

Geometric Theorem Proving

Pedro Quaresma

CISUC, Mathematics Department
University of Coimbra

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AI (synthetic) Methods

Synthetic methods attempt to automate traditional geometry proof methods that produce human-readable proofs.

In 1950s Gelernter created a theorem prover that could find solutions to a number of problems taken from high-school textbooks in plane geometry [Gel59].

It was based on the human simulation approach and has been considered a landmark in the AI area for this time.

In spite of the success and significant improvements [BdC95, CP79, CP86, Gil70, KA90, Nev74, Qua89] with these methods, the results did not lead to the development of a powerful geometry theorem prover.



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Coordinate-free Methods

Instead of coordinates, three basic geometric quantities: the ratio of parallel line segments, the signed area, and the Pythagorean difference.

- ▶ Area method [CGZ93, JNQ11, QJ06a];
- ▶ Full angle method [CGZ94, CGZ96a];
- ▶ Solid geometry [CGZ95].

Geometric proofs, small and human-readable.

But:

- ▶ not the “normal” high-school geometric reasoning;
- ▶ for many conjectures these methods still deal with extremely complex expressions involving certain geometric quantities.



Other approaches

- ▶ An approach based on a deductive database and forward chaining works over a suitably selected set of higher-order lemmas and can prove complex geometry theorems, but still has a smaller scope than algebraic provers [CGZ94, CGZ00].
- ▶ Quaife used a resolution theorem prover to prove theorems in Tarski's geometry [Qua89].
- ▶ A GATP based on coherent-logic capable of producing both readable and formal proofs of geometric conjectures of certain sort [SPJ10].
- ▶ Probabilistic verification of elementary geometry statements [CFGG97, RGK99].
- ▶ Visual Reasoning/Proofs [Kim89, YCG10a, YCG10b].



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Mechanical Geometric Formula Derivation

- ▶ Locus Generation: to determine the implicit equation of a locus set [BAE07, BA12].

The set of points determined by the different positions of a point, the tracer, as a second point in which the tracer depends on, called the mover, runs along the one dimensional object to which it is restrained.

- ▶ Deriving Geometry Formulas: automatic derivation of geometry formulas [cCsG90, KSY94, RV99].

Example: find the formula for the area of a triangle ABC in terms of its three sides.



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Geometric Tools & Integration Issues

Geometric tools: DGS & GATP & CAS & RGP.

- ▶ DGS - Dynamic Geometry Software [Gro11, Hoh02, Jac01, Jan06, RGK99] - “visual proofs” [CGY04];
- ▶ GATP - Geometry Automated Theorem Provers
 - ▶ verification of the soundness of a geometric construction [JQ07].
 - ▶ reason about a given DGS construction [CGZ96b, JQ06, Nar07a, QP06, QJ06b].
 - ▶ human-readable proofs [JNQ11, QJ06a, QJ09].
- ▶ RGP - Repositories of Geometric Problems [QJ07, Qua11].
- ▶ eLearning [ABY86, HLY86, QJ06b, SQ08, SQ10, SQ12]

Integration: Intergeo Project [SHK⁺10] — Deduction STREP Proposal [WSA⁺12].



Geometric Tools & Integration Issues

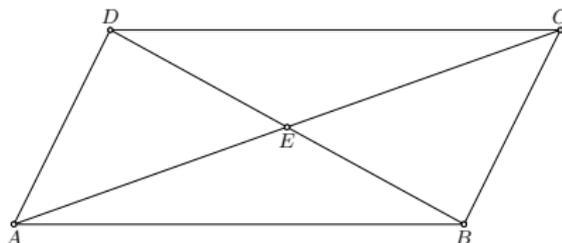
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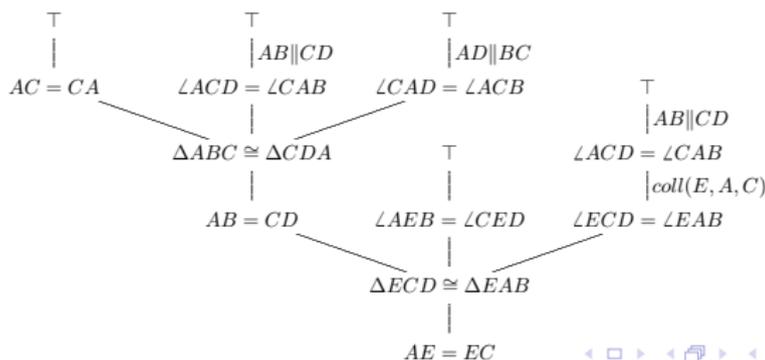
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Example 1 - Gelernter

