

1. Tilings; n -Crosses and Lee Spheres
2. Perfect error correcting Lee codes
3. Golomb-Welch Conjecture
4. Known results on $PL(n, 2)$ codes
5. $PL(7, 2)$ codes

$PL(7, 2)$ and the Golomb-Welch Conjecture

5th Combinatorics Day
University of Beira Interior

April 17, 2015
Covilhã

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University of Aveiro

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1.2. n -Crosses

1.3. Lee-Spheres

2. Perfect error correcting Lee codes

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3. Golomb-Welch Conjecture

4. Known results on $PL(n, 2)$ codes

5. $PL(7, 2)$ codes

5.1 $3 \leq |G_i| \leq 8$

5.2 $3 \leq |G_i| \leq 7$

5.3 $4 \leq |G_i| \leq 7$

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Tiling R^n by unit cubes

$\mathcal{T} = \{T_i, i \in I\}$, $T_i \subset R^n$ **tiles** R^n if

- ▶ $\bigcup_{i \in I} T_i = R^n$
- ▶ $\text{int}(T_i) \cap \text{int}(T_j) = \emptyset$ for all $i \neq j$.

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In what follows T_i will denote a unit n -cube or a cluster of unit n -cubes.

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Two unit n -cubes are **twins** if they share a complete $n - 1$ face

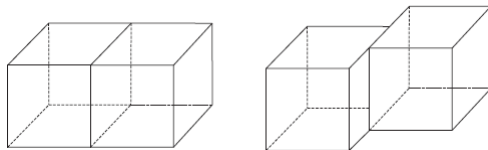


Figure: Twins and no twins.

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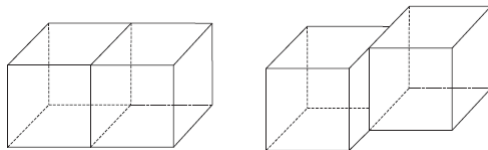


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Lattice Tiling R^n by unit cubes

Lattice: L is a **lattice** in R^n of dimension k if.

$$L = \{m_1 v_1 + m_2 v_2 + \dots + m_k v_k, m_i \in \mathbb{Z}, i = 1, \dots, k\}$$

with v_1, v_2, \dots, v_k linearly independent.

\hookrightarrow **group under vector addition** such that each of its points is the center of a ball that contains no other points.

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\leftrightarrow **group under vector addition** such that each of its points is the center of a ball that contains no other points.

Being L a lattice in R^n of **dimension n** , the set

$$F = \{x_1 v_1 + x_2 v_2 + \dots + x_n v_n, 0 \leq x_i \leq 1, i = 1, \dots, n\}$$

is called a **fundamental parallelepiped** of the lattice.

The **volume** of F is independent of the chosen basis and is called the **determinant of L** .

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Lattice Tiling R^n by unit cubes

A **Lattice tiling** of R^n **by unit cubes** is a tiling where the centers of the cubes form a lattice.

Minkowski's Conjecture (1896): Each lattice tiling of R^n by unit cubes contains twins.

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Geometrie der Zahlen p.105; 1896

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1907: Minkowski settled the case $n = 3$ *Diophantische Approximationen p.67-74.*

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1941: Hajós settled the Minkowski's Conjecture proving an equivalent conjecture about finite abelian groups.

Hajós's version of the Minkowski's Conjecture (1896):

let G be a finite abelian group.

If a_1, a_2, \dots, a_n are elements of G and r_1, r_2, \dots, r_n are positive integers such that each element of G is uniquely expressed in the form:

$$a_1^{x_1} \dots a_n^{x_n}, \quad 0 \leq x_1 \leq r_1 - 1, \dots, 0 \leq x_n \leq r_n - 1,$$

then $a_i^{r_i} = e$ for some $i \in \{1, 2, \dots, n\}$.

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Tilings by crosses

n -cross: Cluster consisting of $2n + 1$ unit cubes; a central cube where at each facet another unit cube is attached.



Figure: A 2-cross and a 3-cross.

Golomb & Welch (1968): there is a tiling of R^n by crosses for all $n \geq 2$.

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Szabó (1981): If $2n + 1$ is not a prime, then there exists a \mathbb{Q} -tiling of \mathbb{R}^n by crosses that is neither a \mathbb{Z} -tiling nor a lattice tiling.

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Lee spheres

By a Lee sphere centered at $W = (w_1, \dots, w_n) \in \mathbb{Z}^n$ of radius r , denoted by $S(W, r)$ we mean the set

$$S(W, r) = \{V = (v_1, \dots, v_n) \in \mathbb{Z}^n : \rho_L(V, W) = \sum_{i=1}^n |w_i - v_i| \leq r\}.$$

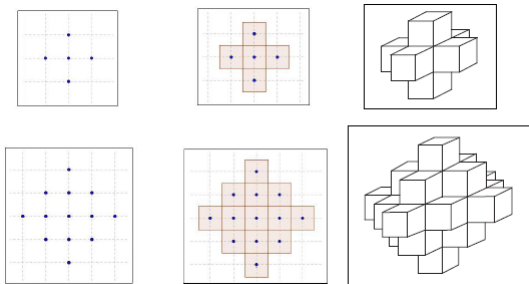


Figure: Lee sphere of radius 1 and 2 in \mathbb{Z}^2 and the corresponding unit cube clusters in \mathbb{R}^2 and \mathbb{R}^3 respectively.

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2.1. Perfect error correcting Lee codes

Perfect error correcting Lee codes

Let us consider the metric space (\mathbb{Z}^n, ρ_L) . Any subset \mathcal{M} of \mathbb{Z}^n , $|\mathcal{M}| \geq 2$, is called a **code**. The elements of \mathbb{Z}^n will be referred as **words** and in particular, the elements of \mathcal{M} will be called **codewords**.

A code \mathcal{M} is a **r -error correcting Lee code** if

$$(i) \quad \forall W, V \in \mathcal{M}, S(W, r) \cap S(V, r) = \emptyset.$$

If, in addition,

$$(ii) \quad \cup_{W \in \mathcal{M}} S(W, r) = \mathbb{Z}^n,$$

then \mathcal{M} is a **perfect r -error correcting Lee code of word length n over \mathbb{Z}** , shortly a **$PL(n, r)$ code**.

The elements $V \in S(W, r)$, with $W \in \mathcal{M}$ are said to be **covered** by W .

A $PL(n, r)$ code is a tiling of \mathbb{Z}^n by Lee spheres of radius r

inducing a tiling of \mathbb{R}^n by the corresponding cluster of unit cubes.

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2.1. Perfect error correcting Lee codes

$PL(2, r)$ codes

For each $r \geq 2$ there is a tiling of \mathbb{R}^2 by clusters of unit 2-cubes associated to Lee Spheres of radius r .

$PL(2, r)$ codes do exist for any $r \geq 2$.

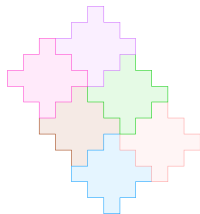


Figure: Tiling of \mathbb{R}^2 by Lee Spheres of radius 2.

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Golomb-Welch Conjecture

Golomb-Welch Conj. (1969): For any $n \geq 3$ there is no tilings of \mathbb{R}^n by clusters of unit cubes associated with Lee spheres of radius $r \geq 2$.

Equivalently, there are no $PL(n, r)$ codes for $n \geq 3$ and $r \geq 2$.

Gravier, Molard, Payan (1998): There are no $PL(3, r)$ codes for any $r \geq 2$.

Spacapan (2007): There are no $PL(4, r)$ codes for any $r \geq 2$.

P. Horak (2009): There are no $PL(5, r)$ codes for any $r \geq 2$;
There are no $PL(6, 2)$ codes.
There are no $PL(6, r)$ codes for any $r \geq 2$.

