

$$f: (A, (A_x)_{x \in X}) \to (B, (B_x)_{x \in X})$$

in  $\mathcal{C}_2$  is a morphism  $f: A \to B$  in  $\mathcal{C}_0$  with  $f(A_x) \subseteq B_x$  for each  $x \in X$ .

•  $S_1$  and  $S_2$  to be the underlying set functors. Note that in all our categories here the pullbacks and coequalizers are *chosen as for sets* and preserved by the functors  $S_1$  and  $S_2$ ; moreover, one could easily choose suitable right inverses of  $S_1$  and  $S_2$ .

We can write the objects of  $C_1 \times_{C_0} C_2$  simply as  $A = (A, (A_x)_{x \in X})$  assuming that A is a set equipped with a preorder relation and a family  $(A_x)_{x \in X}$  of subsets satisfying the condition above. A morphism  $p: E \to B$  in  $C_1 \times_{C_0} C_2$  is then a map  $p: E \to B$  that preserves the preorder relation and has  $p(E_x) \subseteq B_x$  for each  $x \in X$ . It is easy to show that p is an effective descent morphism in  $C_2$  if and only if p and all induced maps  $p_x: E_x \to B_x$  are surjective. Then, putting this together with the description of effective descent morphisms in Ord in Proposition 3.4 of [2], and applying Corollaries 2.3 and 2.6, we conclude that

**Theorem 2.8.** A morphism  $p: E \to B$  in  $C_1 \times_{C_0} C_2$  is an effective descent morphism if and only if:

- (a) for every  $b_2 \leq b_1 \leq b_0$  in B, there exist  $e_2 \leq e_1 \leq e_0$  in E with  $p(e_i) = b_i$  for each i = 0, 1, 2;
- (b) p and all induced maps  $p_x \colon E_x \to B_x$  are surjective.

## 3. X-FILTERED PREORDERS

Definition 3.1. An X-filtered preorder is a pair (A, a), where  $a: X \to A$  is a  $\{0, 1\}$ -profunctor; that is,  $a \subseteq X \times A$  is a relation satisfying

$$(x' \leqslant x \And a \leqslant a') \Rightarrow ((x, a) \in \mathbf{a} \Rightarrow (x', a') \in \mathbf{a})$$

for all  $x, x' \in X$  and  $a, a' \in A$ . A morphism  $f: (A, \mathbf{a}) \to (B, \mathbf{b})$  in the category  $\operatorname{Ord}_X$  of X-filtered preorders is a morphism  $f: A \to B$  in  $\operatorname{Ord}$  with  $f\mathbf{a} \leq \mathbf{b}$ , that is, with

$$(x,a) \in \mathsf{a} \Rightarrow (x,f(a)) \in \mathsf{b}$$

for all  $x \in X$  and  $a \in A$ .

The reason for calling such pairs  $(A, \mathbf{a})$  X-filtered preorders is that a profunctor  $\mathbf{a} \colon X \to A$ can be equivalently described as an X-filtration on A defined as a family  $(A_x)_{x \in X}$  of upclosed subsets of A with  $x' \leq x \Rightarrow A_x \subseteq A_{x'}$ . The relationship between these two types of structure is straightforward and well known at various levels of generality; it is given by

$$A_x = \mathsf{a}(x, -) = \{a \in A \mid (x, a) \in \mathsf{a}\}$$

or, equivalently, by

$$\mathsf{a} = \{(x, a) \in X \times A \mid a \in A_x\}$$

If there is no danger of confusion, we will simply write

$$(A, \mathsf{a}) = A = (A, (A_x)_{x \in X})$$

(the second equality here is what we have already used in a more general situation in Example 2.7).

We could also similarly describe these structures as families of A-indexed (for various A) families of subsets

$$\mathsf{a}(-,a) = \{ x \in X \mid (x,a) \in \mathsf{a} \}$$

of X, but that would be less useful since we are considering a fixed X, not a fixed A.

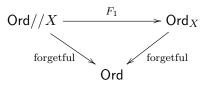
We will use the fully faithful functors

$$\operatorname{Ord}//X \xrightarrow{F_1} \operatorname{Ord}_X \xrightarrow{F_2} \operatorname{\mathcal{C}}_1 \times_{\operatorname{\mathcal{C}}_0} \operatorname{\mathcal{C}}_2$$

where  $F_1$  is defined by

$$F_1(A,\alpha) = (A, \{(x,a) \mid x \leq \alpha(a)\})$$

and by requiring the diagram



to commute,  $C_1 \times_{C_0} C_2$  is as in Example 2.7, and  $F_2$  is the inclusion functor. The following proposition is also well known at various levels of generality, but instead of explaining that we give a (simple) direct proof:

**Proposition 3.2.** The image of  $F_1: \operatorname{Ord}//X \to \operatorname{Ord}_X$ , which is the same as its replete image, consists of those filtered preorders A in which, for every  $a \in A$ , the subset a(-, a) of X has a largest element; equivalently, there is a largest  $x \in X$  with  $a \in A_x$ .

