

# Skew RSK and coincidence of Littlewood-Richardson commutators

Olga Azenhas, CMUC, University of Coimbra

September 10-13, 2017

79th Séminaire Lotharingien de Combinatoire

Joint session with

XXI Incontro Italiano di Combinatoria Algebrica

Bertinoro

# Littlewood-Richardson coefficients as structure constants

- *The ring of symmetric polynomials: the product of Schur polynomials.* Let  $x = (x_1, \dots, x_d)$ . Schur polynomials  $s_\lambda(x)$  for all partitions with  $\ell(\lambda) \leq d$  form a  $\mathbb{Z}$ -linear basis for the ring  $\Lambda_d := \mathbb{Z}[x]^{\mathfrak{S}_d}$  of symmetric polynomials in  $x$ ,

$$s_\mu s_\nu = \sum_{\lambda} c_{\mu\nu}^\lambda s_\lambda, \quad c_{\mu\nu}^\lambda \in \mathbb{Z}_{\geq 0}^+.$$

- *Schubert calculus of Grassmannians: the product in the cohomology of Grassmannians.* Schur polynomials  $s_\lambda(x)$  with  $\lambda$  inside a rectangle  $d \times (n-d)$  ( $0 < d < n$ ) may be interpreted as representatives of Schubert classes  $\sigma_\lambda$



Schubert classes  $\{\sigma_\lambda\}_{\lambda \subseteq n \times (n-d)}$  form a  $\mathbb{Z}$ -linear basis for the cohomology ring  $H^*(G(d, n))$  of the Grassmannian  $G(d, n)$  (the set of all complex  $d$ -dimensional linear subspaces of  $\mathbb{C}^n$ ), and

$$\sigma_\mu \sigma_\nu = \sum_{\lambda \subseteq d \times (n-d)} c_{\mu\nu}^\lambda \sigma_\lambda.$$

# Littlewood-Richardson coefficients

- These numbers  $c_{\mu, \nu}^{\lambda}$  also arise as
  - ▶ *tensor product multiplicities*. Schur polynomials  $s_{\lambda}(x)$  may be interpreted as irreducible characters of the general linear group  $GL_d(\mathbb{C})$ . The decomposition of the tensor product of two irreducible polynomial representations  $V^{\mu}$  and  $V^{\nu}$  of the general linear group  $GL_d(\mathbb{C})$  into irreducible representations of  $GL_d(\mathbb{C})$ , is given by

$$V^{\mu} \otimes V^{\nu} = \bigoplus_{\ell(\lambda) \leq d} V^{\lambda} \oplus c_{\mu \nu}^{\lambda}.$$

- *Positivity of Littlewood-Richardson coefficients in existence problems*. There exist  $d \times d$  non singular matrices  $A$ ,  $B$  and  $C$ , over a *local principal ideal domain*, with Smith invariants  $\mu$ ,  $\nu$  and  $\lambda$  respectively, such that  $AB = C$  iff  $c_{\mu \nu}^{\lambda} > 0$ .

# Littlewood-Richardson coefficients

- These numbers  $c_{\mu,\nu}^{\lambda}$  also arise as
  - ▶ *tensor product multiplicities*. Schur polynomials  $s_{\lambda}(x)$  may be interpreted as irreducible characters of the general linear group  $GL_d(\mathbb{C})$ . The decomposition of the tensor product of two irreducible polynomial representations  $V^{\mu}$  and  $V^{\nu}$  of the general linear group  $GL_d(\mathbb{C})$  into irreducible representations of  $GL_d(\mathbb{C})$ , is given by

$$V^{\mu} \otimes V^{\nu} = \bigoplus_{\ell(\lambda) \leq d} V^{\lambda} \oplus c_{\mu\nu}^{\lambda}.$$

- *Positivity of Littlewood-Richardson coefficients in existence problems*. There exist  $d \times d$  non singular matrices  $A$ ,  $B$  and  $C$ , over a *local principal ideal domain*, with Smith invariants  $\mu$ ,  $\nu$  and  $\lambda$  respectively, such that  $AB = C$  iff  $c_{\mu\nu}^{\lambda} > 0$ .
- *The commutativity of Littlewood-Richardson coefficients*

$$c_{\mu\nu}^{\lambda} = c_{\nu\mu}^{\lambda}$$

# Our structure coefficients are combinatorial numbers

- The structure coefficient  $c_{\mu,\nu}^\lambda$  is  
the cardinality of an explicit set of combinatorial objects.
- It is possible to determine  $c_{\mu,\nu}^\lambda > 0$  without determining its exact value.
- Exhibiting the commutation symmetry of these combinatorial objects means:  
Fix  $\mu, \nu \subset \lambda$  with  $|\mu| + |\nu| = |\lambda|$ , and your set  $\mathbb{LR}_{\mu,\nu,\lambda}$  of combinatorial objects counted by  $c_{\mu,\nu}^\lambda$ . An LR commutor is any bijection

$$\mathbb{LR}_{\mu,\nu,\lambda} \rightarrow \mathbb{LR}_{\nu,\mu,\lambda}$$

$$c_{\mu\nu}^\lambda = c_{\nu\mu}^\lambda.$$

# Coincidence of LR commutators

- **Problem** (Pak, Vallejo'04). Show that all LR commutators are involutions and **coincide**.