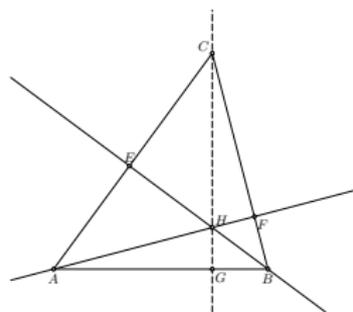


$$points(A, B, C) \wedge AB \parallel CD \wedge AD \parallel BC \wedge coll(E, A, C) \wedge coll(E, B, D) \Rightarrow AE = EC$$



Geometry Deductive Database - The Orthocenter Theorem

$points(A, B, C) \wedge coll(E, A, C) \wedge perp(B, E, A, C) \wedge coll(F, B, C) \wedge$
 $perp(A, F, B, C) \wedge coll(H, A, F) \wedge coll(H, B, E) \wedge coll(G, A, B) \wedge$
 $coll(G, C, H)$



The fix-point contains two of the most often encountered properties of this configuration:

- ▶ $perp(C, G, A, B)$;
- ▶ $\angle FGC = \angle CGE$



Quaife's GATP

Tarski axiomatic system: is, or rather its algebraic equivalent, complete and decidable.

Quaife developed a GATP for Euclidean plane geometry within the automated reasoning system OTTER (a resolution theorem prover) [Qua89].

(A1) Reflexivity axiom for equidistance.

$$\rightarrow u \cdot v \equiv v \cdot u$$

(A2) Transitivity axiom for equidistance.

$$u \cdot v \equiv w \cdot x, u \cdot v \equiv y \cdot z \rightarrow w \cdot x \equiv y \cdot z$$

(A4) Segment construction axiom, two clauses.

$$(A4.1) \rightarrow B(u, v, Ext(u, v, w, x))$$

$$(A4.2) \rightarrow v \cdot Ext(u, v, w, x) \equiv w \cdot x$$

(...)



Quaife's GATP

Heuristics

- ▶ maximum weight for retained clauses at 25,
- ▶ first attempt to obtain a proof in which no variables are allowed in any generated and retained clause.

The provers based upon Wu's algorithm, are able to prove quite more difficult theorems in geometry those by Quaife's GATP.

However Wu's method only works with hypotheses and theorems that can be expressed as equations, and not with inequalities as correspond to the relation B in Quaife's resolution prover.



Visual Reasoning/Representation

Visual Reasoning extend the use of diagrams with a method that allows the diagrams to be perceived and to be manipulated in a creative manner [Kim89].

Visually Dynamic Presentation of Proofs linking the proof done by a synthetic method (full-angle) with a visual presentation of the proof [YCG10a, YCG10b].



Wu's Method

An elementary version of Wu's method is simple:
Geometric theorem T transcribed as polynomial equations and inequations of the form:

- ▶ H: $h_1 = 0, \dots, h_s = 0, d_1 \neq 0, \dots, d_t \neq 0$;
- ▶ C: $c=0$.

Proving T is equivalent to deciding whether the formula

$$\forall_{x_1, \dots, x_n} [h_1 = 0 \wedge \dots \wedge h_s \wedge d_1 \neq 0 \wedge \dots \wedge d_t \neq 0 \Rightarrow c = 0] \quad (1)$$

is valid.



Wu's Method

Computes a characteristic set C of $\{h_1, \dots, h_s\}$ and the pseudo-remainder r of c with respect to C .

If r is identically equal to 0, then T is proved to be true.

The subsidiary condition $J \neq 0$, where J is the product of initials of the polynomials in C are the ndg conditions [WT86, Wu00].

This is a decision procedure.



Area Method — Basic Geometric Quantities¹

Definition (Ratio of directed parallel segments)

For four collinear points P , Q , A , and B , such that $A \neq B$, the ratio of directed parallel segments, denoted $\frac{\overline{PQ}}{\overline{AB}}$ is a real number.

Definition (Signed Area)

The *signed area* of triangle ABC , denoted S_{ABC} , is the area of the triangle with a sign depending on its orientation in the plane.

