

Topological spaces, categorically

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The talk is based on joint work with M.M. Clementino and W. Tholen.

“The kinds of structures which actually arise in the practice of geometry and analysis are far from being ‘arbitrary’ . . . , as concentrated in the thesis that *fundamental* structures are themselves categories.”



F. William Lawvere.

Metric spaces, generalized logic, and closed categories.

Rend. Sem. Mat. Fis. Milano, 43:135–166 (1974), 1973.

Also in: *Repr. Theory Appl. Categ.* 1:1–37, 2002.

Examples

Metric spaces, $(P_+ = [0, \infty]^{\text{op}}, +, 0)$

X with $d : X \times X \rightarrow P_+$ such that

$$0 \geq d(x, x), \quad d(x, y) + d(y, z) \geq d(x, z).$$

Categories, $(\text{Set}, \times, 1)$

X with $\text{hom} : X \times X \rightarrow \text{Set}$ such that

$$1 \rightarrow \text{hom}(x, x), \quad \text{hom}(x, y) \times \text{hom}(y, z) \rightarrow \text{hom}(x, z)$$

and ... (commutative diagrams in Set).

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and ... (commutative diagrams in Set).

Ordered sets, $(2 = \{\text{false}, \text{true}\}, \&, \text{true})$

X with $\leq : X \times X \rightarrow 2$ such that

$$\text{true} \models (x \leq x), \quad (x \leq y \& y \leq z) \models x \leq z.$$

The ordered category V-Rel

Quantale

$V = (V, \otimes, k)$ commutative unital quantale with $u \otimes - \dashv \text{hom}(u, -)$.

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- ▶ Involution: $r^\circ : Y \dashv\rightarrow X$ where $r^\circ(y, x) = r(x, y)$ for $r : X \dashv\rightarrow Y$.
- ▶ For each Set-map $f: f \dashv f^\circ$.

V-categories

A **V-category** is a pair $(X, a : X \multimap X)$ such that

$$k \leq a(x, x)$$

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respectively

$$\text{id}_X \leq a$$

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V-functors

A **V-functor** $f : (X, a) \longrightarrow (Y, b)$ is a Set-map such that

$$a(x, x') \leq b(f(x), f(x')) \qquad \text{respectively} \qquad f \cdot a \leq b \cdot f.$$

Topological spaces $2 = (2, \&, \text{true})$, $\mathbb{U} = (U, e, m)$

X with $\longrightarrow: UX \dashrightarrow X$ such that

$$\text{true} \models (\dot{x} \longrightarrow x), \quad (\ddot{x} \longrightarrow x \& x \longrightarrow x) \models m_X(\ddot{x}) \longrightarrow x.$$

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In fact, $U: \text{Set} \longrightarrow \text{Set}$ can be extended to a functor
 $U: \text{Rel} \longrightarrow \text{Rel}$ such that e and m become oplax.

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3. $\varphi : X \multimap Y$ is a \mathbf{V} -module iff $\varphi : X^{\text{op}} \otimes Y \longrightarrow \mathbf{V}$ is a \mathbf{V} -functor.
4. In particular $a : X^{\text{op}} \otimes X \longrightarrow \mathbf{V}$ is a \mathbf{V} -functor. Its mate $y = \ulcorner a \urcorner : X \longrightarrow \mathbf{V}^{X^{\text{op}}}$ is fully faithful. More general, we have

$$[y(x), \varphi] = \varphi(x).$$

5. ...