One has several LR commutators and all are involutions.

One involution needs to be included in the coincidence list :

Goal: Internal insertion LR commutator coincides with switching LR commutator.

- **The Littlewood-Richardson (LR) rule** (*D.E. Littlewood and A. Richardson 34; M.-P. Schützenberger 77; G.P. Thomas 74*). In the linear expansion of a product of two Schur polynomials

$$s_\mu(x) s_\nu(x) = \sum_{\lambda} c_{\mu\nu}^{\lambda} s_{\lambda}(x)$$

the coefficients are given by

$$c_{\mu\nu}^{\lambda} = \#\{\text{ballot SSYT of shape } \lambda/\mu \text{ and content } \nu\}.$$

The coefficients  $c_{\mu\nu}^{\lambda}$  are known as Littlewood–Richardson (LR) coefficients, and the ballot SSYT's are also known as Littlewood-Richardson tableaux.

- **LR tableau, ballot tableau** The content of each initial segment of the reading word, right to left across rows and top to bottom, is a partition.

<table style="border-collapse: collapse; text-align: left;"> <tr><td style="border: 1px solid black; width: 20px; height: 20px;"></td><td style="border: 1px solid black; width: 20px; height: 20px;"></td><td style="border: 1px solid black; width: 20px; height: 20px;">1</td><td style="border: 1px solid black; width: 20px; height: 20px;">1</td><td style="border: 1px solid black; width: 20px; height: 20px;">2</td></tr> <tr><td style="border: 1px solid black;"></td><td style="border: 1px solid black;">1</td><td style="border: 1px solid black;">2</td><td colspan="2"></td></tr> <tr><td style="border: 1px solid black;">1</td><td style="border: 1px solid black;">3</td><td colspan="3"></td></tr> </table>			1	1	2		1	2			1	3				2112131	<table style="border-collapse: collapse; text-align: left;"> <tr><td style="border: 1px solid black; width: 20px; height: 20px;"></td><td style="border: 1px solid black; width: 20px; height: 20px;"></td><td style="border: 1px solid black; width: 20px; height: 20px;">1</td><td style="border: 1px solid black; width: 20px; height: 20px;">1</td><td style="border: 1px solid black; width: 20px; height: 20px;">1</td></tr> <tr><td style="border: 1px solid black;"></td><td style="border: 1px solid black;">2</td><td style="border: 1px solid black;">2</td><td colspan="2"></td></tr> <tr><td style="border: 1px solid black;">1</td><td style="border: 1px solid black;">3</td><td colspan="3"></td></tr> </table>			1	1	1		2	2			1	3				1112231	<table style="border-collapse: collapse; text-align: left;"> <tr><td style="border: 1px solid black; width: 20px; height: 20px;"></td><td style="border: 1px solid black; width: 20px; height: 20px;"></td><td style="border: 1px solid black; width: 20px; height: 20px;">1</td><td style="border: 1px solid black; width: 20px; height: 20px;">1</td><td style="border: 1px solid black; width: 20px; height: 20px;">1</td></tr> <tr><td style="border: 1px solid black;"></td><td style="border: 1px solid black;">1</td><td style="border: 1px solid black;">2</td><td colspan="2"></td></tr> <tr><td style="border: 1px solid black;">2</td><td style="border: 1px solid black;">3</td><td colspan="3"></td></tr> </table>			1	1	1		1	2			2	3				1112132
		1	1	2																																														
	1	2																																																
1	3																																																	
		1	1	1																																														
	2	2																																																
1	3																																																	
		1	1	1																																														
	1	2																																																
2	3																																																	

one has 3 candidates 1,2,3 each receiving 4,2,1 votes respectively. A particular ordering of the votes is then a **ballot sequence** of length 7 where at any stage candidate 1 has at least many votes as candidate 2, and candidate 2 has at least many votes as candidate 3.

# Tableau Switching (Benkart-Sotille-Stroomer)

- A *perforated tableau*  $U$  of shape  $\lambda$  is a filling of some of the boxes in  $\lambda$  with integers satisfying some restrictions:  $x' \underset{x}{\quad}$ ,  $x' \ x \ x \geq x'$ ; and  $x' \underset{x}{\quad} x > x'$ .

Switch an integer with the neighbour empty box  $\blacksquare$  to the south, east, north or west, in a perforated tableau, so that the result is still a perforated tableau:

$$\begin{array}{l} \text{contracts } U \quad \begin{array}{c} \blacksquare \\ x \end{array} \xrightarrow{s} \begin{array}{c} x \\ \blacksquare \end{array} \qquad \blacksquare \ x \xrightarrow{s} x \ \blacksquare \\ \text{expands } U \quad \begin{array}{c} x \\ \blacksquare \end{array} \xrightarrow{s} \begin{array}{c} \blacksquare \\ x \end{array} \qquad x \ \blacksquare \xrightarrow{s} \blacksquare \ x. \end{array}$$

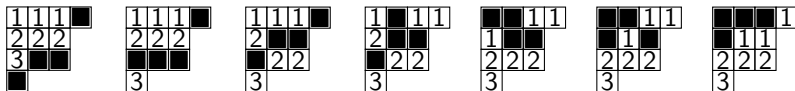
A *perforated tableau pair*  $U \cup V$  of shape  $\lambda$  is the superimposing of two perforated tableaux  $U$  and  $V$  of shape  $\lambda$ , so that together they completely fill  $\lambda$  and no two letters are in a same box.