*Proof.* For  $(A, \alpha)$  and  $(B, \beta)$  in Ord/X, a morphism  $f: A \to B$  in Ord, and  $a \in A$ , we have

$$\alpha(a) \leqslant \beta f(a) \Leftrightarrow \forall_x (x \leqslant \alpha(a) \Rightarrow x \leqslant \beta f(a)),$$

which means that f is a morphism from  $(A, \alpha)$  to  $(B, \beta)$  in  $\operatorname{Ord}//X$  if and only if it is a morphism from  $F_1(A, \alpha)$  to  $F_1(B, \beta)$  in  $\operatorname{Ord}_X$ . That is,  $F_1$  is indeed fully faithful.

Let  $(A, \mathbf{a})$  be an object in  $\operatorname{Ord}_X$  such that each  $\mathbf{a}(-, a)$  has a largest element. Let  $\alpha(a)$  be such a largest element (for each  $a \in A$ ), and consider the so defined  $(A, \alpha)$ . We have

$$a \leqslant a' \Rightarrow \alpha(a) \in \mathsf{a}(-,a') \Rightarrow \alpha(a) \leqslant \alpha(a'),$$

and so  $(A, \alpha)$  is an object in  $\operatorname{Ord}//X$ . We also have  $x \leq \alpha(a) \Leftrightarrow (x, a) \in a$ , and so  $F_1(A, \alpha) = (A, a)$ .

Conversely, if  $(A, \mathbf{a}) = F_1(A, \alpha) = (A, \{(x, a) \mid x \leq \alpha(a)\})$ , then, obviously, for each  $a \in A$ ,  $\alpha(a)$  is the largest  $x \in X$  with  $x \leq \alpha(a)$ , which is equivalent to  $x \in \mathbf{a}(-, a)$ .

#### 4. Effective descent morphisms in $Ord_X$

We will often use a pullback diagram

$$E \times_B A \xrightarrow{\pi_2} A$$
$$\begin{array}{c} \pi_1 \\ \mu \\ E \xrightarrow{n} B \end{array} \xrightarrow{\pi_2} B$$

in a given category, and refer to it simply as the pullback of p and f. Recall, e.g. from part 2 of Corollary 2.7 in [3]:

**Proposition 4.1.** Let  $\mathcal{D}$  be a category with pullbacks and  $\mathcal{C}$  a full subcategory in  $\mathcal{D}$  closed under pullbacks. For an effective descent morphism  $p: E \to B$  in  $\mathcal{D}$ , the following conditions are equivalent:

- (i) p is an effective descent morphism in C;
- (ii) for every morphism  $f: A \to B$  in  $\mathcal{D}$ , we have  $E \times_B A \in \mathfrak{C} \Rightarrow A \in \mathfrak{C}$ .

Using this proposition applied to the inclusion  $F_2: \operatorname{Ord}_X \to \mathcal{C}_1 \times_{\mathcal{C}_0} \mathcal{C}_2$  and the results of Section 2, we will prove

**Theorem 4.2.** A morphism  $p: E \to B$  in  $Ord_X$  is an effective descent morphism if and only if:

- (a) for each  $b_2 \leq b_1 \leq b_0$  in B, there exist  $e_2 \leq e_1 \leq e_0$  in E with  $p(e_i) = b_i$  for each i = 0, 1, 2;
- (b) for each  $x \in X$  and each  $b_1 \leq b_0$  in  $B_x$ , there exist  $e_1 \leq e_0$  in  $E_x$  with  $p(e_i) = b_i$  for each i = 0, 1.

*Proof.* "Only if": Suppose p is an effective descent morphism in  $\operatorname{Ord}_X$ . In Corollary 2.3, take  $S_1$  to be the forgetful functor  $\operatorname{Ord}_X \to \operatorname{Ord}$  and  $J_1: \operatorname{Ord} \to \operatorname{Ord}_X$  to be defined by

$$J_1(A) = (A, (A_x)_{x \in X})$$
 with  $A_x = A$  for all  $x \in X$ .

According to the description of effective descent morphisms in Ord, it follows that condition (a) is satisfied.