Topological theory

Definition

A **topological theory** \mathcal{T} is a triple $\mathcal{T} = (\mathbb{T}, \mathbb{V}, \xi)$ consisting of
a monad $\mathbb{T} = (T, e, m)$, a quantale $\mathbb{V} = (V, \otimes, k)$ and
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such that

$$(M_e) \text{id}_V \leq \xi \cdot e_V,$$

$$(M_m) \xi \cdot T\xi \leq \xi \cdot m_V,$$

$$(Q_\otimes) \quad \begin{array}{ccc} T(V \times V) & \xrightarrow{T(\otimes)} & TV \\ \langle \xi \cdot T\pi_1, \xi \cdot T\pi_2 \rangle \downarrow & \leq & \downarrow \xi \\ V \times V & \xrightarrow{\otimes} & V, \end{array}$$

$$(Q_k) \quad \begin{array}{ccc} T1 & \xrightarrow{Tk} & TV \\ ! \downarrow & \leq & \downarrow \xi \\ 1 & \xrightarrow{k} & V, \end{array}$$

$(Q_V) (\xi_X)_X : P_V \longrightarrow P_V T$ is a natural transformation.

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- ▶ $\mathcal{U}_{P_+} = (\mathbb{U}, P_+, \xi_{P_+})$ is a strict topological theory, where

$$\xi_{P_+} : UP_+ \longrightarrow P_+, \quad x \longmapsto \inf\{v \in P_+ \mid x \in T([0, v])\}.$$

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- ▶ $\mathcal{L}_V^\otimes = (\mathbb{I}, V, \xi_\otimes)$ is a strict topological theory where

$$\xi_\otimes : LV \longrightarrow V, \quad (v_1, \dots, v_n) \longmapsto v_1 \otimes \dots \otimes v_n.$$

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Given $r : X \times Y \rightarrow V$, we put

$$T_\xi r : TX \times TY \rightarrow V$$

$$(\mathfrak{x}, \mathfrak{y}) \mapsto \bigvee \left\{ \xi \cdot Tr(w) \mid w \in T(X \times Y), w \mapsto \mathfrak{x}, \mathfrak{y} \right\},$$

that is,

$$\begin{array}{ccc} T(X \times Y) & \xrightarrow{\tau_{X,Y}} & TX \times TY \\ & \searrow \xi \cdot Tr & \swarrow T_\xi r \\ & V & \end{array} \quad \leq$$

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Theorem

The following statements hold.

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3. $T_\xi s \cdot T_\xi r \leq T_\xi(s \cdot r)$ provided that T satisfies (BC), and $\overline{T_\xi s} \cdot \overline{T_\xi r} \geq \overline{T_\xi(s \cdot r)}$ provided that (Q_\otimes^-) holds.

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4. The natural transformations e and m become op-lax, that is, for every V-relation $r : X \dashrightarrow Y$ we have the inequalities:

$$\begin{array}{ccc} X & \xrightarrow{e_X} & TX \\ r \downarrow & \leq & \downarrow T_\xi r \\ Y & \xrightarrow{e_Y} & TY \end{array}$$

$$\begin{array}{ccc} TTX & \xrightarrow{m_X} & TX \\ T_\xi T_\xi r \downarrow & \leq & \downarrow T_\xi r \\ TTY & \xrightarrow{m_Y} & TY \end{array}$$

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$$b \circ a = b \cdot T_{\xi} a \cdot m_X^{\circ}.$$

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- ▶ $a \circ (b \circ c) \geq a \circ b \circ c \leq (a \circ b) \circ c$.
- ▶ If \mathcal{T} is a strict theory, then Kleisli convolution is associative.

V-Rel vs. \mathcal{T} -Rel

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We consider now

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We have

- ▶ $(1_Y)_\# \circ a = e_Y^\circ \circ a$ and $a \circ (1_X)_\# = a \circ e_X^\circ$.
- ▶ $r_\#$ is unitary.
- ▶ T satisfies (BC) $\Rightarrow s_\# \circ r_\# \leq (s \cdot r)_\#$.
- ▶ $(Q_{\otimes}^-) \Rightarrow s_\# \circ r_\# \geq (s \cdot r)_\#$.