- If  $\mathbf{u}$  and  $\mathbf{v}$  are vertically or horizontally adjacent letters from  $U$  and  $V$  respectively, then an interchanging of  $\mathbf{u}$  with  $\mathbf{v}$  is a *switch*, written  $\mathbf{u} \underset{s}{\leftrightarrow} \mathbf{v}$ , provided it produces a new perforated tableau pair,

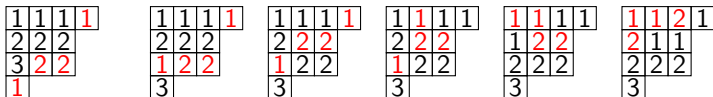
$$\begin{array}{c} \mathbf{u} \\ \mathbf{v} \end{array} \underset{s}{\leftrightarrow} \begin{array}{c} \mathbf{v} \\ \mathbf{u} \end{array} \qquad \mathbf{u} \ \mathbf{v} \underset{s}{\leftrightarrow} \mathbf{v} \ \mathbf{u}.$$



- *switching* in a perforated tableau: local moves preserve the perforated condition.

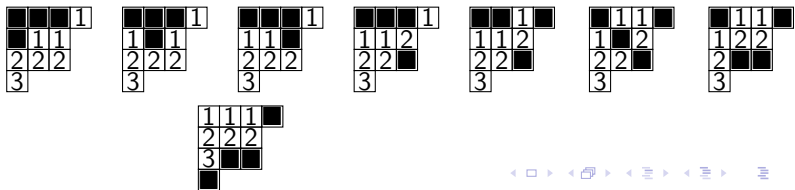


- *switching* in a perforated pair: local moves preserve the perforated condition.



- *Jeu de taquin* in a skew tableau is a particular case of *switching* where an order of the switches has been imposed. The game can be reduced to the local moves:

$$\begin{array}{c} \blacksquare \\ a \end{array} b \rightarrow \begin{array}{c} a \\ \blacksquare \end{array} b, \quad a \leq b \qquad \begin{array}{c} \blacksquare \\ a \end{array} b \rightarrow \begin{array}{c} b \\ \blacksquare \end{array}, \quad a > b$$



## Switching LR commutor $\rho_1$

- $\rho_1^{(n)} : \mathbb{LR}_{\mu,\nu,\lambda} \rightarrow \mathbb{LR}_{\nu,\mu,\lambda}, \ell(\lambda) \leq n.$

$$Y_{32} \cup T = \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 1 & 2 & 3 & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 1 & 2 & 1 & 2 \\ \hline 2 & 2 & 3 & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 1 & 1 & 2 & 2 \\ \hline 2 & 2 & 3 & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 1 & 1 & 2 & 2 \\ \hline 2 & 3 & 2 & \\ \hline \end{array} \xrightarrow{s}$$

$$\begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 1 & 1 & 2 & 2 \\ \hline 2 & 3 & 2 & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 1 & 2 & 2 \\ \hline 1 & 3 & 2 & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 1 & 3 & 2 & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 3 & 1 & 2 & \\ \hline \end{array} = Y_{321} \cup H, H \equiv T.$$

# Switching LR commutor $\rho_1$

- $\rho_1^{(n)} : \mathbb{LR}_{\mu,\nu,\lambda} \rightarrow \mathbb{LR}_{\nu,\mu,\lambda}, \ell(\lambda) \leq n.$

$$Y_{32} \cup T = \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 1 & 2 & 3 & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 1 & 2 & 1 & 2 \\ \hline 2 & 2 & 3 & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 1 & 1 & 2 & 2 \\ \hline 2 & 2 & 3 & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 1 & 1 & 2 & 2 \\ \hline 2 & 3 & 2 & \\ \hline \end{array} \xrightarrow{s}$$

$$\begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 1 & 1 & 2 & 2 \\ \hline 2 & 3 & 2 & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 1 & 2 & 2 \\ \hline 1 & 3 & 2 & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 1 & 3 & 2 & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 3 & 1 & 2 & \\ \hline \end{array} = Y_{321} \cup H, H \equiv T.$$

- Tableau sliding LR commutor (Thomas-Yong infusion). *To each entry of the inner shape (in red) associate a jeu de taquin slide.*

$$Y_{32} \cup T = \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 1 & 2 & 3 & \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 1 & 2 & 2 \\ \hline 1 & 2 & 3 & \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 1 & 2 & 2 \\ \hline 1 & 2 & 3 & \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 1 & 1 & 2 & 2 \\ \hline 2 & 2 & 3 & \\ \hline \end{array} \rightarrow$$

