Algebraic Structure From Non-Algebraic Proofs

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Definition (Grandis, Tholen)

An algebraic weak factorisation system (awfs) (a.k.a natural weak factorisation system) on a category \mathbb{C} consists of a comonad L: $\mathbb{C}^2 \to \mathbb{C}^2$ and monad R: $\mathbb{C}^2 \to \mathbb{C}^2$ where the underlying copointed and pointed endofunctors arise from a functorial factorisation, satisfying a "distributive law."

We can used awfs's to get a structured notion of (trivial) cofibrations and fibrations. What about weak equivalences?

Definition

Suppose we are given a morphism of awfs's $\xi : (C^t, F) \rightarrow (C, F^t)$.

We say a *structured weak equivalence* is a map $f: X \to Y$ together with an F^t-algebra structure on Ff.

A *morphism of structured weak equivalences* is a commutative square



such that the induced map $Ff \to Ff'$ is a morphism of F^t -algebras. Write the resulting category as W-**Map**

Consider diagrams of the following form:



A *functorial* 3-*for*-2 *operator* takes such a diagram together with weak equivalence structures on two of the maps, and returns a weak equivalence structure on the third map, preserving morphisms of structured weak equivalences.

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Definition

An *algebraic model structure with structured weak equivalences* consists of

- 1. A morphism of algebraic weak factorisation systems $\xi : (C^t, F) \rightarrow (C, F^t).$
- 2. A functorial 3-for-2 operator.

Functors of the form below, that commute with forgetful functors.

- 1. C-coalg $\times_{\mathbb{C}^2} W$ -Map $\rightarrow C^t$ -coalg
- 2. C^t -coalg \rightarrow W-Map
- 3. $F-Alg \times_{\mathbb{C}^2} W-Map \rightarrow F^t-Alg$
- 4. F^{t} -Alg \rightarrow W-Map

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- The original motivation for this definition was the construction of identity types in cubical sets.

Theorem (S)

Suppose that \mathbb{C} is a category with an ams with structured weak equivalences and a "stable functorial choice of path objects." Then \mathbb{C} also has a "stable functorial choice of very good path objects" (which can be used for identity types, giving explicit definitions for *J*-terms).

- This is based on the notion of algebraic model structure due to Riehl.
- The original motivation for this definition was the construction of identity types in cubical sets.
- Unpublished results by Sattler suggest a wide range of categories (including CCHM cubical sets) can be made in algebraic model structures with structured weak equivalences.

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Theorem (Sattler)

Suppose (\mathbb{C}, \otimes) is a finitely complete and finitely cocomplete symmetric affine monoidal closed category and we are given

- 1. an interval object $\delta_0, \delta_1 \colon 1 \to \mathbb{I}$
- 2. wfs's $(\mathcal{C}, \mathcal{F}^t)$ and $(\mathcal{C}^t, \mathcal{F})$
- 3. $C^t \subseteq C$
- 4. C is closed under pullbacks
- 5. f is a fibration if and only if $\delta_0 \hat{\oslash} f$ and $\delta_1 \hat{\oslash} f$ are trivial fibrations.
- 6. $[\delta_0, \delta_1] \hat{\oslash} preserves trivial fibrations.$
- 7. Trivial cofibrations are stable under pullback along fibrations
- 8. (Trivial) Fibrations extend along trivial cofibrations

Define W to be maps of the form $f \circ m$ where $m \in C^t$ and $f \in \mathcal{F}^t$. Then (C, \mathcal{F}, W) is a model structure.

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Usually, the most natural way to prove 5 is to show C is generated by a set I and (C^t, \mathcal{F}) is cofibrantly generated by maps $\delta_i \hat{\otimes} m$ where $m \in I$ and i = 0, 1.

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Usually, the most natural way to prove 6 is show that $[\delta_0, \delta_1] \hat{\otimes} m \in \mathcal{C}$ for $m \in I$

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We will make this idea precise using Grothendieck fibrations.

- 1. Making it easy to formalise (without missing out too many details, or leaving them as exercises for the reader).
- 2. (Hopefully) The same ideas apply to variants based on realizability, allowing us to extract computational information telling us how to compute the operators.

Definition

Let \mathbb{C} be a category. We define the Grothendieck fibration of category indexed families $p: \operatorname{Fam}(\mathbb{C}) \to \operatorname{CAT}$ as follows. An object of $\operatorname{Fam}(\mathbb{C})$ consists of $\mathbb{A} \in \operatorname{CAT}$ together with a functor $\mathbb{A} \to \mathbb{C}$. *p* is defined to be the projection functor.

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For the proof to work, we need **CAT** to contain categories the same cardinality as \mathbb{C} (so in particular $\mathbb{C} \in CAT$).

By a well known result due to Freyd, \mathbb{C} cannot have colimits of shape \mathbb{A} for all $\mathbb{A} \in \mathbf{CAT}$ unless it is a poset. Hence in this case p is not a bifibration.

