

A new SDP approach to the Max-Cut problem

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The Stable Set Problem

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Stable Set Problem - LP Formulation

Given a graph $G = (\{1, \dots, n\}, E)$ and a weight vector $\omega \in \mathbb{R}^n$, solve the linear program

$$\alpha(G, \omega) := \max_{x \in \text{STAB}(G)} \langle \omega, x \rangle.$$

Definition of Theta Body

Definition (Lovász ~ 1980)

Given a graph $G = (\{1, \dots, n\}, E)$ we define its theta body, $\text{TH}(G)$, as the set of all vectors $x \in \mathbb{R}^n$ such that

$$\begin{bmatrix} 1 & x^t \\ x & U \end{bmatrix} \succeq 0$$

for some symmetric $U \in \mathbb{R}^{n \times n}$ with $\text{diag}(U) = x$ and $U_{ij} = 0$ for all $(i, j) \in E$.

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Theorem (Lovász ~ 1980)

The relaxation is tight, i.e. $\text{TH}(G) = \text{STAB}(G)$, if and only if the graph G is perfect.

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We say a polynomial f is ***k*-sos modulo the ideal I** if and only if

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In particular, for any \mathbf{p} in the zero set $\mathcal{Z}(I)$ we have

$$f(\mathbf{p}) = h_1^2(\mathbf{p}) + \dots + h_m^2(\mathbf{p}) \geq 0,$$

so any *k*-sos polynomial is a nonnegative on the zero-set of the ideal.

Connection to Algebra

Theorem (Lovász ~ 1993)

TH(G) equals the intersection of all half-spaces

$$H_f = \{x \in \mathbb{R}^n : f(x) \geq 0\}$$

where f ranges over all affine polynomials that are 1-sos modulo $\mathcal{I}(S_G)$.

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This definition does not depend directly on the combinatorics of the graph, but only on the ideal $\mathcal{I}(S_G)$.

Theta Bodies of Ideals

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Given an ideal $I \subset \mathbb{R}[x_1, \dots, x_n]$ we define its k -th theta body, $\text{TH}_k(I)$, as the intersection of all half-spaces

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Remarks:

- $\overline{\text{conv}(\mathcal{Z}(I))} \subseteq \dots \subseteq \text{TH}_k(I) \subseteq \text{TH}_{k-1}(I) \subseteq \dots \subseteq \text{TH}_1(I)$.

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Remarks:

- $\overline{\text{conv}(\mathcal{Z}(I))} \subseteq \dots \subseteq \text{TH}_k(I) \subseteq \text{TH}_{k-1}(I) \subseteq \dots \subseteq \text{TH}_1(I)$.
- If $S \subset \mathbb{R}^n$ is a finite set and $I = \mathcal{I}(S)$ then for some k , we have $\text{TH}_k(I) = \text{conv}(S)$.

Combinatorial Moment Matrices

Let I be a polynomial ideal and

$$\mathcal{B} = \{1 = f_0, x_1 = f_1, \dots, x_n = f_n, f_{n+1}, \dots\}$$

be a basis of $\mathbb{R}[\mathbf{x}]/I$ and $\mathcal{B}_k = \{f_i : \deg(f_i) \leq k\}$ for all k .

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$$(f^k(\mathbf{x}))(f^k(\mathbf{x}))^t = \sum_{f_i \in \mathcal{B}} A_i f_i(\mathbf{x})$$

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for some symmetric matrices A_i . Given a vector y indexed by the elements in \mathcal{B} we define the **combinatorial moment matrix** of y as

$$M_{\mathcal{B},k}(y) = \sum_{f_i \in \mathcal{B}} A_i y_{f_i}.$$

Theta Bodies and Moment Matrices

Theorem (GPT)

Let I be a polynomial ideal and $\mathcal{B} = \{1, x_1, \dots, x_n, \dots\}$ a basis for $\mathbb{R}[\mathbf{x}]/I$. Let

$$\mathcal{M}_{\mathcal{B},k}(I) = \{y \in \mathbb{R}^{\mathcal{B}} : y_0 = 1; M_{\mathcal{B},k}(y) \succeq 0\}$$

then

$$TH_k(I) = \overline{\pi_{\mathbb{R}^n}(\mathcal{M}_{\mathcal{B},k}(I))}$$

where $\pi_{\mathbb{R}^n} : \mathbb{R}^{\mathcal{B}} \rightarrow \mathbb{R}^n$ is just the projection over the coordinates indexed by the degree one monomials.

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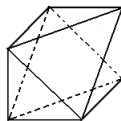
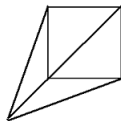
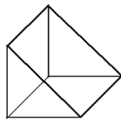
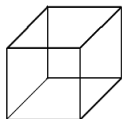
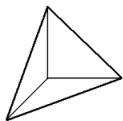
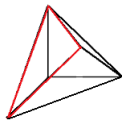
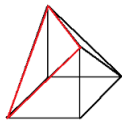
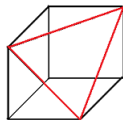
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A full characterization is possible for $k = 1$ in the case of vanishing ideals of finite sets in \mathbb{R}^n .

Theorem (GPT)

Let $S \subset \mathbb{R}^n$ be finite then $\mathcal{I}(S)$ is TH₁-exact if and only if for every facet defining hyperplane H of the polytope $\text{conv}(S)$ we have a parallel translate H' of H such that $S \subseteq H' \cup H$.