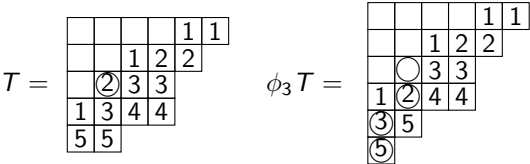
$$\begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 1 & 1 & 2 & 2 \\ \hline 2 & 3 & 2 & \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 1 & 1 & 2 & 2 \\ \hline 2 & 3 & 2 & \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 1 & 2 & 1 & 2 \\ \hline 2 & 3 & 2 & \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 1 & 2 & 1 & 2 \\ \hline 2 & 3 & 2 & \\ \hline \end{array} \rightarrow \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 1 & 3 & 2 & \\ \hline \end{array}$$

$$\rightarrow \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 3 & 1 & 2 & \\ \hline \end{array} = Y_{321} \cup H, H \equiv T.$$

# Internal and External insertion on skew-tableaux

For skew tableaux there are two types of row insertion: **external** and **internal**. External insertion is similar to Schensted's original procedure.

Sagan-Stanley internal insertion operator  $\phi_i$  on  $T$  a skew-tableau:



- Definition of internal insertion operator  $\phi_i$  on  $T$ . Need  $(i, \mu_i + 1)$  to be an inner corner of  $T$ :
  - ▶  $\phi_i$  bumps the entry, say  $x$ , in the inner corner cell  $(i, \mu_i + 1)$  of  $T$ , and then inserts (externally as in the Schensted procedure) the bumped element  $x$  in the subtableau consisting of the last  $n - i$  rows of  $T$ .

## Special case of Sagan-Stanley skew RSK correspondence

- *Internal row insertion correspondence*. Fix partitions  $\mu \subseteq \alpha, \beta$ . There is a bijection, defined below,

$$\begin{aligned}
 YT(\alpha/\mu) \times YT(\beta/\mu) &\longrightarrow \bigcup_{\substack{\lambda \\ |\lambda|=|\alpha|+|\beta|-|\mu|}} YT(\lambda/\beta) \times YT(\lambda/\alpha) \\
 (T, U) &\longrightarrow (P, Q),
 \end{aligned}$$

where  $P \equiv T$  and  $Q \equiv U$ . The  $P$ -tableau  $P$  is the internal row insertion of  $T$  whose sequence of inner corners containing the entries of  $T$  to be internally inserted is dictated by the entries of  $U$  in the standard order.

- $T = \begin{array}{|c|c|c|} \hline & 1 & 3 \\ \hline 2 & 3 & \\ \hline \end{array}$ ,  $U = \begin{array}{|c|c|c|c|} \hline & 1 & 1 & 2 \\ \hline 2 & 2 & & \\ \hline \end{array}$   $stdU = \begin{array}{|c|c|c|c|} \hline & 1 & 2 & 5 \\ \hline 3 & 4 & & \\ \hline \end{array}$ .
- *internal insertion order word*  $R(U) = 12211 = R(stdU)$ .
- $P = \phi_{R(U)} T = \phi_{R(stdU)} T = \phi_{12211} T = \phi_1 \phi_2 \phi_2 \phi_1 \phi_1 T$ .

# $P$ -tableau in the Internal row insertion correspondence

- $\phi_{R(U)} = \phi_{R(\text{std}U)} = \phi_{12211} = \phi_1\phi_2\phi_2\phi_1\phi_1$ ,

$$\begin{aligned}
 P = \phi_{12211} T &= \phi_{12211} \begin{array}{|c|c|c|} \hline & 1 & 3 \\ \hline 2 & 3 & \\ \hline \end{array} = \phi_{1221} \begin{array}{|c|c|c|} \hline & & 3 \\ \hline 1 & 3 & \\ \hline 2 & & \\ \hline \end{array} = \phi_{122} \begin{array}{|c|c|c|} \hline & & \\ \hline 1 & 3 & 3 \\ \hline 2 & & \\ \hline \end{array} \\
 &= \phi_{12} \begin{array}{|c|c|c|} \hline & & \\ \hline & 3 & 3 \\ \hline 1 & & \\ \hline 2 & & \\ \hline \end{array} = \phi_1 \begin{array}{|c|c|c|} \hline & & \\ \hline & & 3 \\ \hline 1 & 3 & \\ \hline 2 & & \\ \hline \end{array} = \begin{array}{|c|c|c|c|} \hline & & & \\ \hline & & & 3 \\ \hline 1 & 3 & & \\ \hline 2 & & & \\ \hline \end{array}
 \end{aligned}$$

- Under which conditions do we have coincidence of the  $P$ -tableau?

$$U = \begin{array}{|c|c|} \hline & 1 \\ \hline 2 & \\ \hline \end{array} \text{ and } V = \begin{array}{|c|c|} \hline & 3 \\ \hline 1 & \\ \hline \end{array}, \phi_2\phi_1 T = \begin{array}{|c|c|c|} \hline & & 3 \\ \hline & 3 & \\ \hline 1 & & \\ \hline 2 & & \\ \hline \end{array} \neq \phi_1\phi_2 T = \begin{array}{|c|c|c|} \hline & & 3 \\ \hline & 1 & \\ \hline 2 & 3 & \\ \hline \end{array}.$$

$$U' = \begin{array}{|c|c|} \hline & 2 \\ \hline 1 & \\ \hline 3 & \\ \hline \end{array} \text{ and } V' = \begin{array}{|c|c|} \hline & 3 \\ \hline 1 & \\ \hline 2 & \\ \hline \end{array}, R(U') = 312 \equiv R(V') = 132, \text{ and}$$

$$\phi_{312} T = \phi_{132} T = \begin{array}{|c|c|c|} \hline & & 3 \\ \hline & 1 & \\ \hline & 3 & \\ \hline 2 & & \\ \hline \end{array}.$$

