

A pair of monads

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Motivation

Adjunctions and Monads

Definition

A *monad* \mathbb{T} on a category \mathbf{C} is a triple (T, m, e) , where $m : TT \rightarrow T$ and $e : 1 \rightarrow T$ are natural transformations satisfying the identities

$$\begin{array}{ccc} TTT & \xrightarrow{Tm} & TT \\ mT \downarrow & & \downarrow m \\ TT & \xrightarrow{m} & T \end{array}$$

$$\begin{array}{ccccc} T & \xrightarrow{eT} & TT & \xleftarrow{Te} & T \\ & \searrow & \downarrow & \swarrow & \\ & 1 & T & 1 & \end{array}$$

Example

If $(L \dashv R, \eta, \epsilon) : \mathbf{C} \rightarrow \mathbf{B}$ is an adjunction, then $(RL, R\epsilon L, \eta)$ is a monad on \mathbf{C} , and $(LR, L\eta R, \epsilon)$ is a comonad on \mathbf{B} .

Definition

A \mathbb{T} -algebra is a pair (X, a) , where $X \in \mathbf{C}$ and $a : TX \rightarrow X$ a morphism such that

$$a \cdot Ta = a \cdot m_X \text{ and } a \cdot e_X = 1_X.$$

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Definition

If (X, a) and (Y, b) are \mathbb{T} -algebras, then a \mathbb{T} -algebra homomorphism $f : (X, a) \rightarrow (Y, b)$ is a morphism $f : X \rightarrow Y$ in \mathbf{C} such that $b \cdot Tf = f \cdot a$.

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Definition

The category of \mathbb{T} -algebras and \mathbb{T} -algebra homomorphisms are denoted by $\mathbf{C}^{\mathbb{T}}$ or by $\mathbf{Alg}(\mathbb{T})$. The forgetful functor $G^{\mathbb{T}} : \mathbf{C}^{\mathbb{T}} \rightarrow \mathbf{C} : (X, a) \mapsto X$ admits a left adjoint $F^{\mathbb{T}} : \mathbf{C} \rightarrow \mathbf{C}^{\mathbb{T}} : X \mapsto (TX, m_X), f \mapsto Tf$.

A couple of triples (H. Simmons, 1982)

The adjunction

$$\mathbf{Top} \begin{array}{c} \xrightarrow{O} \\ \xleftarrow[\sigma]{\perp} \end{array} \mathbf{DLat}^{op}$$

where $\sigma(L) = \mathbf{DLat}(L, 2)$ and $OX = \mathbf{Top}(X, 2)$, induces two monads

$$M : \mathbf{Top} \rightarrow \mathbf{Top} \text{ and } \mathfrak{J} : \mathbf{DLat} \rightarrow \mathbf{DLat}.$$

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$$M : \mathbf{Top} \rightarrow \mathbf{Top} \text{ and } \mathfrak{I} : \mathbf{DLat} \rightarrow \mathbf{DLat}.$$

- $MX = \{\mathcal{F} \mid \mathcal{F} \text{ open prime filter on } X\}.$
- $\mathfrak{I}L \cong \{I \mid I \text{ ideal on } L\}.$

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Definition

A subset $J \subseteq D$ is an ideal if

- For any $a \in D$ and $i \in J$, if $a \leq i$ then $a \in J$.
- $0 \in J$ and for any $a, b \in J$, $a \vee b \in J$.

Distributive lattices and frames

Remark

The unit and multiplication that come with $\mathfrak{J} : \mathbf{DLat} \rightarrow \mathbf{Dlat}$ are given by

$$\downarrow : D \rightarrow \mathfrak{J}D \text{ and } \bigcup : \mathfrak{J}\mathfrak{J}D \rightarrow \mathfrak{J}D$$

Definition

A distributive lattice L is a frame if and only if the following equation holds: $a \wedge \bigvee S = \bigvee \{a \wedge s \mid s \in S\}$.

Distributive lattices and frames

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Definition

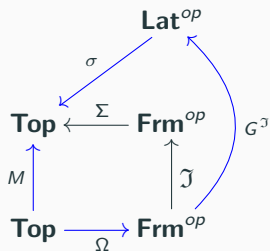
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Theorem

1. (H. Simmons, 1982) $\mathbf{Alg}(M) = \mathbf{StKSp}$.
2. (H. Simmons, 1982 - P. Johnstone, 1982) $\mathbf{Alg}(\mathfrak{J}) = \mathbf{Frm}$.

