## Distant words

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(1) Combinatorics on infinite words
(2) Three notions of complexity
(3) Distant words

## Finite and infinite words

- Let $A$ be a finite set (alphabet). Sequences $w=w_{1} w_{2} \cdots$, finite or infinite, with $w_{i} \in A$ are called words.
- Let $A^{*}$ be the set of all finite words. $A^{*}$ is the free monoid generated by $A$ with the operation of concatenation of words. The empty word is the identity.
- The set of one-way infinite words over $A$ is denoted by $A^{\mathbb{N}}$.
- We will say that $u$ is a factor of the infinite word $w$ if $u=w_{i} w_{i+1} \cdots w_{j}$, with $i \leq j$.


## Thue-Morse infinite word

Let $h: A^{*} \rightarrow A^{*}$ satisfy $h(u v)=h(u) h(v)$ for all words $u$ and $v$. $h$ is a morphism from the free monoid into itself.

## Thue-Morse word:

Let $A=\{a, b\}$ and $\phi: A^{*} \rightarrow A^{*}$ be defined by
$\phi(a)=a b, \phi(b)=b a$.
Iterating this morphism, we obtain an infinite word $\phi(a)=a b, \phi(a b)=a b b a, \phi(a b b a)=a b b a b a a b$,

$$
t=\lim _{i \rightarrow \infty} \phi^{i}(a)=\text { abbabaabbaababba } \cdots
$$

## A second example: Fibonacci word

Here is another example of an infinite word obtained by morphism iteration.

Fibonacci word:

Let $\psi: A^{*} \rightarrow A^{*}$ be defined by $\psi(a)=a b, \psi(b)=a$.

$$
f=\lim _{i \rightarrow \infty} \phi^{i}(a)=\text { abaababaabaab } \cdots
$$

## Unavoidable regularities

An important topic in Combinatorics on Words is to decide which patterns are avoidable or not. Sometimes the answer depends on the size of the alphabet.

## Definition

A power of order $k$ is a pattern of the form $u^{k}$ for some finite word $u$ (e.g., a square if $k=2$, a cube if $k=3$ ).

Can we avoid squares, or cubes, inside some infinite word?

## Avoiding Powers

The Thue-Morse word is an example of infinite binary word which avoids cubes: does not contain 3 equal consecutive blocks.

We can say that the pattern $u^{3}$ is avoidable in a two letter alphabet already.

The property must be preserved by the defining morphism $\phi$.

## A new definition

We introduce now a dual notion of repetition. Why not look for anti-repetitions instead?

## Definition

An anti-power of order $k$ is a pattern of the form $v_{1} v_{2} \cdots v_{k}$, with $v_{i}$ all distinct and having the same length.

The finite word

$$
u=10110100
$$

is an example of anti-power of order 4.

## Powers or Anti-Powers

## Theorem (Fici, Restivo, Zamboni, Silva)

Let $w$ be an infinite word that avoids anti-powers of order $k$, for some $k>1$. Then, there exists a finite word $u$ such that $u^{n}$ is a factor of $w$ for each $n \geq 1$.

## Corollary

Every infinite word w admits powers or anti-powers of arbitrarily large order.

## Recurrence and uniform recurrence

## Definition

An infinite word $w$ is recurrent if all its factors show up twice.

## Definition

An infinite word $w$ is uniformly recurrent (almost periodic) if the distance between two consecutive occurrences of a given factor is bounded.

## The importance of being uniformly recurrent

## Theorem (Furstenberg - 1981)

Given an infinite word w, there exists an uniformly recurrent word $x$ such that $F(x) \subset F(w)$.

- It is enough to understand the structure of the class of uniformly recurrent words.
- Periodic words, $w=u^{\infty}$, are (non-interesting) examples of uniformly recurrent words.


## Aperiodic uniformly recurrent

Observe that an aperiodic uniformly recurrent word can never contain arbitrarily large powers of some factor $u$.

