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Commutative Algebras in Fibonacci Categories

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Let C be a braided monoidal category with braiding denoted $c_{X,Y} : X \otimes Y \rightarrow Y \otimes X$.

If C comes equipped with a natural family of isomorphisms $\theta = \{\theta_X : X \rightarrow X \mid X \in \text{Ob}C\}$ (satisfying suitable axioms) then C is said to be **balanced** and the members of the family θ are referred to as **ribbon twists**.

If every object X of C has a dual object X^* together with unit and counit morphisms then C is said to be **rigid** (or autonomous).

Recall that a tortile (rigid and balanced) monoidal category is said to be **fusion** when it is semi-simple k -linear together with a k -linear tensor product, finite dimensional hom-spaces and a finite number of simple objects (up to isomorphism).

A fusion category is called **modular** when it satisfies a certain non-degeneracy (modularity) condition.

An algebra A in a braided monoidal category C is said to be **commutative** when $\mu_{C_{A,A}} = \mu$.

A commutative algebra A in a balanced category C is called **ribbon** if $\theta_A = 1_A$.

A set R is called a **fusion rule** if its integer span $\mathbb{Z}R$ has the structure of an associative unital ring such that the unit element of $\mathbb{Z}R$ belongs to R and

$$r \cdot s \in \mathbb{Z}_{\geq 0}R$$

for any $r, s \in R$.

We also require R to have a **rigidity condition** which we formulate as follows:

Equip $\mathbb{Z}R$ with the symmetric bilinear form $(-, -)$ defined by $(r, s) = \delta_{r,s}$ for $r, s \in R$. The condition is then defined to be the existence of an involution $(-)^* : R \rightarrow R$ such that

$$(r \cdot s, t) = (s, r^* t) \quad r, s, t \in R$$

Let C be a semi-simple ridged monoidal category. The set $Irr(C)$ of isomorphism classes of simple objects has the structure of a fusion rule and $\mathbb{Z}Irr(C) = K_0(C)$.

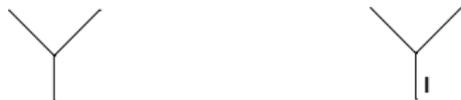
The aim is to describe modular categories with the fusion rule

$$\text{Fib} = \{1, x\} : x^2 = 1 + x$$

We let \mathcal{Fib} be a semi-simple k -linear category with simple objects I and X and tensor product defined as

$$X \otimes X = I \oplus X$$

for which there are two fundamental hom-spaces $\mathcal{Fib}(X^2, X)$ and $\mathcal{Fib}(X^2, I)$. We use a tree / string (of sorts) notation for the two corresponding basis vectors:



We exploit this notational convenience to investigate the categorical properties of \mathcal{Fib} .

The only non-trivial component of the associativity constraint for \mathcal{Fib} is

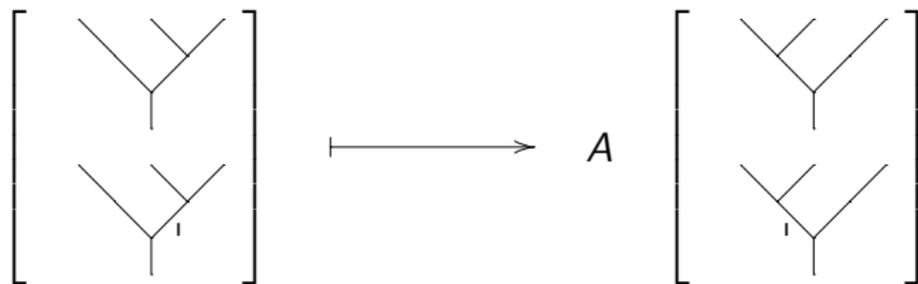
$$\alpha_{X,X,X} : (X \otimes X) \otimes X \rightarrow X \otimes (X \otimes X)$$

which on the level of hom-spaces corresponds the two isomorphisms $\mathcal{Fib}(\alpha_{X,X,X}, I)$ and $\mathcal{Fib}(\alpha_{X,X,X}, X)$.

Clearly $\dim(\mathcal{Fib}(X^3, I)) = 1$ and $\dim(\mathcal{Fib}(X^3, X)) = 2$.

Thus the associativity k -linear transformation for $\mathcal{Fib}(X^3, X)$ is given by $A \in GL_2(k)$ and by $\alpha \in k^*$ for $\mathcal{Fib}(X^3, I)$.

