

# Limits and pointwise colimits in accessible categories

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First consider the following problem: given a class  $\mathcal{I}$  of small categories, describe a class  $\mathcal{S}$  of sketches such that a category  $\mathbf{A}$  is accessible and has  $\mathcal{I}$ -limits if and only if it is equivalent to  $\text{Mod}(\mathbf{S})$  for some sketch  $\mathbf{S}$  in  $\mathcal{S}$ . To quote only but the two easiest examples:

**Theorem 1.** [Ageron95] *A category  $\mathbf{A}$  is  $\alpha$ -accessible and has nonempty limits if and only if it is equivalent to  $\text{Mod}(\mathbf{S})$  for some sketch  $\mathbf{S}$  all of whose cones are of size  $< \alpha$  and all of whose cocones are based on the empty category.*

**Theorem 2.** [Lair96] *A category  $\mathbf{A}$  is accessible and has a terminal object if and only if it is equivalent to  $\text{Mod}(\mathbf{S})$  for some sketch  $\mathbf{S}$  all of whose cocones are based on connected categories. (Note: no control of the rank is possible here.)*

Other cases have been treated, either by Lair or by myself: accessible categories with pullbacks and equalizers [Ageron95], with pullbacks [Ageron96], with equalizers [Lair97], etc. In each case the problem amounts to determine the class of colimits (*not* meaning the class of indexations) which commute to the considered type of limits in **Set**.

In general, there is no hope to solve the similar problem for colimits: unlike limits, they cannot be *forced* to be pointwise. Accidentally however, there is an analogue to Theorem 1:

**Theorem 3.** [Ageron97] *A category  $\mathbf{A}$  is  $\alpha$ -accessible and has nonempty colimits if and only if it is equivalent to  $\text{Mod}(\mathbf{S})$  for some sketch  $\mathbf{S}$  all of whose cones are of size  $< \alpha$  and all of whose cocones: (i) share their vertex with a cone based on the empty category; and (ii) are based on categories with non empty limits of size  $< \alpha$  with the corresponding limit cones being distinguished in  $\mathbf{S}$ .*

**Corollary** (of Theorems 1 and 3). *The full subcategories of  $\alpha$ -ACC consisting respectively of  $\alpha$ -accessible categories with nonempty limits and of  $\alpha$ -accessible categories with nonempty colimits are Cartesian closed. (Note:  $\alpha$ -ACC itself is not Cartesian closed.)*

Any analogue to Theorem 2 has to be of a different form, involving a generalised notion of accessibility. First, say that an object  $A$  in a category with an initial object  $0$  is **nonempty** if  $\text{Hom}(A, 0) = \emptyset$ . Say that a category  $\mathbf{A}$  is  **$\alpha$ -positively accessible** if  $\mathbf{A}$  has  $\alpha$ -filtered colimits, an initial object, and a small full subcategory  $\mathbf{B}$  consisting of nonempty  $\alpha$ -presentable objects so that every object of  $\mathbf{A}$  is an empty or  $\alpha$ -filtered colimit of objects in  $\mathbf{B}$ . Then:

**Theorem 4.** [Ageron99] *A category  $\mathbf{A}$  is positively accessible if and only if it is*

equivalent to  $\text{Mod}(\mathbf{S})$  for some sketch  $\mathbf{S}$  all of whose cones are based on nonempty categories.

Now let  $\mathbf{S}$  be a sketch. Define  $\mathbf{S}^+$  as  $\mathbf{S}$  with a terminal object  $S$  added and the following items distinguished: the cone  $(id_S, id_S)$  based on the discrete category with two objects, the cones of  $\mathbf{S}$  completed with the arrows from them to  $S$ , the same cocones as those of  $\mathbf{S}$ .

**Theorem 5.** [Ageron99] (i) *Let  $\mathbf{S}$  be a sketch all of whose cones are based on nonempty  $\alpha$ -small categories. If  $\text{Mod}(\mathbf{S})$  is  $\alpha$ -positively accessible, then  $\text{Mod}(\mathbf{S})$  is  $\alpha$ -accessible.*

(ii) *Let  $\mathbf{S}$  be a sketch all of whose cones are of size  $< \alpha$ . If  $\text{Mod}(\mathbf{S})$  is  $\alpha$ -accessible, then  $\text{Mod}(\mathbf{S}^+)$  is  $\alpha$ -positively accessible.*

Results similar to Theorems 4 and 5 apply for example to sketches all of whose cones are based on *connected* categories: in this case Theorem 5 gives a result announced by S. Lack.