Coherently closed tangent categories The link between SDG and the  $\partial$ - $\lambda$ -calculus

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# CCCs for differential geometry

There are two main approaches that have appeared for axiomatizing differential geometry that is meant to be performed in a CCC.

#### Synthetic differential geometry

- Led to drastic simplifications of many parts of differential geometry
- Also greatly simplified constructions by using internal logic
- Led to "previously undreamed of" opportunities

#### The differential $\lambda$ -calculus

- The simply typed version led to monoidal differential categories
- Then Cartesian differential categories
- And ultimately to tangent categories

# Synthetic differential geometry

Lawvere '67:

- Need a ring R and an ideal D ⊆ R such that d<sup>2</sup> = 0 for all d ∈ D;
- The map R × R → [D, R] given by (a, b) → λd.d ⋅ a + b must be an isomorphism;
- ▶ Need spaces *M* that believe that *D* is well behaved. For example, all spaces must believe that



where  $D(2) = \{(x, y) | x^2 = y^2 = xy = 0\}$  is a pushout.

▶ D is so small that [D, ] has a right adjoint.

# Differential semantics of $\lambda$ -calculi

Wanted: a model of linear logic where

- Proofs are interepreted as smooth (or at least continuous) functions
- Linear proofs are interpreted as linear proofs in a mathematical sense

So that linearity (proof theory) is captured mathematically (by being represented by a line).

One can then obtain a linear (proof theory) approximation to a proof...

In the models, this turned out to be a differential operator [3, 2]

## The differential $\lambda$ -calculus

Interestingly, the differential operator is a special cut or composition:

 $\frac{v : \text{Abstract vector quantity}}{\Gamma, x : A \vdash m : B \quad \Gamma \vdash a : A \quad \Gamma \vdash v : A}{\Gamma \vdash \frac{dm}{dx} (a) \cdot v : B}$ 

Determines how v can be used linearly

Importantly, this differential composition operator is just composition with a tangent vector v. This is reminiscent of SDG:

$$T(f)(v) := D \xrightarrow{v} M \xrightarrow{f} N$$

# The differential $\lambda$ calculus

Important theorems:

Theorem (Ehrhard-Regnier)

The equational theory of the  $\lambda_{\partial}$  may be oriented into a rewriting system that is confluent modulo.

Corollary (Ehrhard-Regnier)

 $\lambda_\partial$  is a conservative extension of  $\lambda$ .

#### Slogan: The $\lambda$ -calculus was smooth all along!

# Today's talk

- These two formal systems are more intimately related than the metaphor of linear composition
- A link will be established using abstract tangent structure studied by [4] and [1]

A tangent structure on  $\mathbb{X}$  is an endofunctor, that sends an object M, naturally, to a commutative monoid in  $\mathbb{X}/M$ , together with two additional natural transformations.

## Tangent structure

Axiom 1 [Additive bundle]: A category has an additive bundle structure, when there is an endofunctor  $\mathbb{X} \xrightarrow{T} \mathbb{X}$ .

Τ There must be a natural transformation 11

There must be natural +, 0 giving p the structure of a commutative monoid in  $\mathbb{X}/M$ .

## Tangent structure

A transverse system is a collection of pullback squares in a category  $\mathcal{Q}$ . An endofunctor  $\mathbb{X} \xrightarrow{T} \mathbb{X}$  is transverse when  $T(q) \in \mathcal{Q}$  for all  $q \in \mathcal{Q}$ .

Axiom 2 [Transerversality]: The tangent functor must be transverse for some transverse system containing the pullback powers of p and another pullback shown below.

The idea is that only pullbacks of transverse maps behave well.

As T preserves transverse pullbacks and equations,  $T^2M$  is a commutative monoid in  $\mathbb{X}/TM$  in two different ways:

$$\begin{array}{ccc}
T^2 M & T^2 M \\
P_{TM} & & \downarrow T(p_M) \\
TM & TM
\end{array}$$

#### Axiom 3 [Canonical flip]: There is an involution



that exchanges the addition.

## Tangent structure

Axoim 4 [Vertical lift]: There is a lift:

$$TM \xrightarrow{I} T^{2}M$$

$$\downarrow Tp$$

$$M \xrightarrow{I} TM$$

that lifts the additive structure of p to the additive structure of Tp.

#### Axiom 5 [Coherence]:

# l and c form an internally commutative cosemigroup in $[\mathbb{X},\mathbb{X}]$ Add string diagrams

Axiom 6 [Universality of lift]:



is a transverse pullback.

## Tangent structure

A tangent category is a category with an endofunctor  $\mathbb{X} \xrightarrow{T} \mathbb{X}$  that satisfies axioms 1-6.