Graphically,

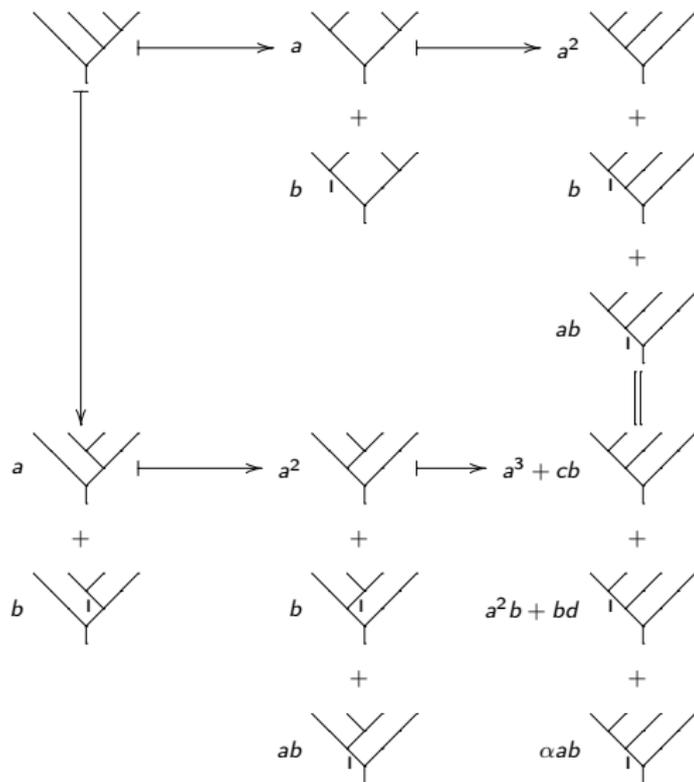


where

$$A = \begin{pmatrix} a & b \\ c & d \end{pmatrix}$$

The pentagon coherence condition for associativity determines a set of equations relating the entries of the matrix A and α .

Clearly $\dim(\mathcal{Fib}(X^4, I)) = 2$ and $\dim(\mathcal{Fib}(X^4, X)) = 3$ such that there are five distinct graphical calculations to perform.



$$\alpha a^2 + bc = \alpha^2$$

$$\alpha ab + bd = 0$$

$$\alpha cb + d^2 = 1$$

$$\alpha ca + dc = 0$$

$$a^3 + bc = a^2$$

$$a^2 b + bd = b$$

$$ca^2 + cd = c$$

$$abc + d^2 = 0$$

$$\alpha ab = ab$$

$$\alpha cb = d$$

$$\alpha ca = ca$$

$$\alpha^2 d = cb$$

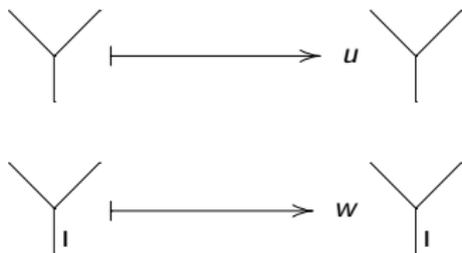
To summarize these calculations, the associativity constraint for \mathcal{Fib} is given by,

$$\alpha = 1, \quad A = \begin{pmatrix} a & b \\ -ab^{-1} & -a \end{pmatrix}$$

where a is a solution of $a^2 = a + 1$.

We note that $\det(A) = -1$ and $A^{-1} = A$.

Similar calculations are performed to obtain a classification of possible braidings and corresponding balanced structures. Braidings come from first considering the form these categorical operations take on the level of fundamental hom-spaces and then running through the axioms.



where $w, u \in k^*$.

It turns out that braidings are completely determined by a number u (where $w = u^2$) satisfying the relation $u^2 = ua - 1$. Together with $a^2 = 1 + a$ these two equations tell us that u is a primitive root of unity of order 10.

Thus the domain of definition for a braided Fibonacci category is the cyclotomic field $\mathbb{Q}(\sqrt[10]{1})$.

Balanced structures require us to consider the endo-hom-spaces of the two simple objects I and X . Since these spaces are one dimensional (basis vectors are the respective identities) the ribbon twists are simply scalar multiples of the basis vectors.

We find that $\theta_1 = id_1$ and $\theta_X = \rho \cdot id_X$ where naturality has forced the scalar multiple corresponding to θ_1 to be unity. The ensuing calculations tells us that balanced structures are completely determined by the braiding as $\rho = u^{-2}$.

By studying possible monoidal equivalences one finds that, up to monoidal equivalence, associativity constraints for Fibonacci categories correspond to solutions of $a^2 = 1 + a$, such that

$$\alpha = 1, \quad A = \begin{pmatrix} a & 1 \\ -a & -a \end{pmatrix}$$

Further more, for each associativity constraint we have that braided balanced structures (up to braided equivalence) correspond to solutions of $u^2 = au - 1$.

The above, taken together with studying rigid structures, yields the following.