Definition (Pythagoras difference)

For three points A , B , and C , the *Pythagoras difference*, is defined in the following way: $\mathcal{P}_{ABC} = \overline{AB}^2 + \overline{CB}^2 - \overline{AC}^2$.

¹ [JNQ11, QJ06a, QJ09]

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Properties of the Ratio of Directed Parallel Segments

$$\blacktriangleright \frac{\overline{PQ}}{\overline{AB}} = -\frac{\overline{QP}}{\overline{AB}} = \frac{\overline{QP}}{\overline{BA}} = -\frac{\overline{PQ}}{\overline{BA}};$$

$$\blacktriangleright \frac{\overline{PQ}}{\overline{AB}} = 0 \text{ iff } P = Q;$$

▶ (...)

EL1 (The Co-side Theorem) Let M be the intersection of two non-parallel lines AB and PQ and $Q \neq M$. Then it holds that

$$\frac{\overline{PM}}{\overline{QM}} = \frac{S_{PAB}}{S_{QAB}}; \quad \frac{\overline{PM}}{\overline{PQ}} = \frac{S_{PAB}}{S_{PAQB}}; \quad \frac{\overline{QM}}{\overline{PQ}} = \frac{S_{QAB}}{S_{PAQB}}.$$



Properties of the Signed Area

- ▶ $S_{ABC} = S_{CAB} = S_{BCA} = -S_{ACB} = -S_{BAC} = -S_{CBA}$.
- ▶ $S_{ABC} = 0$ iff A , B , and C are collinear.
- ▶ $PQ \parallel AB$ iff $S_{PAB} = S_{QAB}$, i.e., iff $S_{PAQB} = 0$.
- ▶ Let $ABCD$ be a parallelogram, P and Q be two arbitrary points. Then it holds that $S_{APQ} + S_{CPQ} = S_{BPQ} + S_{DPQ}$ or $S_{PAQB} = S_{PDQC}$.
- ▶ Let R be a point on the line PQ . Then for any two points A and B it holds that $S_{RAB} = \frac{PR}{PQ} S_{QAB} + \frac{RQ}{PQ} S_{PAB}$.
- ▶ (...)



Properties of the Pythagoras Difference

- ▶ $\mathcal{P}_{AAB} = 0$.
- ▶ $\mathcal{P}_{ABC} = \mathcal{P}_{CBA}$.
- ▶ If A , B , and C are collinear then, $\mathcal{P}_{ABC} = 2\overline{BA} \overline{BC}$.
- ▶ $AB \perp BC$ iff $\mathcal{P}_{ABC} = 0$.
- ▶ Let AB and PQ be two non-perpendicular lines, and Y be the intersection of line PQ and the line passing through A and perpendicular to AB . Then, it holds that

$$\frac{\overline{PY}}{\overline{QY}} = \frac{\mathcal{P}_{PAB}}{\mathcal{P}_{QAB}}, \quad \frac{\overline{PY}}{\overline{PQ}} = \frac{\mathcal{P}_{PAB}}{\mathcal{P}_{PAQB}}, \quad \frac{\overline{QY}}{\overline{PQ}} = \frac{\mathcal{P}_{QAB}}{\mathcal{P}_{PAQB}}.$$

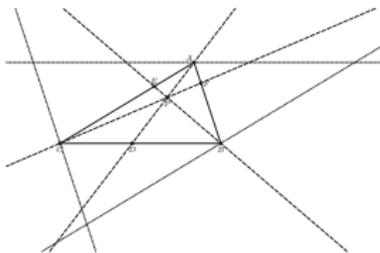
- ▶ (...)



The Area Method — The Proof Algorithm

Express the hypothesis of a theorem using a set of constructive statements.

Each constructive statement introduces a new point.



The conclusion is expressed by a polynomial in some geometry quantities (defined above), without any relation to a given system of coordinates.

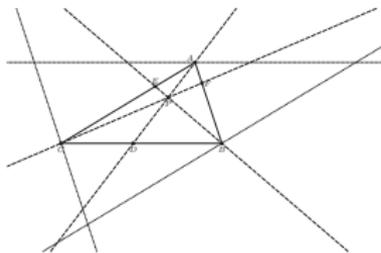
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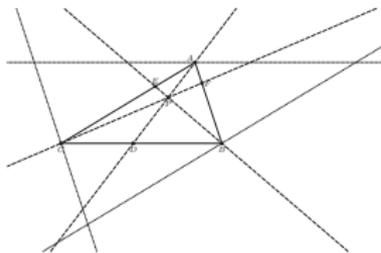
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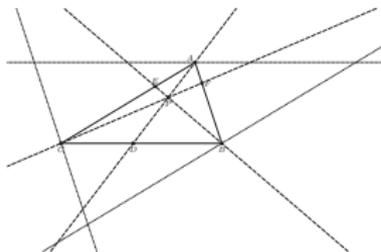
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Constructive Geometric Statements

ECS1 construction of an arbitrary point U ; (...).

