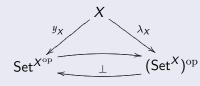
On the Isbell conjugation adjunction for monad-quantale-enriched categories

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 $\mathsf{hom}: X^\mathrm{op} \times X \to \mathsf{Set}$

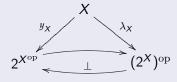


$$\leq: X^{\mathrm{op}} \times X \to 2$$

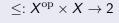
$$X$$

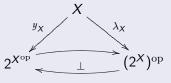
$$2^{X^{\mathrm{op}}} \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad$$

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- Both sides define lax idempotent monads on Ord.
- algebra=(co)complete ordered set.
- complete=cocomplete.
- There is a distributive law.

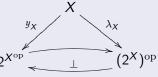




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How to do this on $\mathbb{T}V$?

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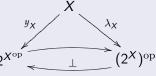


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How to do this on TV?

Now we consider: $a: TX \times X \to V$ with $\begin{cases} 1_X \leq a \cdot e_X, \\ a \cdot Ta \leq a \cdot m_X \end{cases}$

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How to do this on $\mathbb{T}V$?

Now we consider: $a: (TX)^{\mathrm{op}} \otimes X \to V$ with $\begin{cases} 1_X \leq a \cdot e_X, \\ a \cdot Ta \leq a \cdot m_X \end{cases}$

• \mathbb{T} extends to a monad on V-Cat; $(X, a_0 : X \longrightarrow X) \mapsto (TX, Ta_0 : TX \longrightarrow TX).$

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 $\mathsf{V}\text{-}\mathsf{Cat}^{\mathbb{T}} \xrightarrow{\mathsf{T}} (\mathbb{T},\mathsf{V})\text{-}\mathsf{Cat}, \begin{cases} \mathsf{K}(\mathsf{X},\mathsf{a}_0,\alpha) = (\mathsf{X},\mathsf{a}_0\cdot\alpha:\mathsf{TV}\longrightarrow\mathsf{V}), \\ \mathsf{M}(\mathsf{X},\mathsf{a}) = (\mathsf{TX},\mathsf{Ta}\cdot\mathsf{m}_\mathsf{X},\mathsf{m}_\mathsf{X}) \end{cases}$

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ullet induces lax idem. monad $\mathbb T$ on $(\mathbb T,\mathsf V) ext{-}\mathsf{Cat},\,\mathsf V ext{-}\mathsf{Cat}^\mathbb T\simeq (\mathbb T,\mathsf V) ext{-}\mathsf{Cat}^\mathbb T$.

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- (X, a) is called representable if $X \overset{e_X}{\frown} TX$. Then $a = a_0 \cdot \alpha$.
- For X = (X, a) representable: $X^{op} = (X, a_0^{\circ} \cdot \alpha)$.
- X = (X, a) representable $\Rightarrow a \cdot Ta = a \cdot m_X \Rightarrow X \otimes$ -exponentiable.

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• For $f: X \to Y$ and (X, a), (Y, b) representable:

$$f$$
 is (\mathbb{T},V) -functor $\iff egin{cases} f ext{ is V-functor,} \ f(lpha(\mathfrak{x})) \geq eta(\mathit{U}f(\mathfrak{x})). \end{cases}$

$$(TX)^{\operatorname{op}} \xrightarrow{\Lambda_X} V^X$$

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$$\downarrow^{\gamma} \qquad \qquad \downarrow^{\gamma}$$

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$$\downarrow^{\eta_X} \qquad \downarrow^{\gamma} \qquad e_X \qquad \downarrow^{\chi}$$

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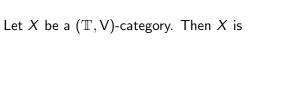
ullet $\omega = (\omega_X)$ is a nat. transform. when restricted to repres. cat's.

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- $oldsymbol{\omega} = (\omega_X)$ is a nat. transform. when restricted to repres. cat's.
- ullet $\mathbb{Q}=(Q,\omega,\lambda)$ is a lax idemp. monad on V-Cat $^{\mathbb{T}}\simeq (\mathbb{T},\mathsf{V})$ -Cat $^{\mathbb{T}}.$



• cocomplete whenever "all weighted colimits exist"

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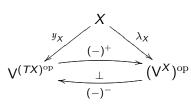
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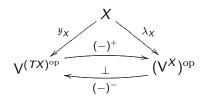
• totally complete if $X \xrightarrow{\top} (V^X)^{op}$ in V-Cat^T.

• We have



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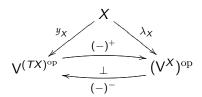
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 \bullet $(-)^-$ is even a (\mathbb{T},V) -functor, but $(-)^+$ in general not.

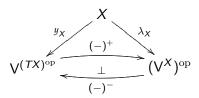
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- $(-)^-$ is even a (\mathbb{T}, V) -functor, but $(-)^+$ in general not.
- For X repres.: X cocomplete $\iff X$ complete.
- However: $[0,\infty]^{\mathrm{op}}$ (in $(\mathbb{U},[0,\infty])$ -Cat) is totally complete but not totally cocomplete.

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- **1** $Qf: QX \to QY$ has a left adjoint (in (\mathbb{T}, V) -Cat).
- $V^f: V^Y \to V^X$ is a homomorphism.
- f is weakly open, i.e. $f^{\circ} \cdot b \leq a \cdot Tf^{\circ} \cdot Tb_0$.

In Top:
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 $\mapsto \mathfrak{y}' \leq \mathfrak{y}$

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- f is weakly open, i.e. $f^{\circ} \cdot b \leq a \cdot Tf^{\circ} \cdot Tb_0$.

Hence, for $f: X \to Y$ (\mathbb{T}, V)-functor with X, Y totally complete:

f is right adjoint in (\mathbb{T}, V) -Cat $\iff \begin{cases} f \text{ preserves infima} \\ \text{and is "weakly open"}. \end{cases}$

$$\varphi: X \to QY \text{ in V-Cat}^{\mathbb{T}} = \left\{ \begin{array}{c} \text{V-module } \varphi: X \longrightarrow Y \text{ where} \\ TX \xrightarrow{T\varphi} TY \\ \alpha_* \phi & \phi \beta_* \\ X \xrightarrow{\varphi} Y \end{array} \right.$$

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 $\left(\qquad \dot{X} \xrightarrow{\circ} \dot{Y} \right)$

Hence $(V-Cat^{\mathbb{T}})_{\mathbb{Q}} \simeq \mathbb{Q}$ -Mod.

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For Top:
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$$(-)_* \qquad \mathsf{OrdCompHaus}$$

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- δ_X iso \iff X is Priestley. (Hence: X Priestley) \Rightarrow QX Priestley)

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- ullet Priest $_{\mathbb Q} \simeq \mathsf{DLat}^{\mathrm{op}}_{\vee,\perp}$. (Hence: $\mathsf{Stone}_{\mathbb V} \simeq \mathsf{Bool}^{\mathrm{op}}_{\vee,\perp}$)