# The Knuth class of an internal insertion order word

- Elementary Knuth transformations on words:

$$kij \equiv kji = \begin{array}{|c|c|} \hline i & j \\ \hline k & \\ \hline \end{array}, \quad i < k \leq j, \quad ijk \equiv jik = \begin{array}{|c|c|} \hline i & k \\ \hline j & \\ \hline \end{array}, \quad i \leq k < j.$$

- Sufficient conditions for the coincidence of the  $P$ -tableau in the internal insertion correspondence.

## Theorem

*Knuth commutation of internal insertion operators* A. (2016). Let  $u$  be an internal insertion order word of  $T$  and  $v \equiv u$ . Then

- (a)  $v$  is an internal insertion order word of  $T$ .
- (c)  $\phi_u T = \phi_v T = P$ .

# Extension of Internal insertion operators

Sagan-Stanley internal insertion operator  $\bar{\phi}_i$  on  $Y \cup T$ ,  $Y$  Yamanouchi tableau,  $T$  a skew-tableau:

$$Y \cup T = \begin{array}{cccccc} 1 & 1 & 1 & 1 & 1 & 1 \\ 2 & 2 & 1 & 2 & 2 & \\ 3 & 2 & 3 & 3 & & \\ 1 & 3 & 4 & 4 & & \\ 5 & 5 & & & & \end{array}$$

$$\bar{\phi}_3(Y \cup T) = \begin{array}{cccccc} 1 & 1 & 1 & 1 & 1 & 1 \\ 2 & 2 & 1 & 2 & 2 & \\ 3 & 3 & 3 & 3 & & \\ 1 & 2 & 4 & 4 & & \\ 3 & 5 & & & & \\ 5 & & & & & \end{array}$$



# Internal insertion LR commutator A.'99; A., King, Terada'16

- $Y_{32} \cup T =$ 

1	1	1	1
2	2	1	2
1	2	3	

 $\quad \bar{\phi}_{12}\bar{\phi}_1$

- $\emptyset \xrightarrow{\chi_1^3 \omega_1^1}$ 

1	1	1	1
---	---	---	---

 $\xrightarrow{\chi_2^2 \bar{\phi}_1 \omega_2^1}$ 

1	1	1	1
2	1	2	2

 $\xrightarrow{\chi_3^0 \bar{\phi}_1 \bar{\phi}_2 \omega_3^1}$ 

1	1	1	1
2	2	1	2
3	1	2	

 $= Y_{321} \cup H$

$$\chi_3^0 \bar{\phi}_{12} \omega_3^1 [\chi_2^2 \bar{\phi}_1 \omega_2^1 [\chi_1^3 \omega_1^1 (\emptyset)]] = Y_{321} \cup H$$

# Internal insertion LR commutor A.'99; A., King, Terada'16

- $Y_{32} \cup T = \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 1 & 2 & 3 & \\ \hline \end{array} \quad \bar{\phi}_{12}\bar{\phi}_1$

- $\emptyset \xrightarrow{\chi_1^3\omega_1^1} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline \end{array} \xrightarrow{\chi_2^2\bar{\phi}_1\omega_2^1} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 1 & 2 & 2 \\ \hline \end{array} \xrightarrow{\chi_3^0\bar{\phi}_1\bar{\phi}_2\omega_3^1} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 3 & 1 & 2 & \\ \hline \end{array} = Y_{321} \cup H$

$$\chi_3^0\bar{\phi}_{12}\omega_3^1[\chi_2^2\bar{\phi}_1\omega_2^1[\chi_1^3\omega_1^1(\emptyset)]] = Y_{321} \cup H$$

- $(Y \cup T)^{(1)} = \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline \end{array} = \rho_1^{(1)}[(Y \cup T)^{(1)}] = \chi_1^3\omega_1^1\rho_1^{(0)}[\emptyset]$

# Internal insertion LR commutor A.'99; A., King, Terada'16

- $Y_{32} \cup T = \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 1 & 2 & 3 & \\ \hline \end{array} \quad \bar{\phi}_{12}\bar{\phi}_1$

- $\emptyset \xrightarrow{\chi_1^3 \omega_1^1} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline & & & \\ \hline & & & \\ \hline \end{array} \xrightarrow{\chi_2^2 \bar{\phi}_1 \omega_2^1} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 1 & 2 & 2 \\ \hline & & & \\ \hline \end{array} \xrightarrow{\chi_3^0 \bar{\phi}_1 \bar{\phi}_2 \omega_3^1} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 3 & 1 & 2 & \\ \hline \end{array} = Y_{321} \cup H$

$$\chi_3^0 \bar{\phi}_{12} \omega_3^1 [\chi_2^2 \bar{\phi}_1 \omega_2^1 [\chi_1^3 \omega_1^1 (\emptyset)]] = Y_{321} \cup H$$