Examples in \mathbb{R}^3 TH₁-exactNot TH₁-exact

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The Problem

Given edge weights α we want to find which cut C maximizes

$$\alpha(C) := \sum_{(i,j) \in C} \alpha_{i,j}.$$

The Cut Polytope

Definition

The cut polytope of G , $\text{CUT}(G)$, is the convex hull of the characteristic vectors $\chi_C \subseteq \mathbb{R}^E$ of the cuts of G , where $(\chi_C)_{ij} = -1$ if $(i, j) \in C$ and 1 otherwise.

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LP formulation

Given a vector $\alpha \in \mathbb{R}^E$ solve the optimization problem

$$\text{mcut}(G, \alpha) = \max_{x \in \text{CUT}(G)} \frac{1}{2} \langle \alpha, \mathbf{1} - x \rangle.$$

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Theorem

If G is connected then the set

$$\{x_e^2 - 1 : e \in E\} \cup \{1 - \mathbf{x}^A : A \subseteq E, A \text{ circuit in } G\}$$

generates I_G , and

$$\mathcal{B} := \{\mathbf{x}^{F_T} : T \subseteq [n], |T| \text{ even}\}$$

is a basis for $\mathbb{R}[\mathbf{x}]/I_G$.

General Cut Theta body

Let \mathcal{B}_k be the set of all even $T \subseteq V$ such that $d_T \leq k$.

Theorem

The set $TH_k(I_G)$ is given by

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Cut Theta Body

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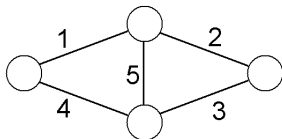
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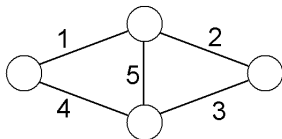
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for some a symmetric $U \in \mathbb{R}^{E \times E}$ with $\text{diag}(U) = \mathbf{1}$, if (e, f, g) is a triangle in G , $U_{e,f} = x_g$, and if $\{e, f, g, h\}$ forms a 4-cycle $U_{e,f} = U_{g,h}$.

Example

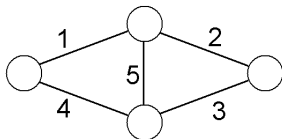


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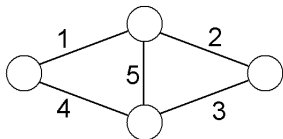
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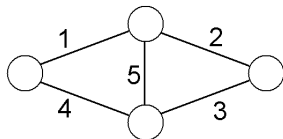
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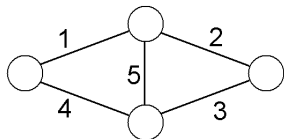
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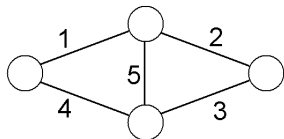
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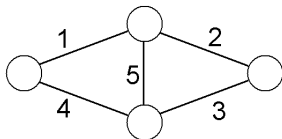
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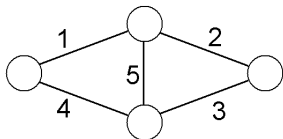
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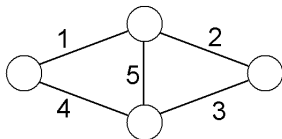
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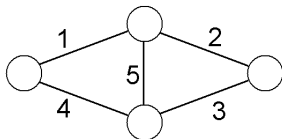
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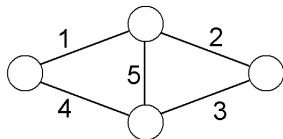
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 x_2 & & 1 & x_5 & & x_3 \\
 x_3 & & & 1 & ? & x_2 \\
 x_4 & & & & 1 & x_1 \\
 x_5 & & & & & 1
 \end{bmatrix}
 \stackrel{!}{=} 0$$

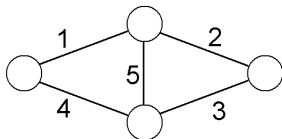
Example



$\text{TH}_1(I_G)$ is the set of $x \in \mathbb{R}^5$ such that

$$\begin{array}{c}
 0 \\
 1 \\
 2 \\
 3 \\
 4 \\
 5
 \end{array}
 \begin{bmatrix}
 1 & x_1 & x_2 & x_3 & x_4 & x_5 \\
 x_1 & 1 & y_1 & & x_5 & x_4 \\
 x_2 & & 1 & x_5 & & x_3 \\
 x_3 & & & 1 & y_1 & x_2 \\
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 x_5 & & & & & 1
 \end{bmatrix}
 \succeq 0$$

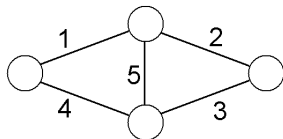
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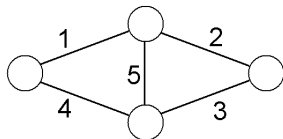
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Cut-Perfect Graphs

In analogy with the stable set results, it makes sense to have the following definition:

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Theorem (GLPT)

A graph is cut-perfect if and only if it has no K_5 minor and no chordless cycle of size larger than 4.

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- This technique can in theory be applied to any combinatorial problem to derive hierarchies. Results may vary.

The End

Thank You