ECS2 construction of a point Y such that it is the intersection of two lines (LINE $U V$) and (LINE $P Q$);
ndg-condition: $UV \nparallel PQ$; $U \neq V$; $P \neq Q$.
degree of freedom for Y : 0

ECS3 construction of a point Y such that it is a foot from a given point P to (LINE $U V$); (...).

ECS4 construction of a point Y on the line passing through point W and parallel to (LINE $U V$), such that $\overline{WY} = r\overline{UV}$, (...).

ECS5 construction of a point Y on the line passing through point U and perpendicular to (LINE $U V$), such that $r = \frac{4S_{UVY}}{P_{UVU}}$, (...).



Forms of Expressing the Conclusion

property	in terms of geometric quantities
points A and B are identical	$\mathcal{P}_{ABA} = 0$
points A, B, C are collinear	$\mathcal{S}_{ABC} = 0$
AB is perpendicular to CD	$\mathcal{P}_{ABA} \neq 0 \wedge \mathcal{P}_{CDC} \neq 0 \wedge \mathcal{P}_{ACD} = \mathcal{P}_{BCD}$
AB is parallel to CD	$\mathcal{P}_{ABA} \neq 0 \wedge \mathcal{P}_{CDC} \neq 0 \wedge \mathcal{S}_{ACD} = \mathcal{S}_{BCD}$
O is the midpoint of AB	$\mathcal{S}_{ABO} = 0 \wedge \mathcal{P}_{ABA} \neq 0 \wedge \frac{\overline{AO}}{\overline{AB}} = \frac{1}{2}$
AB has the same length as CD	$\mathcal{P}_{ABA} = \mathcal{P}_{CDC}$
points A, B, C, D are harmonic	$\mathcal{S}_{ABC} = 0 \wedge \mathcal{S}_{ABD} = 0 \wedge \mathcal{P}_{BCB} \neq 0 \wedge \mathcal{P}_{BDB} \neq 0 \wedge \frac{\overline{AC}}{\overline{CB}} = \frac{\overline{DA}}{\overline{DB}}$
angle ABC has the same measure as DEF	$\mathcal{P}_{ABA} \neq 0 \wedge \mathcal{P}_{ACA} \neq 0 \wedge \mathcal{P}_{BCB} \neq 0 \wedge \mathcal{P}_{DED} \neq 0 \wedge \mathcal{P}_{DFD} \neq 0 \wedge \mathcal{P}_{EFE} \neq 0 \wedge \mathcal{S}_{ABC} \cdot \mathcal{P}_{DEF} = \mathcal{S}_{DEF} \cdot \mathcal{P}_{ABC}$
A and B belong to the same circle arc CD	$\mathcal{S}_{ACD} \neq 0 \wedge \mathcal{S}_{BCD} \neq 0 \wedge \mathcal{S}_{CAD} \cdot \mathcal{P}_{CBD} = \mathcal{S}_{CBD} \cdot \mathcal{P}_{CAD}$



Elimination Lemmas

EL2 Let $G(Y)$ be a linear geometric quantity and point Y is introduced by the construction $(\text{PRATIO } Y \ W \ (\text{LINE } U \ V) \ r)$. Then it holds

$$G(Y) = G(W) + r(G(V) - G(U)).$$

EL3 Let $G(Y)$ be a linear geometric quantity and point Y is introduced by the construction $(\text{INTER } Y \ (\text{LINE } U \ V) \ (\text{LINE } P \ Q))$. Then it holds

$$G(Y) = \frac{S_{UPQ}G(V) - S_{VPQ}G(U)}{S_{UPVQ}}.$$

► (...)