O. Grosek & P. Horak (2014): There are no **lattice** tilings of \mathbb{R}^n for $7 \leq n \leq 12$ by Lee spheres of radius 2.

What about $PL(7, 2)$ codes?

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$PL(n, 2)$ codes

Let \mathcal{M} be a $PL(n, 2)$ code and that $O = (0, 0, \dots, 0) \in \mathcal{M}$, which means that:

- all words $V \in \mathbb{Z}^n$ such that $\rho_L(V, O) \leq 2$ are covered by O .

Next level of words to be covered $\Rightarrow \mathcal{V}_3 = \{V \in \mathbb{Z}^n : \rho_L(O, V) = 3\}$.

Let $\mathcal{T} = \{W \in \mathcal{M} \text{ covering the words of } \mathcal{V}_3\}$. Then, $\mathcal{T} \subset \mathcal{V}_5$,

$$\mathcal{T} = \mathcal{A} \cup \mathcal{B} \cup \mathcal{C} \cup \mathcal{D} \cup \mathcal{E} \cup \mathcal{F} \cup \mathcal{G} \text{ where}$$

- ▶ $\mathcal{A} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 5]\}$;
- ▶ $\mathcal{B} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 4, \pm 1]\}$;
- ▶ $\mathcal{C} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 3, \pm 2]\}$;
- ▶ $\mathcal{D} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 3, \pm 1^2]\}$;
- ▶ $\mathcal{E} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 2^2, \pm 1]\}$;
- ▶ $\mathcal{F} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 2, \pm 1^3]\}$;
- ▶ $\mathcal{G} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 1^5]\}$.

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- ▶ $\mathcal{A} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 5]\}; \quad \Rightarrow \quad \mathbf{a} = |\mathcal{A}|$
- ▶ $\mathcal{B} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 4, \pm 1]\}; \quad \Rightarrow \quad \mathbf{b} = |\mathcal{B}|$
- ▶ $\mathcal{C} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 3, \pm 2]\}; \quad \Rightarrow \quad \mathbf{c} = |\mathcal{C}|$
- ▶ $\mathcal{D} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 3, \pm 1^2]\}; \quad \Rightarrow \quad \mathbf{d} = |\mathcal{D}|$
- ▶ $\mathcal{E} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 2^2, \pm 1]\}; \quad \Rightarrow \quad \mathbf{e} = |\mathcal{E}|$
- ▶ $\mathcal{F} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 2, \pm 1^3]\}; \quad \Rightarrow \quad \mathbf{f} = |\mathcal{F}|$
- ▶ $\mathcal{G} = \{V \in \mathcal{T} : V \text{ is of type } [\pm 1^5]\}; \quad \Rightarrow \quad \mathbf{g} = |\mathcal{G}|.$

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$$\mathcal{E} := [\pm 2^2, \pm 1]; \mathcal{F} := [\pm 2, \pm 1^3]; \mathcal{G} := [\pm 1^5]$$

- All words $V \in \mathbb{Z}^n$ such that $\mu_L(V, O) \leq 2$ are covered by O .

Next level of words to be covered $\Rightarrow \mathcal{V}_3 = \{V \in \mathbb{Z}^n : \rho_L(O, V) = 3\}$.

$$\mathcal{V}_3 = \{[\pm 3]; [\pm 2, \pm 1]; [\pm 1^3]\}.$$

- ▶ Each word of type $[\pm 3]$ is covered by one and only one codeword in $\mathcal{A} \cup \mathcal{B} \cup \mathcal{C} \cup \mathcal{D}$ and $|[\pm 3]| = 2n$. Thus,

$$a + b + c + d = 2n$$

- ▶ Each word of type $[\pm 2, \pm 1]$ is covered by one codeword in \mathcal{B} ; 2 codewords in \mathcal{C} ; 2 codewords in \mathcal{D} ; 4 codewords in \mathcal{E} and 3 codewords in \mathcal{F} . Besides, $|[\pm 2, \pm 1]| = 8 \binom{n}{2}$. Thus,

$$b + 2c + 2d + 4e + 3f = 8 \binom{n}{2}$$

- ▶ Each word of type $[\pm 1^3]$ is covered by one codeword of \mathcal{D} ; one codeword of \mathcal{E} ; 4 codewords of \mathcal{F} and 10 words of \mathcal{G} . Besides, $|[\pm 1^3]| = 8 \binom{n}{3}$. Thus,

$$d + e + 4f + 10g = 8 \binom{n}{3}.$$

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Set of signed coordinates $I = \{+1, +2, \dots, +n, -1, -2, \dots, -n\}$

Given $\mathcal{H} \subset \mathbb{Z}^n$ denote by

- ▶ $\mathcal{H}_i = \{W \in \mathcal{H} : iw_{|i|} > 0\}, i \in I$
- ▶ $\mathcal{H}_{ij} = \{W \in \mathcal{H} : iw_{|i|} > 0 \wedge jw_{|j|} > 0\}, i, j \in I, i \neq j, i \neq -j$
- ▶ $\mathcal{H}_i^{(k)} = \{W \in \mathcal{H} : iw_{|i|} > 0 \wedge |w_{|i|}| = k\}, i \in I, k \in \mathbb{Z}^+.$

Relation between the cardinality of each set of codewords and their index subsets can be derived:

$$\text{▶ } g = |\mathcal{G}| = \frac{1}{5} \sum_{i \in I} |\mathcal{G}_i|; |\mathcal{G}_i| = \frac{1}{4} \sum_{j \in I \setminus \{i, -i\}} |\mathcal{G}_{ij}|; |\mathcal{G}_{ij}| = \frac{1}{3} \sum_{k \in I \setminus \{i, -i, j, -j\}} |\mathcal{G}_{ijk}|.$$

Similar equalities for the other subsets of \mathcal{T} can be derived. ◻

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Proposition:

- ▶ $|\mathcal{A}_i \cup \mathcal{B}_i^{(4)} \cup \mathcal{C}_i^{(3)} \cup \mathcal{D}_i^{(3)}| = 1$, for each $i \in I$;
- ▶ $|\mathcal{B}_i^{(4)} \cap \mathcal{B}_j^{(1)}| + |\mathcal{C}_i \cap \mathcal{C}_j| + |\mathcal{D}_i^{(3)} \cap \mathcal{D}_j^{(1)}| + |\mathcal{E}_i^{(2)} \cap \mathcal{E}_j| + |\mathcal{F}_i^{(2)} \cap \mathcal{F}_j^{(1)}| = 1$, for each $i, j \in I$, with $|i| \neq |j|$
- ▶ $|\mathcal{D}_{ijk} \cup \mathcal{E}_{ijk} \cup \mathcal{F}_{ijk} \cup \mathcal{G}_{ijk}| = 1$, for each $i, j, k \in I$, with $|i|, |j|$ and $|k|$ distinct between them.
- ▶ $\forall i \in I, |\mathcal{B}_i^{(4)} \cup \mathcal{C}_i^{(2)} \cup \mathcal{C}_i^{(3)}| + 2|\mathcal{D}_i^{(3)} \cup \mathcal{E}_i^{(2)}| + 3|\mathcal{F}_i^{(2)}| = 2(n-1)$ and
 $|\mathcal{B}_i^{(1)} \cup \mathcal{C}_i^{(2)} \cup \mathcal{C}_i^{(3)} \cup \mathcal{D}_i^{(1)} \cup \mathcal{E}_i^{(2)} \cup \mathcal{F}_i^{(1)}| + 2|\mathcal{E}_i^{(1)}| = 2(n-1)$.