- $(Y \cup T)^{(1)} = \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline & & & \\ \hline & & & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline & & & \\ \hline & & & \\ \hline \end{array} = \rho_1^{(1)}[(Y \cup T)^{(1)}] = \chi_1^3 \omega_1^1 \rho_1^{(0)}[\emptyset]$

$$\begin{aligned} (Y \cup T)^{(2)} &= \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline & & & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 1 & 2 & 2 \\ \hline & & & \\ \hline \end{array} = \rho_1^{(2)}[(Y \cup T)^{(2)}] \\ &= \chi_2^2 \bar{\phi}_1 \omega_2^1 \rho_1^{(1)}[(Y \cup T)^{(1)}] \end{aligned}$$

# Internal insertion LR commutor A.'99; A., King, Terada'16

- $Y_{32} \cup T = \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 1 & 2 & 3 & \\ \hline \end{array} \quad \bar{\phi}_{12}\bar{\phi}_1$

- $\emptyset \xrightarrow{\chi_1^3 \omega_1^1} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline & & & \\ \hline & & & \\ \hline \end{array} \xrightarrow{\chi_2^2 \bar{\phi}_1 \omega_2^1} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 1 & 2 & 2 \\ \hline & & & \\ \hline \end{array} \xrightarrow{\chi_3^0 \bar{\phi}_1 \bar{\phi}_2 \omega_3^1} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 3 & 1 & 2 & \\ \hline \end{array} = Y_{321} \cup H$

$$\chi_3^0 \bar{\phi}_{12} \omega_3^1 [\chi_2^2 \bar{\phi}_1 \omega_2^1 [\chi_1^3 \omega_1^1 (\emptyset)]] = Y_{321} \cup H$$

- $(Y \cup T)^{(1)} = \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline & & & \\ \hline & & & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline & & & \\ \hline & & & \\ \hline \end{array} = \rho_1^{(1)}[(Y \cup T)^{(1)}] = \chi_1^3 \omega_1^1 \rho_1^{(0)}[\emptyset]$

$$\begin{aligned} (Y \cup T)^{(2)} &= \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline & & & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 1 & 2 & 2 \\ \hline & & & \\ \hline \end{array} = \rho_1^{(2)}[(Y \cup T)^{(2)}] \\ &= \chi_2^2 \bar{\phi}_1 \omega_2^1 \rho_1^{(1)}[(Y \cup T)^{(1)}] \end{aligned}$$

$$\begin{aligned} (Y \cup T)^{(3)} &= \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 1 & 2 & 3 & \\ \hline \end{array} \xrightarrow{s} \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 1 & 2 \\ \hline 3 & 1 & 2 & \\ \hline \end{array} = Y_{321} \cup H = \rho_1^{(3)}[(Y \cup T)^{(3)}] \\ &= \chi_3^0 \bar{\phi}_{12} \omega_3^1 \rho_1^{(2)}[(Y \cup T)^{(2)}]. \end{aligned}$$

# Internal insertion and switching LR commutators coincide

## Theorem

(A., 2017)

Let  $n \geq 1$  and  $Y \cup T \in \mathcal{LR}^{(n)}$  with  $Y = Y_\mu$  Yamanouchi tableau and  $T$  a ballot tableau of shape  $\lambda/\mu$  and weight  $\nu$ . Consider the  $n$ th row word of  $T$  where  $F$  is the row subword restricted to the entries in  $[n-1]$ , and  $\nu_n$  is the number of entries equal to  $n$ . Then

$$\rho_1^{(n)}(Y \cup T) = \chi_n^{\mu_n} \bar{\phi}_F \omega_n^{\nu_n} \rho_1^{(n-1)}[(Y \cup T)^{(n-1)}] = \chi_n^{\mu_n} \omega_n^{\nu_n} \bar{\phi}_F \rho_1^{(n-1)}[(Y \cup T)^{(n-1)}],$$

with  $\rho_1^{(0)}(\emptyset) := \emptyset$ . In particular, all bumping routes of  $\bar{\phi}_F$  are pairwise disjoint and terminate in the  $n$ th row.

$(Y \cup T)^{(n-1)}$  consists of the first  $n-1$  rows of  $T$ .

## A bit more: Sketch of the proof (main steps)

- Reduce to the case  $\mu = (\mu_1, \dots, \mu_{n-1}, 0)$ . By induction on  $n \geq 1$  and  $|F| \geq 0$ .
- $n = 1$ , trivial,

$$\rho_1^{(1)}(\emptyset \cup Y_{(\nu_1)}) = \omega_1^{\nu_1} \emptyset = \omega_1^{\nu_1} \rho_1^{(0)}[\emptyset] = Y_{(\nu_1)} \cup \emptyset.$$

$$\emptyset \cup Y_{(\nu_1)} = \emptyset \cup \boxed{1111} \rightarrow \boxed{1111} \cup \emptyset = Y_{(\nu_1)} \cup \emptyset$$

## Sketch of the proof

- Assume the statement true for  $n$  and prove for  $n + 1$ . Detach the  $n + 1$ th row

$$Y \cup T = F(n + 1)^{\nu_{n+1}} * (Y \cup T)^{(n)},$$

$F$  a word on the alphabet  $[n]$ .