To prove that condition (b) is also satisfied, take any  $x \in X$  and  $b_1 \leq b_0$  in  $B_x$ , and consider the pullback of p and f in  $\mathcal{C}_1 \times_{\mathcal{C}_0} \mathcal{C}_2$  with  $A = \{b_1, b_0\}$  having  $b_i \leq b_j \Leftrightarrow j \leq i$ , f being the inclusion map, and

$$A_{x'} = \begin{cases} \{b_1, b_0\}, & \text{if } x' < x; \\ \{b_1\}, & \text{if } x' \sim x; \\ \emptyset, & \text{if } x' \notin x, \end{cases}$$

where by  $x' \sim x$  we mean  $x' \leq x$  and  $x \leq x'$ , and by x' < x we mean  $x' \leq x$  and  $x' \not\sim x$ . In this case

$$E \times_B A = (p^{-1}\{b_1\} \times \{b_1\}) \cup (p^{-1}\{b_0\} \times \{b_0\})$$

with

$$(E \times_B A)_{x'} = \begin{cases} ((E_{x'} \cap p^{-1}\{b_1\}) \times \{b_1\}) \cup ((E_{x'} \cap p^{-1}\{b_0\}) \times \{b_0\}), & \text{if } x' < x; \\ (E_{x'} \cap p^{-1}\{b_1\}) \times \{b_1\}, & \text{if } x' \sim x; \\ \emptyset, & \text{if } x' \leq x. \end{cases}$$

Since  $A_x$  is not upclosed, A does not belong to  $\operatorname{Ord}_X$ . Hence, since p is an effective descent morphism in  $\operatorname{Ord}_X$ , it follows, by Proposition 4.1, that  $E \times_B A$  does not belong to  $\operatorname{Ord}_X$ . Therefore, at least one  $(E \times_B A)_{x'}$  is not upclosed in  $E \times_B A$ . However, suppose x' < x. Then

$$(E \times_B A)_{x'} = (E_{x'} \cap p^{-1}\{b_1\}) \times \{b_1\}) \cup ((E_{x'} \cap p^{-1}\{b_0\}) \times \{b_0\})$$

and, since  $E_{x'}$  is upclosed in E, it is upclosed in  $E \times_B A$ . Or, suppose  $x' \notin x$ . Then  $(E \times_B A)_{x'}$ is empty and so it is upclosed in  $E \times_B A$ . Hence upclosedness of  $(E \times_B A)_{x'}$  must fail for some  $x' \sim x$ ; that is,  $(E \times_B A)_{x'}$ , and therefore  $(E \times_B A)_x$ , is not upclosed in  $E \times_B A$ . Since  $E_x$  is upclosed in E,

$$(E \times_B A)_x = (E_x \cap p^{-1}\{b_1\}) \times \{b_1\}$$

is upclosed in  $p^{-1}{b_1} \times {b_1}$ , and so there exist  $(e_1, b_1) \in (E_x \cap p^{-1}{b_1}) \times {b_1}$  and  $(e_0, b_0) \in p^{-1}{b_0} \times {b_1}$  with  $(e_1, b_1) \leq (e_0, b_0)$ . This gives  $e_1 \leq e_0$ , both in  $E_x$ , since  $e_1$  belongs to  $E_x$  and  $E_x$  is upclosed in E, with  $p(e_1) = b_1$  and  $p(e_0) = b_0$ .

"If": Suppose conditions (a) and (b) hold. This makes p an effective descent morphism in  $\mathcal{C}_1 \times_{\mathcal{C}_0} \mathcal{C}_2$  (see Example 2.7). Therefore, according to Proposition 4.1, it suffices to prove that, for every morphism  $f: A \to B$  in  $\mathcal{C}_1 \times_{\mathcal{C}_0} \mathcal{C}_2$  with each  $(E \times_B A)_x$  upclosed, each  $A_x$  is also upclosed.

Suppose  $a \leq a'$  in A with  $a \in A_x$ . As follows from condition (b), there exist  $e, e' \in E$  with  $e \leq e', p(e) = a$ , and p(e') = a'. This gives  $(e, a) \leq (e', a')$  in  $E \times_B A$  with  $(e, a) \in (E \times_B A)_x$ . Since  $(E \times_B A)_x$  is upclosed in  $E \times_B A$ , it follows that (e', a') is in  $(E \times_B A)_x$ , and so a' is in  $A_x$ , since the projection  $\pi_2 \colon E \times_B A \to A$  is a morphism in  $\mathcal{C}_1 \times_{\mathcal{C}_0} \mathcal{C}_2$ .