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- ▶ $|\mathcal{A}_i \cup \mathcal{B}_i^{(4)} \cup \mathcal{C}_i^{(3)} \cup \mathcal{D}_i^{(3)}| = 1$, for each $i \in I$;
- ▶ $|\mathcal{B}_i^{(4)} \cap \mathcal{B}_j^{(1)}| + |\mathcal{C}_i \cap \mathcal{C}_j| + |\mathcal{D}_i^{(3)} \cap \mathcal{D}_j^{(1)}| + |\mathcal{E}_i^{(2)} \cap \mathcal{E}_j| + |\mathcal{F}_i^{(2)} \cap \mathcal{F}_j^{(1)}| = 1$, for each $i, j \in I$, with $|i| \neq |j|$
- ▶ $|\mathcal{D}_{ijk} \cup \mathcal{E}_{ijk} \cup \mathcal{F}_{ijk} \cup \mathcal{G}_{ijk}| = 1$, for each $i, j, k \in I$, with $|i|, |j|$ and $|k|$ distinct between them.

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Proof: Let V be a word of type $[\pm 2, \pm 1]$, satisfying $i\nu_{|i}, j\nu_{|j} > 0$, $|\nu_{|i}| = 2$ and $|\nu_{|j}| = 1$.

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V must be covered by a codeword $W \in \mathcal{B} \cup \mathcal{C} \cup \mathcal{D} \cup \mathcal{E} \cup \mathcal{F}$ satisfying one and only one of the following conditions:

$$W \in \mathcal{B}_i^{(4)} \cap \mathcal{B}_j^{(1)}; W \in \mathcal{C}_i \cap \mathcal{C}_j; W \in \mathcal{D}_i^{(3)} \cap \mathcal{D}_j^{(1)}; W \in \mathcal{E}_i^{(2)} \cap \mathcal{E}_j; W \in \mathcal{F}_i^{(2)} \cap \mathcal{F}_j^{(1)}.$$

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Proposition:

- ▶ $|D_i \cup E_i| + 3|F_i| + 6|G_i| = 4\binom{n-1}{2}$,
- ▶ $|\mathcal{D}_{ij} \cup \mathcal{E}_{ij}| + 2|\mathcal{F}_{ij}| + 3|\mathcal{G}_{ij}| = 2(n-2); \quad |\mathcal{D}_i \cup \mathcal{E}_i| \leq 2n-1$
- ▶ $|G_i| > \frac{|D_i \cup E_i| + (n-1)(n-6)}{3} - \frac{1}{6}, \quad i \in I;$
- ▶ $|F_i| \leq \frac{8(n-1)+1}{3} - |\mathcal{D}_i \cup \mathcal{E}_i| - \frac{2}{3}|\mathcal{E}_i|, \quad i \in I;$
- ▶ If $n \equiv 1 \pmod{3}$, then $|G_i| \leq \frac{(n-1)(2n-5)}{6}$ for each $i \in I$.
- ▶ If $n \equiv 0 \pmod{3}$, then $|G_i| \leq \frac{(n-1)(n-3)}{3}$ for each $i \in I$.

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Corollary:

- ▶ If $n \equiv 0 \pmod{3}$, then $g \leq \frac{2n(n-1)(n-3)}{15}$;
- ▶ If $n \equiv 1 \pmod{3}$, then $g \leq \frac{n(n-1)(2n-5)}{15}$.

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► "Good Solutions" of
$$\begin{cases} a + b + c + d = 14 \\ b + 2c + 2d + 4e + 3f = 168 \\ d + e + 4f + 10g = 280. \end{cases}$$

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- ▶ If $n \equiv 1 \pmod{3}$, then $|G_i| \leq \frac{(n-1)(2n-5)}{6}$ for each $i \in I.$
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Proposition:

- ▶ $|D_i \cup E_i| + 3|F_i| + 6|G_i| = 60,$
- ▶ $|\mathcal{D}_{ij} \cup \mathcal{E}_{ij}| + 2|\mathcal{F}_{ij}| + 3|\mathcal{G}_{ij}| = 10; |\mathcal{D}_{ij} \cup \mathcal{E}_{ij}|$ and $|\mathcal{G}_{ij}|$ have the same parity,
 $|D_i \cup E_i| \leq 13$
- ▶ $3 \leq |G_i| \leq 9$ for each $i \in I.$
- ▶ $|\mathcal{F}_i| \leq \frac{49}{3} - |D_i \cup E_i| - \frac{2}{3}|E_i|, i \in I;$
- ▶ $6 \leq g \leq 25.$
- ▶ $\forall i \in I, |\mathcal{B}_i^{(4)} \cup \mathcal{C}_i^{(2)} \cup \mathcal{C}_i^{(3)}| + 2|\mathcal{D}_i^{(3)} \cup \mathcal{E}_i^{(2)}| + 3|\mathcal{F}_i^{(2)}| = 12$ and
 $|\mathcal{B}_i^{(1)} \cup \mathcal{C}_i^{(2)} \cup \mathcal{C}_i^{(3)} \cup \mathcal{D}_i^{(1)} \cup \mathcal{E}_i^{(2)} \cup \mathcal{F}_i^{(1)}| + 2|\mathcal{E}_i^{(1)}| = 12.$

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For each $i \in I$, $3 \leq |\mathcal{G}_i| \leq 8$; $9 \leq g \leq 22$.

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$$\triangleright |\mathcal{G}_i| = \frac{1}{4} \sum_{j \in I \setminus \{i, -i\}} |\mathcal{G}_{ij}| \implies \sum_{j \in I \setminus \{i, -i\}} |\mathcal{G}_{ij}| = 36.$$

$$\triangleright |\mathcal{G}_{ij}| = 3 \text{ for any } j \in I \setminus \{i, -i\}.$$

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Let $W_1 \in \mathcal{G}_{i\alpha\beta\gamma\delta}$, $\alpha, \beta, \gamma, \delta \in I \setminus \{i, -i\}$ and $|\alpha|, |\beta|, |\gamma|, |\delta|$ pairwise distinct.

\triangleright For any $W \in \mathcal{G}_i \setminus \{W_1\}$ there exists, at most, one element $\varepsilon \in \{\alpha, \beta, \gamma, \delta\}$ so that $W \in \mathcal{G}_{i\varepsilon}$.

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W_1	i	α	β	γ	δ
W_2	i	α			
W_3	i	α			
W_4	i	β			
W_5	i	β			
W_6	i	γ			
W_7	i	γ			
W_8	i	δ			
W_9	i	δ			

Partial index distribution of the codewords of \mathcal{G}_i



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For each $i \in I$, $3 \leq |G_i| \leq 8$; $9 \leq g \leq 22$.

Proof:

Let $\mathcal{J} = \{\alpha, \beta, \gamma, \delta\}$, $\mathcal{J}^- = \{-\alpha, -\beta, -\gamma, -\delta\}$, $\mathcal{X} = I \setminus (\{i, -i\} \cup \mathcal{J} \cup \mathcal{J}^-) = \{x, -x, y, -y\}$.

W_1	i	α	β	γ	δ
W_2	i	α	k_1	k_2	
W_3	i	α	k_3		
W_4	i	β	k_4	k_5	
W_5	i	β	k_6		
W_6	i	γ	k_7	k_8	
W_7	i	γ	k_9		
W_8	i	δ	k_{10}	k_{11}	
W_9	i	δ	k_{12}		

For each $\varepsilon \in \mathcal{J}$ and $W, W' \in G_{i\varepsilon} \setminus \{W_1\}$,
 there are, **at least**, three distinct elements
 $k, k', k'' \in \mathcal{X}$ so that
 $W \in G_{i\varepsilon k k'}$ and $W' \in G_{i\varepsilon k''}$.

Figure: Partial index distribution of the codewords of G_i .