## Examples of tangent categories

- Smooth manifolds. The usual tangent bundle of a manifold.
- ▶ Finitely presented affine schemes. Not generally cartesian closed; however, Z[x]/x<sup>2</sup> is exponentiable.
   TA := [Z[x]/x<sup>2</sup>, A].
- Any model of synthetic differential geometry.
- Any cartesian differential category.
- The classifying category of the differential  $\lambda$ -calculus.
- The idempotent splitting of any tangent category.
- The manifold construction applied to a tangent category.
- Presheaves
- Functor categories into a tangent category.

The differential  $\lambda\text{-calculus}$  provides the internal logic of Euclidean vector spaces in SDG.

# Differential objects

- SMan contains the category CartSp
- A category of microlinear spaces contains the category of euclidean vector spaces
- ▶ Both cases: determined as monoids V for which  $TV \xrightarrow{p} V$  has a left projection in a product diagram

$$V \stackrel{\hat{p}}{\longleftarrow} TV \stackrel{p}{\longrightarrow} V$$

 Objects that behave like vector spaces are called differential objects.

## **Differential Objects**

An object V is a **differential object** when there is a map  $TV \xrightarrow{\hat{p}} V$ , so product, and it is a commutative monoid  $(V, \sigma, \zeta)$ , such that



# The differential of differential objects

There is a 'combinator'

$$\frac{A \xrightarrow{f} B}{D[f] := A \times A \xrightarrow{\simeq} T(A) \xrightarrow{T(f)} T(B) \xrightarrow{\hat{p}} B}$$
linear

Theorem

The differential objects of a tangent category are always a Cartesian differential category.

A tangent category with finite products is **cartesian** when pullbacks along  $A \stackrel{!}{\longrightarrow} 1$  are transverse.

In particular, we have  $T(A) \times T(B) \simeq T(A \times B)$ .

Then tangent categories are strong by:

$$\theta := A \times TB \xrightarrow{0 \times 1} TA \times TB \xrightarrow{\simeq} T(A \times B)$$

When  $\ensuremath{\mathbb{X}}$  is a closed tangent category, strength gives:

$$\frac{A \times T[A, B] \xrightarrow{\theta} T(A \times [A, B]) \xrightarrow{Tev} TB}{T[A, B] \xrightarrow{\psi = \lambda(\theta Tev)} [A, TB]}$$

as an exponentiable strength.

Strength

Moreover, all the structural transformations are strong and exponentially strong.

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Coherently differential closed tangent categories

When V is a monoid, so is [M, V].

Differential objects are **additively coherent** when the canonical monoid structure, is the only one.

#### Proposition

In a cartesian closed category with additively coherent monoids, the differential objects are a cartesian closed left additive category.

### Corollary

In this case, the differential objects model the algebraic  $\lambda$ -calculus.

Coherently differential closed tangent categories

The differential objects are **differentially coherent** when the strength

 $T[A, B] \rightarrow [A, TB]$ 

is an isomorphism.

Differential coherence implies additive coherence

Proposition

The classifying category  $C[\lambda_{\partial}]$  may given the structure of a tangent category with coherent differential objects.

## Proposition

The differential objects of a differentially coherent tangent category are a model of the differential  $\lambda$ -calculus.

## Euclidean vector spaces

#### Corollary

The only models of the differential  $\lambda$ -calculus arise as the differential objects of a closed tangent category with coherent differential objects.

# Euclidean vector spaces

Representable tangent category: TM = [D, M]

$$T[M, V] = [D, [M, V]] \xrightarrow{\psi} [M, [D, V]] = [M, TV]$$

is the swap isomorphism

#### Corollary

The Euclidean vector spaces in SDG are a model of the differential  $\lambda$ -calculus.

## Corollary

Convenient vector spaces are a model of the differential  $\lambda$ -calculus.

#### Proposition

The strength  $A \times TB \xrightarrow{\theta} T(A \times B)$  is a comonad lifting law:

$$T; (A \times \_) \xrightarrow{\theta} (A \times \_); T$$

#### Corollary

The tangent functor lifts to each coKleisli category,  $\mathbb{X}[A]$ .

Proposition Each of the structural transformations lifts against  $\theta$ .

The strength of the structure transformations guarantees they lift. Corollary

All the structure transformations lift to each  $\mathbb{X}[A]$ 

#### Corollary

When X is a tangent category, each X[A] is a tangent category.

Explicitly,

$$\frac{B \xrightarrow{f} C \quad \mathbb{X}[A]}{A \times B \xrightarrow{f} C \quad \mathbb{X}}$$

$$\frac{A \times TB \xrightarrow{\theta} T(A \times B) \xrightarrow{Tf} TC \quad \mathbb{X}}{TB \xrightarrow{\theta T(f)} TC \quad \mathbb{X}[A]}$$

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When  $\mathbb X$  is closed, the simple slice may be presented as the Kleisli category for  $[A,\_].$ 

## Proposition

When  $\mathbb X$  is a closed tangent category, for which the exponential strength:

$$T[A,B] \xrightarrow{\psi} [A,TB]$$

is an isomorphism, then  $\mathbb{X}^A$  is an equivalent tangent category to  $\mathbb{X}[A].$ 

When  $\psi$  is an isomorphism, a tangent category is **coherently closed**.