## Theorem (Fici, Restivo, Zamboni, Silva)

Every aperiodic uniformly recurrent word w contains anti-powers of arbitrarily large order.

## Factor complexity

- Given an infinite word w , let $F_{w}(n)$ be the set of factors of w with length $\mathrm{n} .1 \leq\left|F_{w}(n)\right| \leq|A|^{n}, n \leq 1$.
- If $w$ is periodic, then the function $F_{w}(n)$ is bounded.


## Theorem (Hedlund, Morse - 1938)

Given an infinite aperiodic word $w$ and $n \in \mathbb{N}$, we have

$$
\left|F_{w}(n)\right| \geq n+1
$$

The Thue-Morse word has linear factor complexity. A normal word has maximal complexity.

## Abelian compelxity

- Given a finite word $u \in A^{*}$, the Parikh vector $\alpha(u)=\left(u_{1}, \cdots, u_{k}\right)$ counts the number of ocorrences of each letter $a_{1}, a_{2}, \cdots a_{k}$ from the finite alphabet $A$.
- We say that $u \sim v$ (abelian equivalent) if they share the same Parikh vector. $v$ can be obtained from $u$ by permutating the letters of $u$.
- We define the function $G_{w}(n)$ which counts the number of distinct Parikh vector for factors of length $n$.


## Bounded abelian complexity

- $u$ is a k -abelian power if it admits a decomposition $u=v_{1} v_{2} \cdots v_{k}$ with $\left|v_{i}\right|=\left|v_{j}\right|$ and $v_{i} \sim v_{j}, 1 \leq i<j \leq k$.
- A word $w$ has bounded abelian complexity if any two factors of the same length $u$ e $v$ have at most $L$ distinct letters.


## Theorem (Richomme, Saari, Zamboni - 2011)

Let $w$ be an infinite word with bounded abelian complexity. Then, $w$ contains abelian powers of arbitrarily large order

## Cyclic complexity

- Two words $u$ and $v$ are conjugated if $u=w_{1} w_{2}$ and $v=w_{2} w_{1}$. This is an equivalence relation.
- Given an infinite word w, the function $H_{w}(n)$ counts the number of conjugacy classes of factors of length $n$.


## Theorem (Cassaigne, Fici, Sciortino, Zamboni - 2011)

Every aperiodic word w has unbounded cyclic complexity.

## Complexity notions

Each one of the three notions of complexity is associated to actions of different subgroups of the symmetric group.

- Factor complexity is associated with the trivial group,
- Abelian complexity is associated with the symmetric group
- Cyclic complexity is associated with the cyclic subgroup generated by the permutation $(1,2,3, \cdots, n)$.


## Far away and close by

Let $w$ be an infinite word in a finite alphabet A and $a \in A$.
$S_{1}=\left\{\left||u|_{a}-|v|_{a}\right|:|u|=|v|\right.$ and uv is a factor of $\left.w\right\}$.
$S_{2}=\left\{\left||u|_{a}-|v|_{a}\right|:|u|=|v|\right.$ and $u$ and $v$ are factors of $\left.w\right\}$.

## Proposition

$S_{1}$ is bounded if and only if $S_{2}$ is bounded.

This implies that an infinite word with unbounded abelian complexity will contain two consecutive blocks which are far apart.

## Consecutive factors distant from each other

The distance (Hamming) between two equal-length words is the number of positions at which the corresponding symbols differ.

## Proposition (Girão, Silva)

Given an infinite binary word w with an infinite number of zeros and ones and some $L>0$, we can find two same length consecutive factors $u$ and $v$ such that the distance is at least $L$.

## Theorem (Girão, Silva) <br> For every positive integers $k, d$, there exists $m(k, d)$ such that any word (on any alphabet) of size at least $m(k, d)$ contains either $k$-powers or $k$-anti-powers whose consecutive blocks are at distance at least $d$ between each two.

## The End

Thank you for your attention!

## References

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