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W_3	i	α	k_3			
W_4	i	β	k_4	k_5		
W_5	i	β	k_6			
W_6	i	γ	k_7	k_8		
W_7	i	γ	k_9			
W_8	i	δ	k_{10}	k_{11}		
W_9	i	δ	k_{12}			

W_1	i	α	β	γ	δ	
W_2	i	α	x	y	j_1	
W_3	i	α	$-x$	j_2	j_3	
W_4	i	β	$-x$	y	j_4	
W_5	i	β	k_6			
W_6	i	γ	k_7	k_8		
W_7	i	γ	k_9			
W_8	i	δ	k_{10}	k_{11}		
W_9	i	δ	k_{12}			

$j_1, j_2, j_3 \in \mathcal{J}^-$
 $j_4 \neq j_1, j_2, j_3$

Figure: Partial index distribution of the codewords of G_i .

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$$\mathcal{A} := [\pm 5]; \mathcal{B} := [\pm 4, \pm 1]; \mathcal{C} := [\pm 3, \pm 2]; \mathcal{D} := [\pm 3, \pm 1^2];$$

$$\mathcal{E} := [\pm 2^2, \pm 1]; \mathcal{F} := [\pm 2, \pm 1^3]; \mathcal{G} := [\pm 1^5]$$

Proposition:

For each $i \in I$, $3 \leq |G_i| \leq 8$; $9 \leq g \leq 22$.

Proof:

Let $\mathcal{J} = \{\alpha, \beta, \gamma, \delta\}$, $\mathcal{J}^- = \{-\alpha, -\beta, -\gamma, -\delta\}$, $\mathcal{K} = I \setminus (\{i, -i\} \cup \mathcal{J} \cup \mathcal{J}^-) = \{x, -x, y, -y\}$.

W_1	i	α	β	γ	δ
W_2	i	α	x	y	j_1
W_3	i	α	$-x$	j_2	j_3
W_4	i	β	$-x$	y	j_4
W_5	i	β	k_6		
W_6	i	γ	k_7	k_8	
W_7	i	γ	k_9		
W_8	i	δ	k_{10}	k_{11}	
W_9	i	δ	k_{12}		

$$j_1, j_2, j_3 \in \mathcal{J}^-$$

$$j_4 \neq j_1, j_2, j_3,$$

$$y \neq k_6$$

$$\text{If } y \in \{k_7, k_8, k_{10}, k_{11}\}$$



$$W_2 \in \mathcal{G}_{ixy} \text{ and } W_4 \in \mathcal{G}_{i,-x,y}$$

contradiction

$$y = k_9 \text{ or } y = k_{12}$$

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W_3	i	α	$-x$	j_2	j_3
W_4	i	β	$-x$	y	j_4
W_5	i	β	k_6		
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W_7	i	γ	k_9	j_5	j_6
W_8	i	δ	k_{10}	k_{11}	
W_9	i	δ	k_{12}		

$j_1, j_2, j_3 \in \mathcal{J}^-$
 $j_4 \neq j_1, j_2, j_3,$

$y = k_9$ or $y = k_{12}$

Then $W_7 \in \mathcal{G}_{i\gamma j_5 j_6}$ $j_5, j_6 \neq j_1, j_4$

$j_5 = j_2$ and $j_6 = j_3.$

contradiction

Figure: Partial index distribution of the codewords of \mathcal{G}_i .

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$j_1, j_2, j_3 \in \mathcal{J}^-$
 $j_4 \neq j_1, j_2, j_3,$

$y = k_9$ or $y = k_{12}$

Then $W_7 \in G_{i\gamma y j_5 j_6}$ $j_5, j_6 \neq j_1, j_4$

$j_5 = j_2$ and $j_6 = j_3.$

contradiction

$y = k_{12}$ \Rightarrow contradiction

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Lemma

$\forall i \in I, |\mathcal{F}_i^{(2)}| \leq 4$. If $|\mathcal{F}_i^{(2)}| = 4$, then $|\mathcal{F}_i^{(2)} \cap \mathcal{F}_j| = 1$ for all $j \in I \setminus \{i, -i\}$

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Lemma

- ▶ $|G_i| = 3 \implies |\mathcal{A}_i| = 1, |\mathcal{B}_i \cup \mathcal{C}_i \cup \mathcal{E}_i| = 0, |\mathcal{D}_i| = |\mathcal{D}_i^{(1)}| = 3$ ($|\mathcal{D}_i^{(3)}| = 0$);
 $|\mathcal{F}_i| = 13$ with $|\mathcal{F}_i^{(2)}| = 4$ and $|\mathcal{F}_i^{(1)}| = 9$.
- ▶ $|G_i| = 4 \implies |\mathcal{D}_i \cup \mathcal{E}_i| = 3$ and $|\mathcal{F}_i| = 11$; *or*
 $|\mathcal{D}_i| = 6 = |\mathcal{D}_i^{(1)}|, |\mathcal{E}_i| = 0, |\mathcal{F}_i| = 10, |\mathcal{A}_i| = 1, |\mathcal{B}_i \cup \mathcal{C}_i| = 0,$
 $|\mathcal{F}_i^{(2)}| = 4$ and $|\mathcal{F}_i^{(1)}| = 6$.
- ▶ $|G_i| = 5 \implies |\mathcal{D}_i \cup \mathcal{E}_i| = 0$ and $|\mathcal{F}_i| = 10$ *or* $|\mathcal{D}_i \cup \mathcal{E}_i| = 3$ and $|\mathcal{F}_i| = 9$ *or*
 $|\mathcal{D}_i \cup \mathcal{E}_i| = 6, |\mathcal{F}_i| = 8$ and $|\mathcal{D}_i| \geq 3$ *or*
 $|\mathcal{D}_i| = 9 = |\mathcal{D}_i^{(1)}|, |\mathcal{A}_i| = 1, |\mathcal{B}_i \cup \mathcal{C}_i| = |\mathcal{E}_i| = 0$ and $|\mathcal{F}_i| = 7$.
with $|\mathcal{F}_i^{(2)}| = 4$ and $|\mathcal{F}_i^{(1)}| = 3$.

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Conditions satisfied by the index subsets of $\mathcal{A} \cup \mathcal{B} \cup \mathcal{C} \cup \mathcal{D} \cup \mathcal{E} \cup \mathcal{F}$ for a specific value of $|\hat{G}_i|$.

Lemma

- ▶ $|\hat{G}_i| = 6 \implies |\mathcal{D}_i \cup \mathcal{E}_i| = 0, |\mathcal{F}_i| = 8$ *or*; $|\mathcal{D}_i \cup \mathcal{E}_i| = 3, |\mathcal{F}_i| = 7;$ *or*
 $|\mathcal{D}_i \cup \mathcal{E}_i| = 6, |\mathcal{F}_i| = 6$ *or* $|\mathcal{D}_i \cup \mathcal{E}_i| = 9, |\mathcal{D}_i| \geq 6, |\mathcal{F}_i| = 5$ *or*
 $|\mathcal{A}_i| = 1, |\mathcal{B}_i \cup \mathcal{C}_i \cup \mathcal{E}_i| = 0, |\mathcal{D}_i| = |\mathcal{D}_i^{(1)}| = 12, |\mathcal{F}_i| = |\mathcal{F}_i^{(2)}| = 4$
- ▶ $|\hat{G}_i| = 7 \implies |\mathcal{D}_i \cup \mathcal{E}_i| = 3, |\mathcal{F}_i| = 5$ *or* $|\mathcal{D}_i \cup \mathcal{E}_i| = 6, |\mathcal{F}_i| = 4;$ *or*
 $|\mathcal{D}_i \cup \mathcal{E}_i| = 9, |\mathcal{D}_i| \geq 3, |\mathcal{F}_i| = 63$ *or*
 $|\mathcal{D}_i \cup \mathcal{E}_i| = 12, |\mathcal{D}_i| \geq 9, |\mathcal{F}_i| = 2.$

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Proposition:

For each $i \in I$, $3 \leq |\mathcal{G}_i| \leq 7$; $9 \leq g \leq 19$.