\mathcal{T} -category

A \mathcal{T} -category is a pair $(X, a : TX \rightarrow X)$ such that

$$k \leq a(e_X(x), x), \quad T_\xi a(x, x) \otimes a(x, x) \leq a(m_X(x), x) \quad \text{respectively}$$

$$\text{id}_X \leq a \cdot e_X, \quad a \cdot T_\xi a \leq a \cdot m_X \quad \text{respectively}$$

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A map $f : (X, a) \rightarrow (Y, b)$ is a \mathcal{T} -functor if

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- ▶ $\mathcal{L}_V^\otimes\text{-Cat} \cong V\text{-MultiCat}$.

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- ▶ $\mathcal{L}_V^\otimes\text{-Cat} \cong V\text{-MultiCat}$.

From now on we consider a **strict** theory $\mathcal{T} = (\mathbb{T}, V, \xi)$.

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and we have

$$\begin{array}{ccc} & \mathcal{T}\text{-Cat} & \\ & \nearrow^{(-)_{\#}} & \searrow^M \\ \text{V-Cat} & \xrightarrow{T_{\xi}} & \text{V-Cat} \end{array}$$

where $M : \mathcal{T}\text{-Cat} \longrightarrow \text{V-Cat}$, $(X, a) \longmapsto (TX, T_{\xi}a \cdot m_X^{\circ})$.

The \mathcal{T} -category V

We define

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2. $\text{hom}(v, _) : V \longrightarrow V$ is a \mathcal{T} -functor for each $v \in V$ which satisfies $\xi \cdot Tv \geq v \cdot !$.

The \mathcal{T} -category V

We define

$$\text{hom}_\xi : TV \times V \longrightarrow V, (v, v) \longmapsto \text{hom}(\xi(v), v).$$

Then $V = (V, \text{hom}_\xi)$ is a \mathcal{T} -category.

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3. $v \otimes _ : V \longrightarrow V$ is a \mathcal{T} -functor for each $v \in V$ which satisfies $\xi \cdot Tv \leq v \cdot !$.

Compatible monoidal structures on V

We assume that a monoidal structure (V, \oplus, I) on V is given such that

1. $(u_1 \oplus v_1) \otimes (u_2 \oplus v_2) \leq (u_1 \otimes u_2) \oplus (v_1 \otimes v_2)$,
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$$\begin{array}{ccc}
 3. & T(V \times V) \xrightarrow{T(\oplus)} TV & \text{and} & T1 \xrightarrow{T I} TV \\
 \langle \xi \cdot T\pi_1, \xi \cdot T\pi_2 \rangle \downarrow & \geq & & \downarrow \xi \\
 V \times V \xrightarrow{\oplus} V, & & & 1 \xrightarrow{I} V.
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Examples

- ▶ $\oplus = \otimes$ (since \mathcal{T} is strict).
- ▶ $\oplus = \wedge$.

Monoidal structures on V-Rel

Extending \oplus to V-Rel

- ▶ For sets X and Y we put $X \oplus Y = X \times Y$.
- ▶ For V-relations $r : X \twoheadrightarrow X'$ and $s : Y \twoheadrightarrow Y'$ we define $r \oplus s : X \times Y \twoheadrightarrow X' \times Y'$ by

$$r \oplus s((x, y), (x', y')) = r(x, x') \oplus s(y, y').$$

Then $\oplus : \text{V-Rel} \times \text{V-Rel} \longrightarrow \text{V-Rel}$ is a lax functor, is associative and with $l : 1 \twoheadrightarrow 1$ as neutral element.

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Of course, we obtain a monoidal structure on V-Cat where $(X, a) \oplus (Y, b) = (X \times Y, a \oplus b)$ with neutral element $E = (1, l)$.

Hopf monad

A **Hopf monad** on a monoidal category E is a monad $\mathbb{T} = (T, e, m)$ on E equipped with a natural transformation

$$\tau : T(- \otimes -) \longrightarrow T(-) \otimes T(-)$$

and a map $\theta : T(N) \longrightarrow N$ such that ...

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Theorem

There is a bijective correspondence between such structures τ, θ on \mathbb{T} and liftings of the monoidal structure on E to $E^{\mathbb{T}}$.