Constructive Steps & Elimination Lemmas

		Geometric Quantities					
		\mathcal{P}_{AYB}	\mathcal{P}_{ABY}	\mathcal{P}_{ABCY}	\mathcal{S}_{ABY}	\mathcal{S}_{ABCY}	$\frac{\overline{AY}}{\overline{CD}}$
Constructive Steps	ECS2	EL5	EL3			EL11	EL1
	ECS3	EL6	EL4			EL12	
	ECS4	EL7	EL2			EL13	
	ECS5	EL10	EL9	EL8		EL14	
		Elimination Lemmas					



The Algorithm

- $S = (C_1, C_2, \dots, C_m, (E, F))$ is a statement in \mathbf{C} .
- ← The algorithm tells whether S is true, or not, and if it is true, produces a proof for S .

```

for (i=m;i==1;i--) {
  if (the ndg conditions of Ci is satisfied) exit;
  // Let G1,\ldots,Gn be the geometric quantities in E and F
  for (j=1;j<=n,j++) {
    Hj <- eliminating the point introduced
                by construction Ci from Gj
    E <- E[Gj:=Hj]
    F <- F[Gj:=Hj]
  }
}
if (E==F) S <- true else S<-false

```

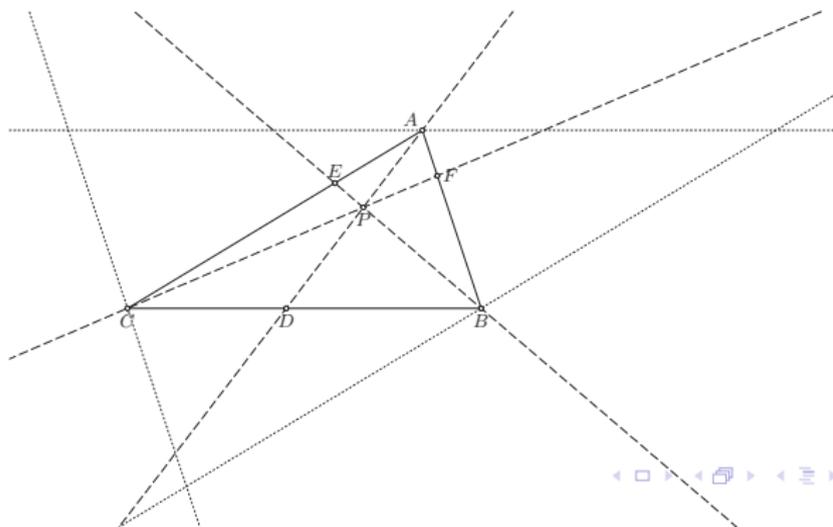
Adding to that it is needed to check the ndg condition of a construction (three possible forms).



An Example (Ceva's Theorem)

Let $\triangle ABC$ be a triangle and P be an arbitrary point in the plane. Let D be the intersection of AP and BC , E be the intersection of BP and AC , and F the intersection of CP and AB . Then:

$$\frac{\overline{AF}}{\overline{FB}} \frac{\overline{BD}}{\overline{DC}} \frac{\overline{CE}}{\overline{EA}} = 1$$



Example — Proof

The proof of a conjecture is based on eliminating all the constructed points, in reverse order, until an equality in only the free points is reached.

$$\begin{aligned}
 \frac{\overline{AF}}{\overline{FB}} \frac{\overline{BD}}{\overline{DC}} \frac{\overline{CE}}{\overline{EA}} &= \frac{S_{APC}}{S_{BCP}} \frac{\overline{BD}}{\overline{DC}} \frac{\overline{CE}}{\overline{EA}} && \text{the point } F \text{ is eliminated} \\
 &= \frac{S_{APC}}{S_{BCP}} \frac{S_{BPA}}{S_{CAP}} \frac{\overline{CE}}{\overline{EA}} && \text{the point } D \text{ is eliminated} \\
 &= \frac{S_{APC}}{S_{BCP}} \frac{S_{BPA}}{S_{CAP}} \frac{S_{CPB}}{S_{ABP}} && \text{the point } E \text{ is eliminated} \\
 &= 1
 \end{aligned}$$

Elimination Steps: 3; Geometric Steps: 6; Algebraic Steps: 23;
Total Steps: 32; CPU Time 0.004s. The GATP also provide the
ndg conditions.



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ndg conditions.