Proof: By contradiction assume $\exists_{i \in I} |\mathcal{G}_i| = 8 \Rightarrow \sum_{j \in I \setminus \{i, -i\}} |\mathcal{G}_{ij}| = 32$.

Recall that $\forall_{j \in I \setminus \{i, -i\}} |\mathcal{G}_{ij}| \leq 3$.

Denote by $\mathcal{J} = \{j \in I \setminus \{i, -i\} : |\mathcal{G}_{ij}| = 3\} (\Rightarrow |\mathcal{J}| \geq 8)$, by

$\mathcal{L} = \{j \in I \setminus \{i, -i\} : |\mathcal{G}_{ij}| \leq 2\} (\Rightarrow |\mathcal{L}| \leq 4)$ and by

$\mathcal{K} = \{j \in I \setminus \{i, -i\} : |\mathcal{G}_{ij}| = 2\}$.

- ▶ There are, at most, four distinct elements $j \in \mathcal{J}$ such that $-j \in \mathcal{L}$
- ▶ There are $x, y \in \mathcal{J}$, distinct between them, so that $-x, -y \in \mathcal{J}$; i. e. $|\mathcal{G}_{ix}| = |\mathcal{G}_{i,-x}| = |\mathcal{G}_{iy}| = |\mathcal{G}_{i,-y}| = 3$.

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Denote by $\mathcal{J} = \{j \in I \setminus \{i, -i\} : |\mathcal{G}_{ij}| = 3\} (\Rightarrow |\mathcal{J}| \geq 8)$, by
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 $\mathcal{K} = \{j \in I \setminus \{i, -i\} : |\mathcal{G}_{ij}| = 2\}.$

$$|\mathcal{G}_{ix}| = |\mathcal{G}_{i,-x}| = |\mathcal{G}_{iy}| = |\mathcal{G}_{i,-y}| = 3.$$

Step 1: $|\mathcal{J}| = 8$; $|\mathcal{K}| = |\mathcal{L}| = 4$ and the partial index distribution of the codewords $W_1, \dots, W_8 \in \mathcal{G}_i$ satisfies:

W_1	i	k_1	x	y
W_2	i	k_2	x	$-y$
W_3	i	k_3	x	
W_4	i	k_4	$-x$	y
W_5	i	k_5	$-x$	$-y$
W_6	i	k_6	$-x$	
W_7	i	k_7	y	
W_8	i	k_8	$-y$	

where $x, -x, y, -y \in \mathcal{J}$,

$x \neq y$ and

$k_1, \dots, k_8 \in \mathcal{K}.$

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Denote by $\mathcal{J} = \{j \in I \setminus \{i, -i\} : |\mathcal{G}_{ij}| = 3\} (\Rightarrow |\mathcal{J}| \geq 8)$, and by

$$\mathcal{K} = \{j \in I \setminus \{i, -i\} : |\mathcal{G}_{ij}| = 2\}.$$

$$|\mathcal{G}_{ix}| = |\mathcal{G}_{i,-x}| = |\mathcal{G}_{iy}| = |\mathcal{G}_{i,-y}| = 3.$$

Step 2: If $k \in \mathcal{K}$, then $-k \in \mathcal{K}$.

W_1	i	k_1	x	y	α
W_2	i	k_2	x	$-y$	β
W_3	i	k_3	x	γ	δ
W_4	i	k_2	$-x$	y	
W_5	i	k_1	$-x$	$-y$	
W_6	i	k_3	$-x$	α	β
W_7	i	k_7	y	β	
W_8	i	k_7	$-y$	α	

$$\mathcal{K} = \{k_1, k_2, k_3, k_7\}$$

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Step 2: If $k \in \mathcal{K}$, then $-k \in \mathcal{K}$.

W_1	i	k_1	x	y	α
W_2	i	k_2	x	$-y$	β
W_3	i	k_3	x	γ	δ
W_4	i	k_2	$-x$	y	γ
W_5	i	k_1	$-x$	$-y$	δ
W_6	i	k_3	$-x$	α	β
W_7	i	k_7	y	β	δ
W_8	i	k_7	$-y$	α	γ

$$\mathcal{K} = \{k_1, k_2, k_3, k_7\}$$

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Denote by $\mathcal{J} = \{j \in I \setminus \{i, -i\} : |\mathcal{G}_{ij}| = 3\} (\Rightarrow |\mathcal{J}| \geq 8)$, and by

$$\mathcal{K} = \{j \in I \setminus \{i, -i\} : |\mathcal{G}_{ij}| = 2\}.$$

Step 3: $|\mathcal{F}_i| = 0$.

Step 4: $\forall j \in \mathcal{J} \quad |\mathcal{D}_{ij} \cup \mathcal{E}_{ij}| = 1$; $\forall k \in \mathcal{K} \quad |\mathcal{D}_{ik} \cup \mathcal{E}_{ik}| = 4$. Besides, if $k \in \mathcal{K}$ there are codewords $V_1, V_2 \in \mathcal{G}_{ik}$ and $U_1, \dots, U_4 \in \mathcal{D}_{ik} \cup \mathcal{E}_{ik}$ whose index distribution satisfies:

V_1	i	k	v_1	v_2	v_3
V_2	i	k	v_4	v_5	v_6

$$V_1, V_2 \in \mathcal{G}_{ik}$$

$$\{v_1, \dots, v_6\} \subset \mathcal{J}$$

U_1	i	k	u
U_2	i	k	$-u$
U_3	i	k	j_1
U_4	i	k	j_2

$$U_1, \dots, U_4 \in \mathcal{D}_{ik} \cup \mathcal{E}_{ik}.$$

$$u, -u \in \mathcal{K} \setminus \{k, -k\}$$

$$j_1, j_2 \in \mathcal{J}, j_1 \neq j_2$$

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Denote by $\mathcal{J} = \{j \in I \setminus \{i, -i\} : |\mathcal{G}_{ij}| = 3\} (\Rightarrow |j| \geq 8)$, and by
 $\mathcal{K} = \{j \in I \setminus \{i, -i\} : |\mathcal{G}_{ij}| = 2\}$.

Step 5: $|\mathcal{G}_i| \neq 8$.

Recall $[\pm 2, \pm 1]$ are covered by codewords of $\mathcal{B}, \mathcal{C}, \mathcal{D}, \mathcal{E}, \mathcal{F}$.

$$V_1, V_2 \in \mathcal{G}_{ik}$$

$$\{v_1, \dots, v_6\} \subset \mathcal{J}$$

V_1	i	k	v_1	v_2	v_3
V_2	i	k	v_4	v_5	v_6

U_1	i	k	u
U_2	i	k	$-u$
U_3	i	k	j_1
U_4	i	k	j_2

$$U_1, \dots, U_4 \in \mathcal{D}_{ik} \cup \mathcal{E}_{ik}$$

$$u, -u \in \mathcal{K} \setminus \{k, -k\}$$

$$j_1, j_2 \in \mathcal{J}, j_1 \neq j_2$$

$\mathcal{H} \Leftrightarrow$ set of words of type $[\pm 2, \pm 1]$

$$P_1 \in \mathcal{H}_i^{(2)} \cap \mathcal{H}_{j_1}^{(1)}$$

$$P_2 \in \mathcal{H}_i^{(2)} \cap \mathcal{H}_{j_2}^{(1)}$$

Recall $[\pm 2, \pm 1]$ are covered by
 codewords of $\mathcal{B}, \mathcal{C}, \mathcal{D}, \mathcal{E}, \mathcal{F}$

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Step 5: $|\hat{G}_i| \neq 8$.