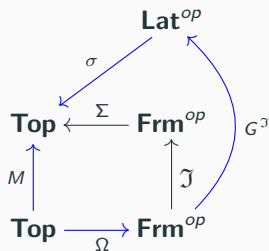
Topological spaces and frames

The monad \mathfrak{J} induces a comonad on **Frm**:



Topological spaces and frames

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Proposition

(B. Banaschewski, 1981) For a given distributive lattice D and a frame L :

$$\mathbf{Frm}(\mathfrak{J}D, L) \cong \mathbf{Lat}(D, L).$$

In particular: completely prime filters on $\mathfrak{J}D$ naturally correspond to prime filters on D .

How do we compare the algebras of M and the coalgebras of \mathfrak{J} if

$$M = \Sigma \cdot \mathfrak{J} \cdot \Omega?$$

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Proposition

(B. Banaschewski and G. C. L. Brümmer, 1988)

$\mathbf{coAlg}(\mathfrak{J}) \cong \mathbf{StKFrm}$.

Comparing algebras

Comparison functor

A monad $\mathbb{T} = (T, m, e)$ on \mathbf{C} and an adjunction $(L \dashv R, \eta, \epsilon) : \mathbf{C} \rightarrow \mathbf{B}$ induces a monad $\mathbb{M} = (M, n, d)$ on \mathbf{B} :

- $M = RTL = RG^{\mathbb{T}}F^{\mathbb{T}}L$;
- $n = RmL \cdot RT\epsilon TL$;
- $d = ReL \cdot \eta$.

There is a comparison functor $K : \mathbf{Alg}(\mathbb{T}) \rightarrow \mathbf{Alg}(\mathbb{M})$:

$$\begin{array}{ccc}
 \mathbf{Alg}(\mathbb{T}) & \xrightarrow{\quad K \quad} & \mathbf{Alg}(\mathbb{M}) \\
 \uparrow F^{\mathbb{T}} \dashv G^{\mathbb{T}} & & \uparrow F^{\mathbb{M}} \dashv G^{\mathbb{M}} \\
 \mathbf{C} & \xleftarrow{\quad L \quad} & \mathbf{B} \\
 & \xrightarrow[\quad R \quad]{\quad \perp \quad} &
 \end{array}$$

Existence of left adjoint

Theorem

If $\mathbf{Alg}(\mathbb{T})$ has coequalizers of reflexive pairs, then K admits a left adjoint.

Ref:

- P. Johnstone, Adjoint lifting theorems for categories of algebras, *Bull. London. Math. Soc.* **7** (1975), 294 -297.
- F. E. J. Linton, Coequalizers in categories of algebras, in *Seminar on Triples and Categorical Homology Theory*, Lecture Notes in Mathematics **80**, Springer Verlag, (1969), 75 -90.

Existence of left adjoint

Theorem

Suppose that \mathbf{C} is cocomplete and \mathcal{E} -well-copowered, where \mathcal{E} is part of an $(\mathcal{E}, \mathcal{M})$ -factorisation system. If T preserves \mathcal{E} , that is $Th \in \mathcal{E}$ for all $h \in \mathcal{E}$, then $\mathbf{Alg}(\mathbb{T})$ is cocomplete.

Ref:

- F. E. J. Linton, Coequalizers in categories of algebras, in *Seminar on Triples and Categorical Homology Theory*, Lecture Notes in Mathematics **80**, Springer Verlag, (1969), 75 -90.
- M. Barr, Coequalizers and free triples, *Math. Z.* **116** (1970), 307 - 322.
- J. Adámek, Colimits of algebras revisited, *Bull. Austral. Math. Soc.* **17**, (1977), 433 - 450.

Existence of left adjoint

Proposition

We have an adjunction $\mathbf{StKSp} \begin{array}{c} \xrightarrow{K^*} \\ \perp \\ \xleftarrow{K} \end{array} \mathbf{StKFrm}^{op}$.

Proof.

- **Frm** is complete.
- **Frm** is \mathcal{M} -well-powered, where $\mathcal{M} = \{f \mid f \text{ is injective}\}$.
- \mathfrak{I} preserves \mathcal{M} .



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- \mathfrak{J} preserves \mathcal{M} .