We claim

$$\rho_1^{(n+1)}(Y \cup T) = \omega_{n+1}^{\nu_{n+1}} \bar{\phi}_F \rho_1^{(n)}[(Y \cup T)^{(n)}].$$

- ▶  $|F| = 0$ . Trivial. Switching on  $Y \cup T$  reduces to  $(Y \cup T)^{(n)}$ , the first  $n$  rows of  $Y \cup T$ , and  $\bar{\phi}_F = id$ ,

$$\rho_1^{(n+1)}(Y \cup T) = \omega_{n+1}^{\nu_{n+1}} \rho_1^{(n)}[(Y \cup T)^{(n)}].$$

## Sketch of the proof

- Let  $|F| \geq 1$ . Apply switching to decompose  $Y \cup T$  and reduce  $|F|$

$$Y \cup T \xrightarrow{s} [Y' \cup S] \cup [D * Q^{(n)}] = [\widehat{F}(n+1)^{\nu_{n+1}} * (Y' \cup S)^{(n)}] \cup [D * Q^{(n)}],$$

$Y' = Y_{(\mu_1, \dots, \mu_{d-1})}$ ,  $\mu_d > 0$ ,  $S \equiv T$ ,  $D * Q^{(n)} \equiv Y_{(\mu_d, \dots, \mu_{n-1})}$ ,  $\widehat{F}$  strict subword of  $F$  and  $D = d^{|D|}$ ,  $|D| > 0$ ,

$$|\widehat{F}| + |D| = |F|.$$

- $0 \leq |\widehat{F}| < |F|$ . By induction on  $|F|$ ,

$$\rho_1^{(n+1)}(Y' \cup S) = \omega_{n+1}^{\nu_{n+1}} \bar{\phi}_{\widehat{F}} \rho_1^{(n)}[(Y' \cup S)^{(n)}],$$

all  $\bar{\phi}_{\widehat{F}}$ -bumping routes terminate in the  $(n+1)$ th row.



- One has so far

$$\rho_1^{(n+1)}(Y \cup T) = \omega_{n+1}^{\nu_{n+1}} \bar{\phi}_{\widehat{F}} \rho_1^{(n)}[(Y' \cup S)^{(n)}] \cup (D * Q^{(n)}),$$

all  $\bar{\phi}_{\widehat{F}}$ -bumping routes terminate in the  $(n+1)$ th row.

- We claim

$$\rho_1^{(n+1)}(Y \cup T) = \omega_{n+1}^{\nu_{n+1}} \bar{\phi}_F \rho_1^{(n)}[(Y \cup T)^{(n)}].$$

- $Y \cup T \xrightarrow{s} [Y' \cup S] \cup [D * Q^{(n)}] = [\widehat{F}(n+1)^{\nu_{n+1}} * (Y' \cup S)^{(n)}] \cup [D * Q^{(n)}]$ .

Detach the  $n$ th row

$$(Y' \cup S)^{(n)} = G C n^{\hat{\nu}_n} * (Y' \cup S)^{(n-1)},$$

$G$  and  $C$  row words in the alphabet  $[n-1]$ ,  $|G| = |\widehat{F}| + |D| = |F|$ .

$$[\widehat{F}(n+1)^{\nu_{n+1}} \cup D] * [G C n^{\hat{\nu}_n} \cup X] =$$

A	B	C	n*	...	n*	X
$\widehat{F}$	n+1*	...	n+1*	D		

 $\xleftrightarrow{s}$

G <sub>1</sub>	F <sub>2</sub>	G <sub>3</sub>	F <sub>4</sub>	...	G <sub>k-1</sub>	F <sub>k</sub>	C	n*	...	n*	X
F <sub>1</sub>	D <sub>2</sub>	F <sub>3</sub>	D <sub>4</sub>	...	F <sub>k-1</sub>	D <sub>k</sub>	n+1*	...	n+1*		

 $\xleftrightarrow{s}$

G <sub>1</sub>	D <sub>2</sub>	G <sub>3</sub>	D <sub>4</sub>	...	G <sub>k-1</sub>	D <sub>k</sub>	C	n*	...	n*	X
F <sub>1</sub>	F <sub>2</sub>	F <sub>3</sub>	F <sub>4</sub>	...	F <sub>k-1</sub>	F <sub>k</sub>	n+1*	...	n+1*		

$\widehat{G} := G_1 G_3 \dots G_{k-1}$  a subword of  $G := AB$ ,  $|\widehat{G}| = |\widehat{F}|$ , and  $\widehat{F}\widehat{G} \equiv F\widehat{G}$  Knuth equivalent.