	i	k	u	$-u$	j_1	j_2
U_1	x	x	x			
U_2	x	x		x		
U_3	x	x			x	
U_4	x	x				x
P_1	± 2				± 1	
P_2	± 2					± 1

Index distribution of $U_1, \dots, U_4 \in \mathcal{D}_{ik} \cup \mathcal{E}_{ik}$

and index value distribution of P_1 and P_2

- ▶ P_1 may only be covered by $U_3 \in (\mathcal{B}_i^{(4)} \cap \mathcal{B}_{j_1}^{(1)}) \cup \mathcal{C}_{j_1} \cup (\mathcal{D}_i^{(3)} \cap \mathcal{D}_{j_1}^{(1)}) \cup (\mathcal{E}_i^{(2)} \cap \mathcal{E}_{j_1})$;
- ▶ P_2 may only be covered by $U_4 \in (\mathcal{B}_i^{(4)} \cap \mathcal{B}_{j_2}^{(1)}) \cup \mathcal{C}_{j_2} \cup (\mathcal{D}_i^{(3)} \cap \mathcal{D}_{j_2}^{(1)}) \cup (\mathcal{E}_i^{(2)} \cap \mathcal{E}_{j_2})$

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- 5.1 $3 \leq |\mathcal{G}_i| \leq 8$
- 5.2 $3 \leq |\hat{\mathcal{G}}_i| \leq 7$
- 5.3 $4 \leq |\mathcal{G}_i| \leq 7$

$$\mathcal{A} := [\pm 5]; \mathcal{B} := [\pm 4, \pm 1]; \mathcal{C} := [\pm 3, \pm 2]; \mathcal{D} := [\pm 3, \pm 1^2];$$

$$\mathcal{E} := [\pm 2^2, \pm 1]; \mathcal{F} := [\pm 2, \pm 1^3]; \mathcal{G} := [\pm 1^5]$$

Step 5: $|\mathcal{G}_i| \neq 8$.

	i	k	u	$-u$	j_1	j_2
U_1	x	x	x			
U_2	x	x		x		
U_3	x	x			x	
U_4	x	x				x
P_1	± 2				± 1	
P_2	± 2					± 1

Index distribution of $U_1, \dots, U_4 \in \mathcal{D}_{ik} \cup \mathcal{E}_{ik}$

and index value distribution of P_1 and P_2

- ▶ P_1 may only be covered by $U_3 \in (\mathcal{B}_i^{(4)} \cap \mathcal{B}_{j_1}^{(1)}) \cup \mathcal{C}_{j_1} \cup (\mathcal{D}_i^{(3)} \cap \mathcal{D}_{j_1}^{(1)}) \cup (\mathcal{E}_i^{(2)} \cap \mathcal{E}_{j_1})$;
- ▶ P_2 may only be covered by $U_4 \in (\mathcal{B}_i^{(4)} \cap \mathcal{B}_{j_2}^{(1)}) \cup \mathcal{C}_{j_2} \cup (\mathcal{D}_i^{(3)} \cap \mathcal{D}_{j_2}^{(1)}) \cup (\mathcal{E}_i^{(2)} \cap \mathcal{E}_{j_2})$

But $|\mathcal{B}_i^{(4)} \cap \mathcal{B}_j^{(1)}| + |\mathcal{C}_i \cap \mathcal{C}_j| + |\mathcal{D}_i^{(3)} \cap \mathcal{D}_j^{(1)}| + |\mathcal{E}_i^{(2)} \cap \mathcal{E}_j| + |\mathcal{F}_i^{(2)} \cap \mathcal{F}_j^{(1)}| = 1$ and

$U_3, U_4 \in (\mathcal{D}_i^{(3)} \cap \mathcal{D}_k^{(1)}) \cup (\mathcal{E}_i^{(2)} \cap \mathcal{E}_k) \Rightarrow$ either P_1 is not covered by U_3 or

P_2 is not covered by U_4 . \square

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$$\mathcal{A} := [\pm 5]; \mathcal{B} := [\pm 4, \pm 1]; \mathcal{C} := [\pm 3, \pm 2]; \mathcal{D} := [\pm 3, \pm 1^2]; \\ \mathcal{E} := [\pm 2^2, \pm 1]; \mathcal{F} := [\pm 2, \pm 1^3]; \mathcal{G} := [\pm 1^5]$$

Step 5: $|\mathcal{G}_i| \neq 8$.

- ▶ Assuming WLOG P_1 not covered by $U_3 \Rightarrow P_1$ covered by $V \in (\mathcal{B}_i^{(4)} \cap \mathcal{B}_j^{(1)}) \cup \mathcal{C}_{ij}$.

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Step 5: $|\mathcal{G}_i| \neq 8$.

- ▶ Assuming WLOG P_1 not covered by $U_3 \Rightarrow P_1$ covered by $V \in (\mathcal{B}_i^{(4)} \cap \mathcal{B}_{j_1}^{(1)}) \cup \mathcal{C}_{ij_1}$.

Both cases $V \in (\mathcal{B}_i^{(4)} \cap \mathcal{B}_{j_1}^{(1)})$ and $V \in \mathcal{C}_{ij_1}$ lead to a contradiction.

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Both cases $V \in (\mathcal{B}_i^{(4)} \cap \mathcal{B}_{j_1}^{(1)})$ and $V \in \mathcal{C}_{ij_1}$ lead to a contradiction.

Result: $\forall_{i \in I} 3 \leq |\mathcal{G}_i| \leq 7, 9 \leq g \leq 19$.

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Proposition:

Let $\mathcal{J} = \{i \in I : |\mathcal{G}_i| = p \wedge |\mathcal{F}_i| = q\}$.

- ▶ $p = 3 \Rightarrow q = 13$ and $|\mathcal{J}| \leq 5$;
- ▶ $p = 4 \Rightarrow q = 10$ and $|\mathcal{J}| \leq 3$;
- ▶ $p = 5 \Rightarrow (q = 10$ and $|\mathcal{J}| \leq 4)$ or $(q = 7$ and $|\mathcal{J}| \leq 2)$;
- ▶ $p = 6 \Rightarrow (q = 8$ and $|\mathcal{J}| \leq 3)$ or $(q = 7$ and $|\mathcal{J}| \leq 8)$ or $(q = 5$ and $|\mathcal{J}| \leq 5)$;
- ▶ $p = 7 \Rightarrow (q = 5$ and $|\mathcal{J}| \leq 5)$ or $(q = 2$ and $|\mathcal{J}| \leq 2)$.

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Let $\mathcal{J} = \{i \in I : |\mathcal{G}_i| = p \wedge |\mathcal{F}_i| = q\}$.

- ▶ $p = 3 \Rightarrow q = 13$ and $|\mathcal{J}| \leq 5$;
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- ▶ $p = 5 \Rightarrow (q = 10$ and $|\mathcal{J}| \leq 4)$ or $(q = 7$ and $|\mathcal{J}| \leq 2)$;
- ▶ $p = 6 \Rightarrow (q = 8$ and $|\mathcal{J}| \leq 3)$ or $(q = 7$ and $|\mathcal{J}| \leq 8)$ or $(q = 5$ and $|\mathcal{J}| \leq 5)$;
- ▶ $p = 7 \Rightarrow (q = 5$ and $|\mathcal{J}| \leq 5)$ or $(q = 2$ and $|\mathcal{J}| \leq 2)$.

Corollary:

$g \neq 9$ and $g \neq 10$.

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Let $\mathcal{J} = \{i \in I : |\mathcal{G}_i| = p \wedge |\mathcal{F}_i| = q\}$.