□

Proposition

If $LRT \cong T$, then:

- There are $\mathbf{Alg}(\mathbb{T}) \rightarrow \mathbf{Fix}(LR)$ and $\mathbf{Alg}(\mathbb{M}) \rightarrow \mathbf{Fix}(RL)$ that are monadic.
- K is an equivalence.

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Remark

- $\Omega\Sigma\mathfrak{J} \cong \mathfrak{J}$ if and only if BUT holds. (Banaschewski, 1983)
- If $\Omega\Sigma\mathfrak{J} \cong \mathfrak{J}$, then every stably compact frame is spatial. (Banaschewski, 1983)

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Remark

- $\Omega\Sigma\mathfrak{I} \cong \mathfrak{I}$ if and only if BUT holds. (Banaschewski, 1983)
- If $\Omega\Sigma\mathfrak{I} \cong \mathfrak{I}$, then every stably compact frame is spatial. (Banaschewski, 1983)
- The left adjoint to the (non-full) embedding $\mathbf{StKSp} \rightarrow \mathbf{Sob}$ is the restriction of Smyth's T_0 stable compactification.

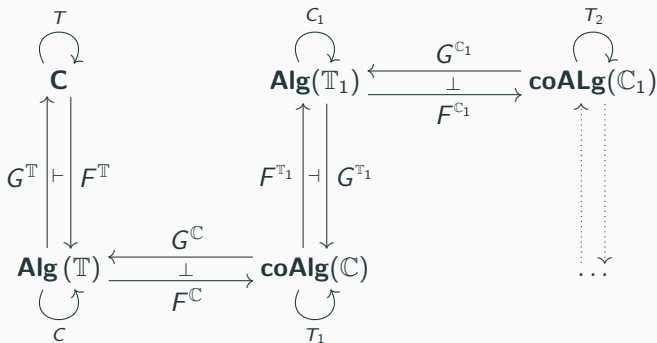
Problem

Since M induces a comonad on **StKSp**, and \mathcal{J} induces a monad on **StKFrm**, can we compare the coalgebras and algebras? Do(es) the sequence(s) stop?

Sequence of algebras

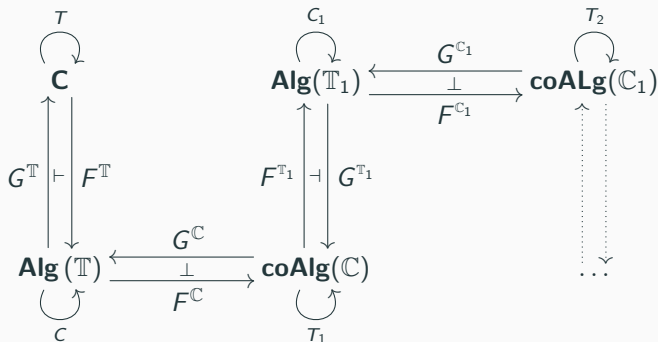
Sequence of algebras and coalgebras

A monad $T : \mathbf{C} \rightarrow \mathbf{C}$ generates the following alternating sequence:



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Theorem

(M. Barr, 1969) The sequence “stops” for general (non-constant) monads on sets, pointed sets and vector spaces over a field. In all cases, at least the comparison functor $\mathbf{C} \rightarrow \mathbf{coAlg}(C)$ is an equivalence.

Theorem

(B. Jacobs, 2013) If \mathbb{T} is a lax idempotent monad, then

$$\mathbf{Alg}(\mathbb{T}) \simeq \mathbf{Alg}(\mathbb{T}_1).$$

Sequence of algebras and coalgebras

Theorem

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$$\mathbf{Alg}(\mathbb{T}) \simeq \mathbf{Alg}(\mathbb{T}_1).$$

In the sequence $\mathbb{T}, \mathbb{C}, \mathbb{T}_1, \mathbb{C}_1, \dots$, we have

(co)Monads	(co)Algebras
$\mathbb{T} = (T, m, e)$	$a \dashv e_X$ and $a \cdot e_X = 1$
$\mathbb{C} = (T, Te, a)$	$c \dashv a$ and $a \cdot c = 1$
$\mathbb{T}_1 = (T, Ta, c)$	$b \dashv c$ and $b \cdot c = 1$
\dots	\dots

Proof.