- $Y \cup T \xrightarrow{s} [Y' \cup S] \cup [D * Q^{(n)}] = [\widehat{F}(n+1)^{\nu_{n+1}} * (Y' \cup S)^{(n)}] \cup [D * Q^{(n)}]$ .  
 Detach the  $n$ th row

$$(Y' \cup S)^{(n)} = G C n^{\hat{\nu}_n} * (Y' \cup S)^{(n-1)},$$

$G$  and  $C$  row words in the alphabet  $[n-1]$ ,  $|G| = |\widehat{F}| + |D| = |F|$ .

$$[\widehat{F}(n+1)^{\nu_{n+1}} \cup D] * [G C n^{\hat{\nu}_n} \cup X] =$$

A	B	C	n*	...	n*	X
$\widehat{F}$	n+1*	...	n+1*	D		

 $\xleftrightarrow{s}$

G <sub>1</sub>	F <sub>2</sub>	G <sub>3</sub>	F <sub>4</sub>	...	G <sub>k-1</sub>	F <sub>k</sub>	C	n*	...	n*	X
F <sub>1</sub>	D <sub>2</sub>	F <sub>3</sub>	D <sub>4</sub>	...	F <sub>k-1</sub>	D <sub>k</sub>	n+1*	...	n+1*		

 $\xleftrightarrow{s}$

G <sub>1</sub>	D <sub>2</sub>	G <sub>3</sub>	D <sub>4</sub>	...	G <sub>k-1</sub>	D <sub>k</sub>	C	n*	...	n*	X
F <sub>1</sub>	F <sub>2</sub>	F <sub>3</sub>	F <sub>4</sub>	...	F <sub>k-1</sub>	F <sub>k</sub>	n+1*	...	n+1*		

$\widehat{G} := G_1 G_3 \dots G_{k-1}$  a subword of  $G := AB$ ,  $|\widehat{G}| = |\widehat{F}|$ , and  
 $\widehat{F}\widehat{G} \equiv F\widehat{G}$  Knuth equivalent.

- $(Y \cup T)^{(n)} \xrightarrow{s} [\widehat{G} C n^{\hat{\nu}_n} * (Y' \cup S)^{(n-1)}] \cup (DX * Q^{(n-1)})$

- By induction on  $n$



$$\begin{aligned} \rho_1^{(n)}(Y' \cup S)^{(n)} &= \rho_1^{(n)}[GCn^{\hat{\nu}_n} * (Y' \cup S)^{(n-1)}] \\ &= \bar{\phi}_G \rho_1^{(n)}(Cn^{\hat{\nu}_n} * (Y' \cup S)^{(n-1)}), \end{aligned}$$

all  $\bar{\phi}_G$ -bumping routes terminate in the  $n$ th row.



$$\begin{aligned} \rho_1^{(n)}[(Y \cup T)^{(n)}] &= \rho_1^{(n)}[\widehat{G}Cn^{\hat{\nu}_n} * (Y' \cup S)^{(n-1)}] \cup (DX * Q^{(n-1)}) \\ &= \bar{\phi}_{\widehat{G}} \rho_1^{(n)}[Cn^{\hat{\nu}_n} * (Y' \cup S)^{(n-1)}] \cup (DX * Q^{(n-1)}). \end{aligned}$$

all  $\bar{\phi}_{\widehat{G}}$ -bumping routes terminate in the  $n$ th row.

$$\widehat{FG} \equiv F\widehat{G} \Rightarrow \bar{\phi}_{\widehat{FG}} = \bar{\phi}_{F\widehat{G}}.$$

$$\rho_1^{(n)}[(Y' \cup S)^{(n)}] = \bar{\phi}_{G}\rho_1^{(n)}[Cn^{\hat{\nu}_n} * (Y' \cup S)^{(n-1)}]$$

$$\rho_1^{(n)}[(Y \cup T)^{(n)}] = \bar{\phi}_{\widehat{G}}\rho_1^{(n)}[Cn^{\hat{\nu}_n} * (Y' \cup S)^{(n-1)}] \cup (DX * Q^{(n-1)})$$

$$\begin{aligned} \rho_1^{(n+1)}(Y \cup T) &= \omega_{n+1}^{\nu_{n+1}} \bar{\phi}_{\widehat{F}} \rho_1^{(n)}[(Y' \cup S)^{(n)}] \cup (D * Q^{(n)}), \\ &= \omega_{n+1}^{\nu_{n+1}} \bar{\phi}_{\widehat{F}} \bar{\phi}_G \rho_1^{(n)}[Cn^{\hat{\nu}_n} * (Y' \cup S)^{(n-1)}] \cup (D * Q^{(n)}) \\ &= \omega_{n+1}^{\nu_{n+1}} \bar{\phi}_F \bar{\phi}_{\widehat{G}} \rho_1^{(n)}[Cn^{\hat{\nu}_n} * (Y' \cup S)^{(n-1)}] \cup (D * X * Q^{(n-1)}) \\ &= \omega_{n+1}^{\nu_{n+1}} \bar{\phi}_F [\bar{\phi}_{\widehat{G}} \rho_1^{(n)}(Cn^{\hat{\nu}_n} * (Y' \cup S)^{(n-1)} \cup (DX * Q^{(n-1)}))] \\ &= \omega_{n+1}^{\nu_{n+1}} \bar{\phi}_F \rho_1^{(n)}[(Y \cup T)^{(n)}]. \end{aligned}$$