- ▶ $p = 3 \Rightarrow q = 13$ and $|\mathcal{J}| \leq 5$;
- ▶ $p = 4 \Rightarrow q = 10$ and $|\mathcal{J}| \leq 3$;
- ▶ $p = 5 \Rightarrow (q = 10$ and $|\mathcal{J}| \leq 4)$ or $(q = 7$ and $|\mathcal{J}| \leq 2)$;
- ▶ $p = 6 \Rightarrow (q = 8$ and $|\mathcal{J}| \leq 3)$ or $(q = 7$ and $|\mathcal{J}| \leq 8)$ or $(q = 5$ and $|\mathcal{J}| \leq 5)$;
- ▶ $p = 7 \Rightarrow (q = 5$ and $|\mathcal{J}| \leq 5)$ or $(q = 2$ and $|\mathcal{J}| \leq 2)$.

Corollary:

$g \neq 9$ and $g \neq 10$.

$$g = 9 = \frac{1}{5} \sum_{i \in I} |\mathcal{G}_i| \Rightarrow \sum_{i \in I} |\mathcal{G}_i| = 45 = 11 \times 3 + 3 \times 4 \Rightarrow |\mathcal{J}| \geq 11 \text{ (contradiction).}$$

$$g = 10 \Rightarrow \sum_{i \in I} |\mathcal{G}_i| = 50 = 6 \times 3 + 8 \times 4 \Rightarrow |\mathcal{J}| \geq 6 \text{ (contradiction).}$$

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$$\begin{aligned} 5.1 \quad & 3 \leq |\mathcal{G}_i| \leq 8 \\ 5.2 \quad & 3 \leq |\tilde{\mathcal{G}}_i| \leq 7 \\ 5.3 \quad & 4 \leq |\hat{\mathcal{G}}_i| \leq 7 \end{aligned}$$

$$\begin{aligned} \mathcal{A} &:= [\pm 5]; \mathcal{B} := [\pm 4, \pm 1]; \mathcal{C} := [\pm 3, \pm 2]; \mathcal{D} := [\pm 3, \pm 1^2]; \\ \mathcal{E} &:= [\pm 2^2, \pm 1]; \mathcal{F} := [\pm 2, \pm 1^3]; \mathcal{G} := [\pm 1^5] \end{aligned}$$

Proposition:

For each $i \in I$ $|\mathcal{G}_i| \neq 3$. More precisely, $\forall i \in I$ $4 \leq |\mathcal{G}_i| \leq 7$, $12 \leq g \leq 19$.

Proof: By contradiction assume $\exists i \in I$ $|\mathcal{G}_i| = 3 \Rightarrow \sum_{j \in I \setminus \{i, -i\}} |\mathcal{G}_{ij}| = 12$.

Step 1: $|\mathcal{G}_i| = 3 \Rightarrow |\mathcal{D}_i| = 3, |\mathcal{E}_i| = 0, |\mathcal{F}_i| = 13$ and $|\mathcal{J}| = \{i \in I : |\mathcal{G}_i| = 3 \wedge |\mathcal{F}_i| = 13\} \leq 5$.

Step 2: $|\mathcal{G}_i| = 3 \Rightarrow \exists_{\alpha, \beta, \gamma \in I \setminus \{i, -i\}}$ $|\mathcal{F}_{i\alpha}| = |\mathcal{F}_{i\beta}| = |\mathcal{F}_{i\gamma}| = 5$, α, β, γ distinct between them;

$$\forall \omega \in I \setminus \{i, -i, \alpha, \beta, \gamma\} \quad |\mathcal{F}_{i\omega}| \leq 3.$$

$\exists u_1, u_2, u_3, u_4 \in \mathcal{F}_i$ satisfying

U_1	i	α	β	x_1
U_2	i	α	γ	x_2
U_3	i	β	γ	x_3
U_4	i	y_1	y_2	y_3

$$x_1, x_2, x_3, y_1, y_2, y_3 \in \mathcal{I} \setminus \{i, -i, \alpha, \beta, \gamma\}$$

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Proposition:

For each $i \in I$ $|G_i| \neq 3$. More precisely, $\forall i \in I$ $4 \leq |G_i| \leq 7$, $12 \leq g \leq 19$.

Step 3:

$$|\mathcal{F}_{i\alpha}| = |\mathcal{F}_{i\beta}| = |\mathcal{F}_{i\gamma}| = 5 \Rightarrow |G_{i\alpha}| = |G_{i\beta}| = |G_{i\gamma}| = 0.$$

$$\exists \delta, \varepsilon, \theta \in I \setminus \{i, -i, \alpha, \beta, \gamma\} \quad |G_{i\delta}| = |G_{i\varepsilon}| = |G_{i\theta}| = 2$$

$$\text{and } \forall \omega \in I \setminus \{i, -i, \alpha, \beta, \gamma, \delta, \varepsilon, \theta\} \quad |G_{i\omega}| = 1$$

The 3 codewords $W_1, W_2, W_3 \in G_i$ satisfy:

W_1	i	δ	ε	z_1	z_2
W_2	i	δ	θ	z_3	z_4
W_3	i	ε	θ	z_5	z_6

$\delta, \varepsilon, \theta, z_1, z_2, \dots, z_6 \in I \setminus \{i, -i, \alpha, \beta, \gamma\}$
distinct between them

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Step 4: Characterization of the index distributions of $\mathcal{G}_i \cup \mathcal{F}_i$ having into account their interaction.

W_1	i	δ	ε	z_1	z_2
W_2	i	δ	ε	z_3	z_4
W_3	i	ε	δ	z_5	z_6

$W_1, W_2, W_3 \in \mathcal{G}_i$

$\delta, \varepsilon, \theta, z_1, z_2, \dots, z_6 \in \mathcal{I} \setminus \{i, -i, \alpha, \beta, \gamma\}$

$$\delta = -j \text{ and } \varepsilon = -k$$

U_1	i	j	k	l
U_2	i	j	γ	x_1
U_3	i	k	γ	x_2
U_4	i	y_1	y_2	y_3

Codewords of \mathcal{F}_i

I) $\delta, \varepsilon, \theta \neq l$;

II) $\exists^1 \omega \in \{\delta, \varepsilon, \theta\}$ such that $\omega = l$.

1.1

$$\theta = -j$$

W_1	i	δ	ε	l	z_1
W_2	i	δ	$-j$	z_2	z_3
W_3	i	ε	$-j$	z_4	z_5

1.2

$$\theta = -l$$

W_1	i	δ	ε	l	z_1
W_2	i	δ	$-l$	z_2	z_3
W_3	i	ε	$-l$	z_4	z_5

2.1

W_1	i	l	$-j$	m	n
W_2	i	l	$-k$	z_1	z_2
W_3	i	$-j$	$-k$	z_3	z_4

2.2

W_1	i	l	$-j$	z_1	z_2
W_2	i	l	m	z_3	z_4
W_3	i	$-j$	m	z_5	z_6

2.3

W_1	i	l	m	z_1	z_2
W_2	i	l	n	z_3	z_4
W_3	i	m	n	z_5	z_6

Each one of these cases are subdivided in several other according with the available possibilities to fill the unknown coordinates.



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Proposition:

For each $i \in I$ $|\hat{G}_i| \neq 3$. More precisely, $\forall i \in I$ $4 \leq |\hat{G}_i| \leq 7$, $12 \leq g \leq 19$.

Step 4: Characterization of the index distributions of $\hat{G}_i \cup \mathcal{F}_i$ having into account their interaction.