Example — Proof

The proof of a conjecture is based on eliminating all the constructed points, in reverse order, until an equality in only the free points is reached.

$$\begin{aligned} \frac{\overline{AF}}{\overline{FB}} \frac{\overline{BD}}{\overline{DC}} \frac{\overline{CE}}{\overline{EA}} &= \frac{S_{APC}}{S_{BCP}} \frac{\overline{BD}}{\overline{DC}} \frac{\overline{CE}}{\overline{EA}} && \text{the point } F \text{ is eliminated} \\ &= \frac{S_{APC}}{S_{BCP}} \frac{S_{BPA}}{S_{CAP}} \frac{\overline{CE}}{\overline{EA}} && \text{the point } D \text{ is eliminated} \\ &= \frac{S_{APC}}{S_{BCP}} \frac{S_{BPA}}{S_{CAP}} \frac{S_{CPB}}{S_{ABP}} && \text{the point } E \text{ is eliminated} \\ &= 1 \end{aligned}$$

Elimination Steps: 3; Geometric Steps: 6; Algebraic Steps: 23;
Total Steps: 32; CPU Time 0.004s. The GATP also provide the
ndg conditions.



Area Method - Formalisation

Formalisation [JNQ11, Nar06, Nar09];

1. $\overline{AB} = 0$ if and only if the points A and B are identical
2. $S_{ABC} = S_{CAB}$
3. $S_{ABC} = -S_{BAC}$
4. If $S_{ABC} = 0$ then $\overline{AB} + \overline{BC} = \overline{AC}$ (Chasles's axiom)
5. There are points A, B, C such that $S_{ABC} \neq 0$ (dimension; not all points are collinear)
6. $S_{ABC} = S_{DBC} + S_{ADC} + S_{ABD}$ (dimension; all points are in the same plane)
7. For each element r of F , there exists a point P , such that $S_{ABP} = 0$ and $\overline{AP} = r\overline{AB}$ (construction of a point on the line)
8. If $A \neq B$, $S_{ABP} = 0$, $\overline{AP} = r\overline{AB}$, $S_{ABP'} = 0$ and $\overline{AP'} = r\overline{AB}$, then $P = P'$ (unicity)
9. If $PQ \parallel CD$ and $\frac{PQ}{CD} = 1$ then $DQ \parallel PC$ (parallelogram)
10. If $S_{PAC} \neq 0$ and $S_{ABC} = 0$ then $\frac{\overline{AB}}{\overline{AC}} = \frac{S_{PAB}}{S_{PAC}}$ (proportions)
11. If $C \neq D$ and $AB \perp CD$ and $EF \perp CD$ then $AB \parallel EF$
12. If $A \neq B$ and $AB \perp CD$ and $AB \parallel EF$ then $EF \perp CD$
13. If $FA \perp BC$ and $S_{FBC} = 0$ then $4S_{ABC}^2 = \overline{AF}^2 \overline{BC}^2$ (area of a triangle)

Using this axiom system all the properties of the geometric quantities required by the area method were *formally verified* (within the *Coq* proof assistant [Tea09]), demonstrating the correctness of the system and eliminating all concerns about provability of the lemmas [Nar09].



Full Angle Method/Solid Geometry

Full Angle Method Full Angle is defined as an ordered pair of line which satisfies the following rules (...) [CGZ96a].

Solid Geometry Method For any points A , B , C and D in the space, the signed volume V_{ABCD} of the tetrahedron $ABCD$ is a real number which satisfies the following properties (...) [CGZ95].



Coherent Logic GATP

Coherent Logic is a fragment of first-order logic with formulae of the following form:

$$A_1(x) \wedge \dots \wedge A_n(x) \rightarrow \exists_{y_1} B(x, y_1) \vee \dots \vee \exists_{y_m} B(x, y_m)$$

with a breath-first proof procedure sound and complete [BC05].

ArgoCLP (**C**oherent **L**ogic **P**rover of the Argo Group²)

- ▶ new proof procedures;
- ▶ proof trace exportable to:
 - ▶ a proof object in Isabelle/Isar;
 - ▶ human readable (English/L^AT_EX).

not aimed at proving complex geometry theorems but rather at proving foundational theorems (close to the axiom level) [SPJ10].

²<http://argo.matf.bg.ac.rs/>



Repositories of Geometric Problems

GeoThms — a Web-based framework for exploring geometrical knowledge that integrates Dynamic Geometry Software (DGS), Automatic Theorem Provers (ATP), and a repository of geometrical constructions, figures and proofs. [JQ06, QJ07].

TGTP — a Web-based library of problems in geometry to support the testing and evaluation of geometric automated theorem proving (GATP) systems [Qua11].



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