Here:

$$(X, \alpha) \otimes (Y, \beta) = (X \otimes Y, (\alpha \otimes \beta) \cdot \tau_{X,Y}).$$

Lax Hopf monad

With $\tau_{X,Y} : T(X \times Y) \rightarrow TX \times TY$ and $! : T1 \rightarrow 1$, in our situation we have

$$\begin{array}{ccc} T(X \oplus Y) & \xrightarrow{\tau_{X,Y}} & TX \oplus TY \\ \downarrow T_{\xi}(r \oplus s) & \leq & \downarrow T_{\xi}r \oplus T_{\xi}s \\ T(X' \oplus Y') & \xrightarrow{\tau_{X',Y'}} & TX' \oplus TY' \end{array} \quad \text{and} \quad \begin{array}{ccc} T1 & \xrightarrow{!} & 1 \\ \downarrow T_{\xi}l & \leq & \downarrow l \\ T1 & \xrightarrow{!} & 1 \end{array}$$

making (T_{ξ}, e, m) a **lax Hopf monad** on V-Rel.

Extending \oplus to \mathcal{T} -Rel. . .

Let $r : X \rightarrow X'$ and $s : Y \rightarrow Y'$ be \mathcal{T} -relations. We put $X \boxplus Y = X \times Y$ and define $r \boxplus s : X \times Y \rightarrow X' \times Y'$ as

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and $l_1 : 1 \rightarrow 1$ as the composite $T1 \xrightarrow{!} 1 \xrightarrow{!} 1$.

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- ▶ $e_X^\circ \boxplus e_Y^\circ \geq e_{X \times Y}^\circ$,
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For $(-)_\# : \mathbf{V}\text{-Rel} \rightarrow \mathcal{T}\text{-Rel}$ we have

- ▶ $(r \oplus r')_\# \leq r_\# \boxplus r'_\#$.
- ▶ $l_\# \leq l_!$.

... and to \mathcal{T} -Cat

Theorem

Each monoidal structure (V, \oplus, I) on V compatible with \mathcal{T} defines a monoidal structure on \mathcal{T} -Cat where $(X, a) \oplus (Y, b) = (X \times Y, a \boxplus b)$ with neutral element $E = (1, I)$.

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$$\tau_{X,Y} : M(X \oplus Y) \rightarrow M(X) \oplus M(Y) \quad \text{and} \quad ! : M(E) \rightarrow E.$$

Closedness of \mathcal{T} -Gph

Assume now that $u \oplus _ : V \rightarrow V$ has right adjoint $u \multimap _ : V \rightarrow V$.

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Let $X = (X, a)$, $Y = (Y, b)$ be \mathcal{T} -graphs. Then

$$X \multimap Y = \{f : X \rightarrow Y \mid f : X \oplus G \rightarrow Y \text{ is a } \mathcal{T}\text{-functor}\}$$

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(where $G = (1, e_X^\circ)$) with structure

$$a \multimap b(p, h) = \bigwedge_{\substack{q \in T(X \times (X \multimap Y)), x \in X \\ q \mapsto p}} (a(T\pi_X(q), x) \multimap b(\text{TeV}(q), h(x))).$$

is a \mathcal{T} -graph as well. In fact, $X \oplus _ \dashv X \multimap _$.

Closed \mathcal{T} -categories

Lemma

$$\begin{array}{ccc} T(V \times V) & \xrightarrow{T(\oplus)} & TV \\ \langle \xi \cdot T\pi_1, \xi \cdot T\pi_2 \rangle \downarrow & & \downarrow \xi \\ V \times V & \xrightarrow{\oplus} & V \end{array} \Rightarrow \begin{array}{ccc} T(X \times Y) & \xrightarrow{\tau_{X,Y}} & TX \times TY \\ T_\xi(r \oplus s) \downarrow & & \downarrow T_\xi r \oplus T_\xi s \\ T(X' \times Y') & \xrightarrow{\tau_{X',Y'}} & TX' \times TY'. \end{array}$$

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Theorem

(V, \oplus, I) closed, strictly compatible with \mathcal{T} ; $X = (X, a) \in \mathcal{T}\text{-Cat}$.