#### 5. Effective descent morphisms in Ord//X

In this section we assume that X is *locally complete*, in the sense that, for each  $x \in X$ , the preorder  $\{x' \in X \mid x' \leq x\}$  is equivalent to a complete lattice. We will also identify the category  $\operatorname{Ord}/X$  with its  $F_1$ -image in  $\operatorname{Ord}_X$ ; note that  $\operatorname{Ord}/X$  is then closed under pullbacks in  $\operatorname{Ord}_X$ , thanks to the local completeness. While working with the morphism  $p: E \to B$  we will write  $B = (B, \beta), E = (E, \varepsilon)$ , and  $E \times_B A = (E \times_B A, \gamma)$  in the notation we used for  $\operatorname{Ord}/X$ .

**Lemma 5.1.** Let  $p: E \to B$  be a morphism in  $\operatorname{Ord}//X$  such that each induced map  $p_x: E_x \to B_x \ (x \in X)$  is surjective. If  $f: A \to B$  is a morphism in  $\operatorname{Ord}_X$  with  $E \times_B A$  in  $\operatorname{Ord}//X$ , then A is in  $\operatorname{Ord}//X$ .

*Proof.* According to Proposition 3.2, we have to prove that, given  $a \in A$ , there is a largest  $x \in X$  with  $a \in A_x$ . Following the proof of Theorem 3.9 in [1], we are going to put

$$\alpha(a) = \bigvee \{ x \in X \mid a \in A_x \}$$

and prove that it is such an element. First we observe that, if a belongs to  $A_x$ , then f(a) belongs  $B_x$  and so  $x \leq \beta f(a)$ . Therefore

$$\{x \in X \mid a \in A_x\} = \{x \leqslant \beta f(a) \mid a \in A_x\},\$$

and so the join above is well defined. By definition of  $\alpha(a)$ , it suffices to prove that  $a \in A_{\alpha(a)}$ . As  $f(a) \in B_{\alpha(a)}$ , we have  $\alpha(a) \leq \beta f(a)$ , and so f(a) belongs to  $B_{\alpha(a)}$ . By our hypotheses, there exist  $e \in E_{\alpha(a)}$  with p(e) = f(a). Hence, for every  $x \leq \alpha(a)$ , (e, a) belongs to  $E_x \times_B A_x = (E \times_B A)_x$ , and so  $x \leq \gamma(e, a)$ . Since  $\alpha(a)$  is the join of such elements x, it follows that  $\alpha(a) \leq \gamma(e, a)$ . Then  $a = \pi_2(c, a) \in A_{\gamma(c, a)} \subseteq A_{\alpha(a)}$ , which completes our proof.

From Proposition 4.1, Theorem 4.2, and Lemma 5.1, we obtain:

**Theorem 5.2.** If X is a locally complete preordered set, then a morphism  $p: E \to B$  in Ord//X is effective for descent in Ord//X provided that:

- (a) for each  $b_2 \leq b_1 \leq b_0$  in B, there exist  $e_2 \leq e_1 \leq e_0$  in E with  $p(e_i) = b_i$  for each i = 0, 1, 2:
- (b') for each  $x \in X$  and each  $b_1 \leq b_0$  with  $x \leq \beta(b_1)$ , there exist  $e_1 \leq e_0$  with  $x \leq \varepsilon(e_1)$ and  $p(e_i) = b_i$  for i = 0, 1.

Our next result shows that conditions (a) and (b') characterize effective descent morphisms in  $\operatorname{Ord}//X$  among those  $p: E \to B$  with  $p_x: E_x \to B_x$  surjective for every  $x \in X$ . **Theorem 5.3.** Let X be a locally complete preordered set with bottom element. A morphism  $p: E \to B$  in Ord//X such that  $p_x: E_x \to B_x$  is surjective for every  $x \in X$  is effective for descent in Ord//X if and only if:

- (a) for each  $b_2 \leq b_1 \leq b_0$  in B, there exist  $e_2 \leq e_1 \leq e_0$  in E with  $p(e_i) = b_i$  for each i = 0, 1, 2;
- (b') for each  $x \in X$  and each  $b_1 \leq b_0$  with  $x \leq \beta(b_1)$ , there exist  $e_1 \leq e_0$  with  $x \leq \varepsilon(e_1)$ and  $p(e_i) = b_i$  for i = 0, 1.

*Proof.* We only need to prove the necessity of conditions (a) and (b'). Its proof follows directly the "only if" proof of Theorem 4.2. Let  $\perp$  be the bottom element of X. As in the proof of Theorem 4.2 we apply Corollary 2.3 to the functors

$$\operatorname{Ord}//X \xrightarrow{S_1} \operatorname{Ord}$$

where  $S_1$  is the forgetful functor and  $J_1$  assigns to each preordered set A the pair  $(A, \perp)$ , with  $\perp(a) = \perp$  for every  $a \in A$ , and conclude that an effective descent morphism in Ord//X is in particular effective for descent in Ord, that is, it satisfies condition (a).