Theorem

Every braided balanced structure on a Fibonacci category is modular. Thus there are four non-equivalent Fibonacci modular categories \mathcal{Fib}_u , parameterized by primitive roots of unity u of order 10.

A set M is a **Non-negative Integer Matrix (NIM)** representation of a fusion rule R if $\mathbb{Z}M$ is equipped with the structure of a $\mathbb{Z}R$ -module such that

$$r \cdot m \in \mathbb{Z}_{\geq 0}M$$

where $r \in R$ and $m \in M$.

Just as for fusion rules the rigidity condition is imposed on $\mathbb{Z}M$.

Let \mathcal{M} be a semi-simple module category (actegory) over a semi-simple rigid monoidal category C . Then $Irr(\mathcal{M})$ is a NIM-representation of the fusion rule $Irr(C)$.

The endgame plan is to study ribbon commutative algebras in $\mathcal{Fib}^{\boxtimes \ell}$ (the tensor powers of \mathcal{Fib}). We do this by classifying the NIM-representations of the Fibonacci fusion rule \mathcal{Fib} and its tensor powers $\mathcal{Fib}^{\times \ell}$.

It is in this way that we avoid having to directly classify all possible module categories of $\mathcal{Fib}^{\boxtimes \ell}$.

We choose to encode NIM-representations of $\text{Fib}^{\times \ell}$ as a certain type of oriented graph.

Nodes correspond to elements of a NIM-set M . Edges are colored in ℓ colours. Two nodes m and n are the source and the target of an i -th coloured edge respectively iff the multiplicity $(x_i * m, n)$ of n in $x_i * m$ is non-zero. Here $x_i = 1 \otimes \dots \otimes 1 \otimes X \otimes 1 \otimes \dots \otimes 1$, where X is in the i -th component.

It turns out that showing there is only one NIM-graph (and so only one irreducible NIM-representation) for Fib is quite straight forward. The graph is

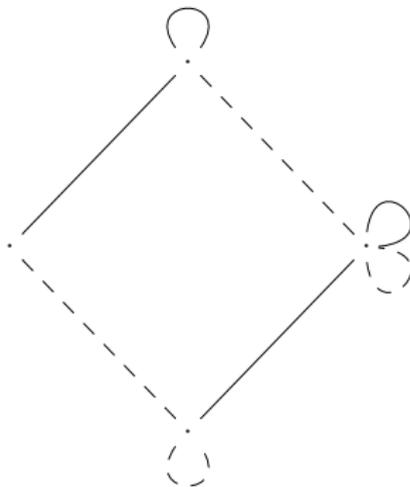


The general case of NIM-representations of $\text{Fib}^{\times \ell}$ is not as easy and requires some fancier footwork.

Theorem

Any indecomposable NIM-representation of $\text{Fib}^{\times \ell}$ is of the form Fib^λ for some set theoretic partition λ of $[\ell] = \{1 \dots \ell\}$.

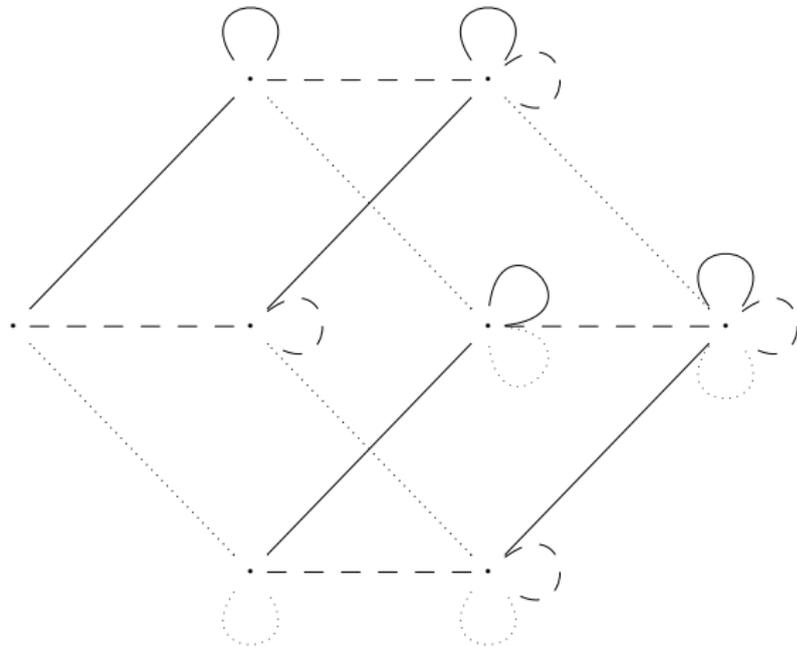
Take for example $\text{Fib}^{\times 2}$. There are two partitions of $[2]$ namely $\{1\} \cup \{2\}$ and $\{1, 2\}$ corresponding to the square



and the double interval



The partition $\{1\} \cup \{2\} \cup \{3\}$ of $[3]$ for $\text{Fib}^{\times 3}$ corresponds to the cube