1. $a \dashv\dashv b$ is transitive for each \mathcal{T} -category $Y = (Y, b)$ if

$$(*) \quad \bigvee_{\mathfrak{x} \in TX} (T_\xi a(\mathfrak{x}, \mathfrak{x}) \oplus u) \otimes (a(\mathfrak{x}, x_0) \oplus v) \geq a(m_X(\mathfrak{x}), x_0) \oplus (u \otimes v).$$

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2. If $a \multimap \text{hom}_\xi$ is transitive, then $(*)$ for all $\mathfrak{x} \in T^2X$, $x_0 \in X$ and $u, v \in V$ with $\xi \cdot Tu = u!$ and $\xi \cdot Tv \leq v!$.

Closed \mathcal{T} -categories

Corollary

Consider $\oplus = \otimes$. Let $X = (X, a)$ be a \mathcal{T} -category. Then

1. If $a \cdot T_{\xi} a = a \cdot m_X$, then $\text{hom}(a, b)$ is transitive for each \mathcal{T} -category $Y = (Y, b)$.
2. $a \cdot T_{\xi} a = a \cdot m_X$ provided that $\text{hom}(a, \text{hom}_{\xi})$ is transitive.

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4. If $Te_X \cdot e_X = m_X^\circ \cdot e_X$, then $X_\# = (X, r_\#)$ is closed for each V -category $X = (X, r)$.

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Let $X = (X, a)$ be a \mathcal{T} -category. TFAE.

- X is \oplus -compact.
- $\bigvee : (X \multimap V) \longrightarrow V$ is a \mathcal{T} -functor (where $X \oplus _ \dashv X \multimap _$).
- $\gamma : |X|_I \longrightarrow V$, $x \mapsto \bigvee_{x \in X} a(x, x)$ is a \mathcal{T} -functor.

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Corollary

A \mathcal{T} -category $X = (X, a)$ is \oplus -compact iff $\pi_Y : Y \oplus X \longrightarrow Y$ is closed for each \mathcal{T} -category $Y = (Y, b)$.

\mathcal{T} -modules

A \mathcal{T} -module $\varphi : (X, a) \rightarrow (Y, b)$ is a \mathcal{T} -relation $\varphi : X \rightarrow Y$ such that

$$b \circ \varphi \leq \varphi \quad \text{and} \quad \varphi \circ a \leq \varphi.$$

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Each \mathcal{T} -functor $f : (X, a) \longrightarrow (Y, b)$ defines \mathcal{T} -modules $f_* \dashv f^*$:

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$f : (X, a) \longrightarrow (Y, b)$ is fully faithful iff $a = (\text{id}_X)_* = f^* \circ f_*$.

Liftings and extensions

In V-Rel

For $\psi : X \dashrightarrow Z$, the composition maps

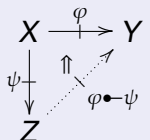
$$- \cdot \psi : \mathbf{V}\text{-Rel}(Z, Y) \longrightarrow \mathbf{V}\text{-Rel}(X, Y) \quad \text{and}$$

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have respective right adjoints

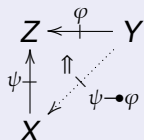
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(extension)

and



(lifting)

Liftings and extensions

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For $\psi : X \rightarrow Z$, the composition maps $_ \circ \psi$ still has a right adjoint but $\psi \circ _$ in general not.

Liftings and extensions

In \mathcal{T} -Rel

For $\psi : X \multimap Z$, the composition maps $_ \circ \psi$ still has a right adjoint but $\psi \circ _$ in general not. We pass from

$$\begin{array}{c} X \xrightarrow{\varphi} Y \\ \psi \downarrow \\ Z \end{array}$$

(in \mathcal{T} -Rel)

to

$$\begin{array}{c} TX \xrightarrow{\varphi} Y \\ m_X^\circ \downarrow \\ TTX \\ T_\xi \psi \downarrow \\ TZ \end{array}$$

(in V-Rel)

and define $\varphi \circ \psi = \varphi \bullet (T_\xi \psi \cdot m_X^\circ)$.