The differential  $\lambda\text{-calculus}$  provides the internal logic of Euclidean vector bundles in SDG.

Euclidean vector spaces in SDG  $\rightsquigarrow$  differential  $\lambda$ -calculus.

1. Euclidean vector bundles are a central focus in differential geometry

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- 2. Locally, they are Euclidean vector spaces...
- 3. We would like to get at these!

Euclidean vector spaces in SDG  $\rightsquigarrow$  differential  $\lambda$ -calculus.

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Locally means in a slice category

# Locally Cartesian Closed Categories

A map is  $A \xrightarrow{f} X$  is exponentiable when for all  $B \xrightarrow{g} A$  there is a terminal pullback extension [5]:



# Locally Cartesian Closed Categories

A category is **locally cartesian closed** when it has finite limits, and every map is exponentiable.

This yields a Cartesian closed structure in each slice by taking the terminal diagram of



yields [f, g].

Need:

- ► The slice of a tangent category by M is a tangent category with T<sub>M</sub>
- ► The coherence isomorphism  $T_M[A, B] \simeq [A, T_M B]$  in every slice

## Theorem (CC'14)

When a tangent category has finite limits<sup>1</sup>, the slice is always a tangent category.

## Slice tangent structure

The slice tangent structure in  $\mathbb{X}/B$  is given by pullback along  $B \xrightarrow{0} TB$ . For  $T_B(h)$ :



## Slice tangent structure

The slice tangent structure in  $\mathbb{X}/B$  is given by pullback along  $B \xrightarrow{0} TB$ . For  $p_B : T_B(E) \to E$ :



# Differential bundles

## Theorem (CC'14)

Suppose X has finite limits. Differential objects in the tangent category X/B are precisely differential bundles over B in X.

This theorem says that in fact differential bundles really do provide local linear algebra.

It also formally reduces what is needed to:  $T_M[A, B] \simeq [A, T_M B]$ in the slice over M.

Proposition

When a tangent functor satisfies

$$T\left[\prod_{x:A}B\right]\simeq\left[\prod_{x:A}TB\right]$$

then in each slice it has

$$T_M[A,B] \simeq [A,T_MB]$$

Follows from Beck-Chevalley

A locally cartesian closed tangent category is **locally coherently** closed when  $T[\prod_{x:A} B] \simeq \prod_{x:A} TB$ 

## $T_M[q_E,q_F]$

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 $T_M[q_E, q_F] \\ \equiv 0^* (T(\prod q_E e_F))$ 

 $T_M[q_E, q_F] \\\equiv 0^* (T(\prod q_E e_F)) \\\simeq 0^* (\prod q_E(Tq_F))$ 

 $T_M[q_E, q_F] \\\equiv 0^* (T(\prod q_E e_F)) \\\simeq 0^* (\prod q_E(Tq_F)) \\\simeq \prod q_E(0^* T(q_F))$ 

 $T_{M}[q_{E}, q_{F}]$   $\equiv 0^{*}(T(\prod q_{E}e_{F}))$   $\simeq 0^{*}(\prod q_{E}(Tq_{F}))$   $\simeq \prod q_{E}(0^{*}T(q_{F}))$   $\equiv [q_{E}, T_{M}q_{F}]$ 

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#### Thus

### Corollary

Suppose X has finite limits and is locally coherently closed, then each X/B is a coherently closed tangent category. The differential objects of X/B is a model of the differential  $\lambda$ -calculus.

# Euclidean Vector Bundles

#### Proposition

In a locally cartesian closed tangent category, the category of differential bundles over B is a model of the differential  $\lambda$ -calculus.

Hence,

Corollary

In any model of Nishimura's axiomatic differntial geometry, the category of Euclidean vector bundles over B are a model of the differential  $\lambda$ -calculus.

# Conclusion

- In SDG and SMan, Euclidean vector spaces are modelled by the differential λ-calculus.
- In SDG and SMan, Euclidean vector bundles are modelled by the differential λ-calculus.
- The more general reasoning in microlinear spaces is still needed for global reasoning about vector bundles.

# Ongoing projects

- Formalizing more logic and differential geometry in bundle categories.
  - ► For example, local monoidal structure with duals...
- Formulating a direct syntax for all tangent categories
  - Including a syntax for the dependently typed differential λ-calculus.
  - For the latter, this is necessary adding equalizers or identity types.
- Physics in tangent categories with substructual aspects of the logic in this talk...

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