1) $\delta, \varepsilon, \theta \neq l$;

1.1

W_1	i	δ	ε	l	z_1
W_2	i	δ	$-j$	z_2	z_3
W_3	i	ε	$-j$	z_4	z_5

U_1	i	j	k	l
U_2	i	j	y	x_1
U_3	i	k	y	x_2
U_4	i	y_1	y_2	y_3

Codewords of \mathcal{F}_i

1.1.1 $\delta = -k$

W_1	i	$-k$	m	l	n
W_2	i	$-k$	$-j$	z_1	z_2
W_3	i	m	$-j$	z_3	z_4

$$|\hat{G}_{i,-k}| = 2$$

1.1.2.1 $W_1 \in \hat{G}_{i,-k}$

W_1	i	m	n	l	$-k$
W_2	i	m	$-j$	z_1	z_2
W_3	i	n	$-j$	z_3	z_4

$$|\hat{G}_{i,-k}| = 1$$

1.1.2.2 $W_1 \notin \hat{G}_{i,-k}$

W_1	i	m	n	l	o
W_2	i	m	$-j$	$-k$	z_1
W_3	i	n	$-j$	z_2	z_3

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For each $i \in I$ $|G_i| \neq 3$. More precisely, $\forall i \in I \quad 4 \leq |G_i| \leq 7, \quad 12 \leq g \leq 19$.

Step 4: Characterization of the index distributions of $G_i \cup F_i$ having into account their interaction.

G_i					
W_1	i	$-k$	m	l	n
W_2	i	$-k$	$-j$	$-o$	$-n$
W_3	i	m	$-j$	o	$-l$

F_j				
U_1	i	j	k	l
U_2	i	j	$-m$	$-n$
U_5	i	j	m	$-o$
U_6	i	j	$-k$	o
U_7	i	j	$-l$	n

F_{ik}				
U_3	i	k	$-m$	o
U_8	i	k	m	$-n$
U_9	i	k	$-j$	n
U_{10}	i	k	$-l$	$-o$

F_{i-m}				
U_{11}	i	$-m$	$-j$	l
U_{12}	i	$-m$	$-k$	$-l$
U_{13}	i	$-m$	$-o$	n

U_4	i	l	$-n$	o
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Codewords of $F_i \cup G_i$

$$|G_{im}| = 2 = |F_{im}| \quad \text{and} \quad 3 \leq |G_m| \leq 7 \Leftrightarrow |G_m \setminus G_i| \geq 1$$

$W_1 \in G_{i,-k,m,l,n}$, $W_3 \in G_{i,m,-j,-n,-o}$, $U_5 \in F_{i,j,m,-l}$ and $U_8 \in F_{ikmo}$

induce a partition \mathcal{P} of $\mathcal{I}\{i, m, -m\}$:

$$\mathcal{P}_1 = \{-k, l, n\}; \quad \mathcal{P}_2 = \{-j, -n, -o\}; \quad \mathcal{P}_3 = \{j, -l\}; \quad \mathcal{P}_4 = \{k, o\}; \quad \mathcal{P}_5 = \{-i\}$$

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Proposition:

For each $i \in I$ $|G_i| \neq 3$. More precisely, $\forall i \in I$ $4 \leq |G_i| \leq 7$, $12 \leq g \leq 19$.

$$W_1 \in G_{i, -k, m, l, n}, \quad W_3 \in G_{i, m, -j, -n, -o}, \quad U_5 \in F_{i, j, m, -l} \quad \text{and} \quad U_8 \in F_{ikmo}$$

$$\mathcal{P} \text{ of } \mathcal{I} \setminus \{i, m, -m\}: \quad \mathcal{P}_1 = \{-k, l, n\}; \quad \mathcal{P}_2 = \{-j, -n, -o\}; \quad \mathcal{P}_3 = \{j, -l\}; \quad \mathcal{P}_4 = \{k, o\}; \quad \mathcal{P}_5 = \{-i\}$$

Possibilities for $W \in G_m \setminus G_i$

1				$-n$	o
2				$-n$	$-l$
3			$-k$	$-o$	j
4				$-o$	$-l$
5				o	j
6				$-j$	k
7	m	$-i$		$-j$	o
8				$-n$	k
9			l	$-o$	k
10				$-n$	j
11				$-o$	j
12				o	j

13				$-j$	o
14				$-o$	k
15				$-j$	$-l$
16			n	$-o$	$-l$
17				k	j
18				k	$-l$
19				o	j
20				o	$-l$
21			$-j$	k	$-l$
22				k	j
23			$-n$	k	$-l$
24				o	j
25				o	$-l$
26			$-o$	k	j

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4. Known results on $PL(n, 2)$ codes
5. $PL(7, 2)$ codes

- 5.1 $3 \leq |G_i| \leq 8$
- 5.2 $3 \leq |G_i| \leq 7$
- 5.3 $4 \leq |G_i| \leq 7$

$$\mathcal{A} := [\pm 5]; \mathcal{B} := [\pm 4, \pm 1]; \mathcal{C} := [\pm 3, \pm 2]; \mathcal{D} := [\pm 3, \pm 1^2];$$

$$\mathcal{E} := [\pm 2^2, \pm 1]; \mathcal{F} := [\pm 2, \pm 1^3]; \mathcal{G} := [\pm 1^5]$$

Proposition:

For each $i \in I$ $|G_i| \neq 3$. More precisely, $\forall i \in I$ $4 \leq |G_i| \leq 7$, $12 \leq g \leq 19$.

Step 5: Identification of a partition of the index set relevant to the index distributions of $G_i \cup \mathcal{F}_i$.

Possibilities for $M \in \mathcal{F}_m \setminus \mathcal{F}_i$

m	-k	j	n	-j	-j	-i	k	-i	-j			
		-o			-o				-o			
		-n			o				-n			
		j			j				j			
	l	-o		-o	-k		-j		-j	-i	-j	-j
		k		k			o		o		o	
		o		j			j		j		j	
		k		k			-o		-o		-o	
	n	-n		-n	l		o		o	j	o	o
		k		k			j		j		j	
		-n		-n			-o		-o		-o	
		o		-j			-j		-j		-j	
o	-j	-j	j	-i	-i	k	-i	-i				
	-i	-i		o	o		o					
	j	j		j	j		j					
	-o	-o		-o	-o		-o					
j	o	o	-i	o	o	-j	o	o				
	-j	-j		-j	-j		-j					
	-n	-n		-n	-n		-n					
	j	j		j	j		j					

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Proposition:

For each $i \in I$ $|G_i| \neq 3$. More precisely, $\forall i \in I \quad 4 \leq |G_i| \leq 7, \quad 12 \leq g \leq 19$.

Step 6: New results about the index distributions of $\mathcal{G} \cup \mathcal{F}$ having into account the known distributions.

- ▶ $3 \leq |G_m| \leq 6$;
- ▶ if $|G_m| = 3$, then $|F_m| = 13$ with $|F_m^{(2)}| = 4$;
- ▶ if $|G_m| = 4$ and $|F_m| = 10$, then $|F_m^{(2)}| = 4$;
- ▶ if $|G_m| = 5$ and $|F_m| = 7$, then $|F_m^{(2)}| = 4$;
- ▶ if $|G_m| = 6$ and $|F_m| = 4$, then $|F_m^{(2)}| = 4$;
- ▶ If $|F_m^{(2)}| = 4$, then $F_m^{(2)} = \{U_8, M, M', M''\}$, where

$U_8 \in \mathcal{F}_{i,k,m,-n}$ and M, M', M'' satisfy one of the following conditions:

M	$m, j, -k, -l$	
M'	$m, -o, l, -j$	$m, -o, n, -j$
M''	$m, -i, n, o$	$m, -i, l, o$

M	m, j, l, o		
M'	$m, -o, -k, -l$	$m, -o, n, -j$	$m, -o, n, -l$
M''	$m, -i, n, -j$	$m, -i, -k, -l$	$m, -i, -k, -j$

M	m, j, n, o	
M'	$m, -o, -k, -l$	$m, -o, l, -j$
M''	$m, -i, l, -j$	$m, -i, -k, -l$

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$PL(7, 2)$ and the Golomb-Welch Conjecture

Thank You