To show the necessity of (b'), let  $x \in X$  and  $b_1 \leq b_0$  in B with  $x \leq \beta(b_1)$ . Given a pullback diagram of p and f in  $\mathcal{C}_0 \times_{\mathcal{C}_1} \mathcal{C}_2$  with A and  $A_x$  as described in the proof of 4.2, both A and  $E \times_B A$  do not belong to  $\operatorname{Ord}_X$  and therefore they do not belong to  $\operatorname{Ord}/X$ .

There is another convenient way to express the surjectivity of each  $p_x \colon E_x \to B_x$ :

**Proposition 5.4.** The following conditions on a morphism  $p: E \to B$  in Ord//X are equivalent:

- (i)  $p_x: E_x \to B_x$  is surjective for every  $x \in X$ ;
- (ii) for every  $b \in B$  there exists  $e \in E$  with p(e) = b and  $\varepsilon(e) \sim \beta(b)$ .

*Proof.* (ii) $\Rightarrow$ (i) is obvious and so we only need to prove (i) $\Rightarrow$ (ii). Suppose (i) holds. Given  $b \in B$ , since  $b \in B_{\beta(b)}$ , there exists  $e \in E_{\beta(b)}$  with p(e) = b. Then we have  $\beta(b) \leq \varepsilon(e)$  since e belongs to  $E_{\beta(b)}$ , and  $\varepsilon(e) \leq \beta p(e) = \beta(b)$  since p is a morphism in Ord//X.

**Lemma 5.5.** Suppose that, for every  $x \in X$ , every subset of  $\{x' \in X \mid x' \leq x\}$  has a largest element. If  $p: E \to B$  is a pullback stable extremal epimorphism in Ord//X, then  $p_x: E_x \to B_x$  is surjective for every  $x \in X$ .

*Proof.* Given  $b \in B$ , consider the pullback diagram

$$p^{-1}(b) \longrightarrow \{b\}$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow$$

$$E \longrightarrow B$$

in  $\operatorname{Ord}//X$ , where the vertical arrow are the inclusion maps, and  $\{b\} = (\{b\}, \beta')$  with  $\beta'(b) = \beta(b)$ ; accordingly  $p^{-1}(b) = (p^{-1}(b), \varepsilon')$  with  $\varepsilon'(e) = \varepsilon(e)$  for each  $e \in p^{-1}(b)$ . We have  $\varepsilon(e) \leq \beta(b)$  for each  $e \in p^{-1}(b)$ , and we define  $\beta'' : \{b\} \to X$  by taking  $\beta''(b)$  to be a largest element in the set  $\{\varepsilon(e) \mid e \in p^{-1}(b)\}$ . Then  $(p^{-1}(b), \varepsilon') \to (\{b\}, \beta')$ , which is the top arrow in our pullback diagram, factors through  $(\{b\}, \beta'') \to (\{b\}, \beta')$ . Since p is a pullback stable extremal epimorphism, it follows that  $(\{b\}, \beta'') \to (\{b\}, \beta')$  is an isomorphism. Therefore  $\beta''(b) \sim \beta'(b)$ , which implies the existence of  $e \in p^{-1}(b)$  with  $\varepsilon(e) \sim \beta(b)$  and so completes the proof.

From Theorem 5.3 and Lemma 5.5, we immediately obtain:

**Theorem 5.6.** Suppose that, for every  $x \in X$ , every subset of  $\{x' \in X \mid x' \leq x\}$  has a largest element. Then a morphism  $p: E \to B$  is effective for descent in Ord//X if and only if:

- (a) for each  $b_2 \leq b_1 \leq b_0$  in B, there exist  $e_2 \leq e_1 \leq e_0$  in E with  $p(e_i) = b_i$  for each i = 0, 1, 2;
- (b') for each  $x \in X$  and each  $b_1 \leq b_0$  with  $x \leq \beta(b_1)$ , there exist  $e_1 \leq e_0$  with  $x \leq \varepsilon(e_1)$ and  $p(e_i) = b_i$  for i = 0, 1.

*Remark* 5.7. Although Theorem 5.3 is stronger than Theorem 3.9 in [1], we still do not know how far it is from a complete characterization of effective descent morphisms in Ord//X. Theorem 5.6 answers this question, but only under a strong additional condition on X.

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