Let A be an indecomposable algebra in $C = \mathcal{Fib}_u^{\boxtimes \ell}$.
 The category C_A of right A modules in C is then an
 indecomposable (left) C -module category. Recall that the forgetful
 functor $F : C_A \rightarrow C$ (forgetting the module structure) is a
 morphism of C -module categories and has a left adjoint
 $G : C \rightarrow C_A$ which is also a C -module category morphism. The left
 adjoint G sends the monoidal unit I to A as module over itself.

It then follows that for the NIM-representation M of C_A we have two maps of NIM-representations (a $\mathbb{Z}R$ -module homomorphism): $f : M \rightarrow \text{Fib}^{\times \ell}$ and $g : \text{Fib}^{\times \ell} \rightarrow M$ which are adjoint in the sense of the rigidity condition

$$(g(y), m)_M = (y, f(m))_{\text{Fib}^{\times \ell}}$$

Since C_A is indecomposable as a C -module category, so is its NIM-representation M .

The afore theorem says we should have $M \simeq \text{Fib}^\lambda$ for some set-theoretic partition λ of $[\ell]$.

In the first instance suppose λ has only one part $\lambda = (\ell)$.

In particular $M = \text{Fib}^{(\ell)}$ has just two simple objects: m and n .

Assume that $m = g(1)$.

Since g is a map of NIM-representations we must have $g(x_i) = n$ for all $i = 1, \dots, \ell$ such that

$$g(x_i * 1) = x_i g(1) = x_i * m = n$$

Hence for an arbitrary element $x_{i_1} \dots x_{i_s}$ of $\text{Fib}^{\times \ell}$

$$g(x_{i_1} \dots x_{i_s} * 1) = f_s * n + f_{s-1} * m$$

where f_s is the s -th Fibonacci number.

Thus the adjoint map f has the form

$$f(m) = 1 + \sum_{s=1}^{\ell} f_{s-1} * \sum_{i_1 < \dots < i_s} x_{i_1} \dots x_{i_s}$$

$$f(n) = \sum_{s=1}^{\ell} f_s * \sum_{i_1 < \dots < i_s} x_{i_1} \dots x_{i_s}$$

We identify $f(m)$ with the class of the algebra A in $K_0(\mathcal{Fib}_u^{\boxtimes \ell}) = \mathbb{Z}[\text{Fib}^{\times \ell}]$. It turns out that since the twist $\theta_{x_{i_1} \dots x_{i_s}} = \theta_{x_{i_1}} \dots \theta_{x_{i_s}}$ of each of the arbitrary elements depends on s in a non-trivial way, A cannot be ribbon.

Now suppose λ is a set-theoretic partition of $[\ell]$ into ordered parts

$$\lambda = [1 \dots \ell_1][\ell_1 \dots \ell_2] \dots [\ell_{n-1} \dots \ell_n]$$

Our theorem says that M , as a NIM-representation of $\mathcal{Fib}_u^{\boxtimes \ell} = \mathcal{Fib}_u^{\boxtimes \ell_1} \boxtimes \dots \boxtimes \mathcal{Fib}_u^{\boxtimes \ell_s}$, has the form

$$M = \text{Fib}^{(\ell_1)} \boxtimes \dots \boxtimes \text{Fib}^{(\ell_s)}$$

By the baby case above each $\mathbb{Z}[\text{Fib}^{(\ell_i)}] = K_0((\mathcal{F}ib_u^{\boxtimes \ell_i})_{A_i})$ for some non-ribbon algebra $A_i \in \mathcal{F}ib_u^{\boxtimes \ell_i}$.

Then $\mathbb{Z}M = K_0(\boxtimes_{j=1}^s (\mathcal{F}ib_u^{\boxtimes \ell_j})_{A_j}) = K_0((\mathcal{F}ib_u^{\boxtimes \ell})_A)$ where $A = \boxtimes_{j=1}^s A_j$.

Since A_j are non-ribbon then so is A .

The case for general λ can be reduced to the above by a permutation of $[n]$ which proves the following.

Theorem

There are no (non-trivial) ribbon commutative algebras in $\mathcal{Fib}_u^{\boxtimes \ell}$.

The argument of the proof of the above theorem works well for $\mathcal{Fib}_u^{\boxtimes \ell} \boxtimes \mathcal{Fib}_v^{\boxtimes m}$ as long as $uv \neq 1$.

Thank you!