Theorem

Suppose (\mathbb{C}, \otimes) is a finitely complete and finitely cocomplete symmetric affine monoidal closed category and we are given

- 1. an interval object $\delta_0, \delta_1 \colon 1 \to \mathbb{I}$
- 2. awfs's (C^t, F) and (C, F^t)
- 3. (C, F^t) is stable under pullback
- 4. (C, F^t) is cofibrantly generated by a functor $M: J \to \mathbb{C}^2$
- 5. (C^t, F) is cofibrantly generated by $\delta_0 \hat{\otimes} M + \delta_1 \hat{\otimes} M$
- There is an endomorphism E : J → J such that M ∘ E ≃ [δ₀, δ₁]⊗M
- 7. Trivial cofibrations are functorially stable under pullback along fibrations

8. (Trivial) Fibrations functorially extend along trivial cofibrations Then we define an ams with structured weak equivalences on \mathbb{C} .

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- 7. Trivial cofibrations are *functorially* stable under pullback along *fibrations*

8. (Trivial) Fibrations functorially extend along trivial cofibrations Then we define an ams with structured weak equivalences on \mathbb{C} . (C^t, F) and (C, F^t) uniquely extend to fibred awfs's over Fam $(\mathbb{C}) \rightarrow CAT$. We just define them pointwise.

(Note that each fibre category $Fam(\mathbb{C})_{\mathbb{A}}$ is the functor category $[\mathbb{A}, \mathbb{C}]$ and the restrictions of (C, F^t) and (C^t, F) are the pointwise awfs's.)

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We then translate all of definitions to properties of the fibration.

Proposition

Suppose G is a vertical map over the category \mathbb{A} . Then

- 1. We can view G as a functor $\mathbb{A} \to \mathbb{C}^2$
- 2. F-algebra structures on G correspond to a choice of F-algebra structure on G(A) for each $A \in \mathbb{A}$ such that $G(\sigma)$ is a morphism of F-algebras for each morphism σ of \mathbb{A} .
- 3. Similarly for F^t-algebra, C-coalgebra, C^t-coalgebra, and weak equivalence structures.

For the cofibrantly generated parts we use the following definition.

Definition

Suppose $p \colon \mathbb{E} \to \mathbb{B}$ is a Grothendieck fibration, and (L, R) is a fibred awfs over p.

Suppose M is a vertical map in the fibre of J and G is a vertical map in the fibre of I.

A family of lifting problems from M to G is a map $\sigma \colon K \to J$ in \mathbb{B} together with a lifting problem from $\sigma^*(M)$ to G.

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A family of lifting problems from M to G is a map $\sigma \colon K \to J$ in \mathbb{B} together with a lifting problem from $\sigma^*(M)$ to G.

We say (L, R) is *cofibrantly generated* by a vertical map M if for all vertical maps G, R-algebra structures on G correspond naturally to coherent choices of fillers for all families of lifting problems from M to G.

We obtain the following structure over $p: \operatorname{Fam}(\mathbb{C}) \to \operatorname{CAT}$.

- 1. an interval object $\delta_0, \delta_1 \colon 1 \to \mathbb{I}$ over 1
- 2. Fibred awfs's (C^t, F) and (C, F^t)
- 3. (C, F^t) is stable under pullback
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- 5. (C^t , F) is cofibrantly generated by $\delta_0 \hat{\otimes} M + \delta_1 \hat{\otimes} M$
- 6. There is a levelwise cartesian square $[\delta_0, \delta_1] \hat{\otimes} M \to M$
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For every A, the fibre category $\mathsf{Fam}(\mathbb{C})_{\mathbb{A}}$ satisfies the conditions to apply Sattler's non-algebraic result. We deduce.

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Lemma

For every category $\mathbb{A} \in \mathbf{CAT}$, $\mathsf{Fam}(\mathbb{C})_{\mathbb{A}}$ is a model structure.

Suppose we want an operator composing weak equivalences.

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We define \mathbb{A} to be the category with objects pairs of structured weak equivalences $X \xrightarrow{g} Y \xrightarrow{h} Z$ in \mathbb{C} .

A morphism of \mathbb{A} is a diagram of the form below, where both squares preserve the weak equivalence structures.



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This defines two vertical maps G and H in $\operatorname{Fam}(\mathbb{C})_{\mathbb{A}}$, and moreover they are weak equivalences in $\operatorname{Fam}(\mathbb{C})_{\mathbb{A}}$. Since $\operatorname{Fam}(\mathbb{C})_{\mathbb{A}}$ is a model structure $H \circ G$ must also be a weak equivalence in $\operatorname{Fam}(\mathbb{C})_{\mathbb{A}}$. We deduce that we can assign $h \circ g$ the structure of a weak equivalence for each object (g, h) in \mathbb{A} , and moreover this assignment is functorial.

- This result can be used to show BCH cubical sets form an ams with structured weak equivalences.
- It may lead to a more efficient proof of Sattler's result that CCHM cubical sets (and many other categories) form an ams with structured weak equivalences.
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Thank you for your attention!