Modules as functors

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Theorem

For \mathcal{T} -categories (X, a) and (Y, b) , and a \mathcal{T} -relation $\psi : X \dashrightarrow Y$, the following assertions are equivalent.

- i. $\psi : (X, a) \dashrightarrow (Y, b)$ is a \mathcal{T} -module.
- ii. Both $\psi : |X| \otimes Y \rightarrow V$ and $\psi : X^{\text{op}} \otimes Y \rightarrow V$ are \mathcal{T} -functors.

L-separatedness/L-completeness

Let $X = (X, a)$ and $Y = (Y, b)$ be \mathcal{T} -categories. We consider

$$\begin{aligned} \alpha_{Y,X} : \mathcal{T}\text{-Cat}(Y, X) &\longrightarrow \mathcal{T}\text{-Map}(Y, X). \\ f &\longmapsto f_* \end{aligned}$$

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Note: It is enough to consider $Y = G = (1, e_1^\circ)$.

Examples

- ▶ In Met: L-complete=Cauchy-complete.
- ▶ In Top: L-complete=weakly sober.

Example: Top

Let X be a topological space. Then

- ▶ $M(X) = (UX, \leq)$ where $x \leq \eta$ if $\bar{x} \subseteq \eta$.

Example: Top

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$$\varphi \text{ is representable by } x \iff A = \overline{\{x\}}.$$

The Yoneda Lemma

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We put $\hat{X} = (\hat{X}, \hat{a})$ where

$$\hat{X} = \{\psi \in \mathbf{V}^{|X|} \mid \psi : X^{\text{op}} \longrightarrow \mathbf{V} \text{ is a } \mathcal{T}\text{-functor}\}$$

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If $T1 = 1$, we have a fully faithful functor $y : X \longrightarrow \hat{X}$.

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From now on we assume $T1 = 1$.

L-closure

Definition

Let $X = (X, a)$ be a \mathcal{T} -category. For $M \subseteq X$ we define

$$\overline{M} = \{x \in X \mid i^* \circ x_* \dashv x^* \circ i_*\}.$$

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Theorem

Then the following assertions are equivalent.

- i. $x \in \overline{M}$.
- ii. *For all \mathcal{T} -functors $\varphi, \psi : X \rightarrow Y$ with L-separated codomain: if $\varphi|_M = \psi|_M$, then $\varphi(x) = \psi(x)$.*
- iii. *For all \mathcal{T} -functors $\varphi, \psi : X \rightarrow V$: if $\varphi|_M = \psi|_M$, then $\varphi(x) = \psi(x)$.*

Further properties

- ▶ $f : X \longrightarrow Y$ is L-dense iff $f_* \circ f^* = (\text{id}_Y)_* = b$.

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$\psi \in \hat{X}$ is a right adjoint \mathcal{T} -module if and only if $\psi \in \overline{y[X]}$.

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Proof.

... $\varphi = (\text{id}_X)_* \circ \psi$ and observe that $\varphi(x) = \hat{a}(e_{\hat{X}}(\psi) y(x))$ and
 $\xi \cdot T\varphi(x) = T_\xi \hat{a}(T e_{\hat{X}} \cdot e_{\hat{X}}(\psi), T y(x)) \dots \quad \square$

L-completeness

We put $\tilde{X} = \overline{y[X]}$, then $y : X \longrightarrow \tilde{X}$ is fully faithful and dense.

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- ▶ X with $a \cdot T_\xi a = a \cdot m_X$, Y L-complete $\Rightarrow Y^X$ L-complete.
- ▶ $V^{|X|}$, \hat{X} , \tilde{X} are L-complete.