- For a \mathbb{T}_1 -algebra $b \dashv c \dashv a \dashv e$ construct the equalizer

$$X_c \xrightarrow{\xi} X \begin{array}{c} \xrightarrow{c} \\ \xleftarrow{e} \end{array} TX$$

- X_c is a \mathbb{T} -algebra: there is $\alpha \dashv e_{X_c}$ with $\alpha \cdot e_{X_c} = 1$.
- $T(X_c) \cong X$ and $T(\alpha) \dashv T(e_{X_c}) \dashv m_{X_c} \dashv e_{TX_c}$ corresponds to $b \dashv c \dashv a \dashv e$.
- This gives a functor $\mathbf{Alg}(\mathbb{T}) \rightarrow \mathbf{Alg}(\mathbb{T}_1)$ that is an equivalence.



Proof.

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- This gives a functor $\mathbf{Alg}(\mathbb{T}) \rightarrow \mathbf{Alg}(\mathbb{T}_1)$ that is an equivalence.



Remark

The equalizer in the proof is a split equalizer, and the assignment $X \mapsto X_c$ is functorial.

Split equalizers

Split coequalizers

Given a monad $\mathbb{T} = (T, m, e)$, $\mathbf{Alg}(\mathbb{T})$ can be described as those objects that are part of a split co-equalizer:

$$\begin{array}{ccccc} TTX & \xrightleftharpoons[m_X]{Ta} & TX & \xrightarrow{a} & X \\ & \nwarrow eT & \nwarrow e & & \\ & & & & \end{array}$$

- $a \cdot m_X = a \cdot Ta$;
- $a \cdot e_X = 1$ and $m_X \cdot eT = 1$; (e and eT are “splits”.)
- $e_X \cdot a = Ta \cdot eT$.

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- $e_X \cdot a = Ta \cdot eT$.

Remark

Dual: split equalizers.

Description of algebras and coalgebras

In the sequence of monads and comonads, the algebras and coalgebras are described as follows

$$\begin{array}{ccccc} TTX & \xrightleftharpoons[Ta]{m_X} & TX & \xrightarrow{a} & X \\ & \nwarrow \quad \nearrow & \nwarrow \quad \nearrow & & \\ & eT & e & & \end{array}$$

$$\begin{array}{ccccc} (X, a) & \xrightarrow{c} & (TX, m_X) & \xrightleftharpoons[Te]{Tc} & (TTX, m_{TX}) \\ & \nwarrow \quad \nearrow & \nwarrow \quad \nearrow & & \\ & a & m_X & & \end{array}$$

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 \end{array}$$

$$\begin{array}{ccccc}
 (X, a) & \xrightarrow{c} & (TX, m_X) & \xrightleftharpoons[Tc]{Te} & (TTX, m_{TX}) \\
 & \nwarrow a & \nwarrow m_X & & \\
 & & & &
 \end{array}$$

and

$$\begin{array}{ccccc}
 (TTX, m_{TX}, Te_{TX}) & \xrightleftharpoons[Ta]{Tb} & (TX, m_X, Te_X) & \xrightarrow{b} & (X, a, c) \\
 & \nwarrow Te & \nwarrow c & & \\
 & & & &
 \end{array}$$

Proposition

B. Jacobs' construction of X_c is part of a split equalizer and a split co-equalizer:

$$\begin{array}{ccccccc} TX & \xrightleftharpoons[b]{a} & X & \xrightarrow{k} & X_c & \xrightarrow{\xi} & X & \xrightleftharpoons[c]{e} & TX \\ & \nwarrow e & & \nwarrow \xi & & \nwarrow k & & \nwarrow b & \\ & & & & & & & & \end{array}$$

Equivalence

Proposition

B. Jacobs' construction of X_c is part of a split equalizer and a split co-equalizer:

$$\begin{array}{ccccccc} TX & \xrightleftharpoons[b]{a} & X & \xrightarrow{k} & X_c & \xrightarrow{\xi} & X & \xrightleftharpoons[e]{c} & TX \\ & \searrow e & & \nwarrow \xi & & \nwarrow k & & \nwarrow b & \\ & & & & & & & & \end{array}$$

Theorem

(Modification of B. Jacobs' proof) By Beck's Theorem, X_c is a \mathbb{T} -algebra and we have $(TX_c, m_{TX_c}, T\xi, Tk) \cong (X, a, c, b)$.

Fakir construction and Stone duality

Fakir construction

Given a monad $\mathbb{T} = (T, m, e)$, construct the equalizer

$$\Phi(X) \xrightarrow{\varphi} TX \rightrightarrows_{eT}^{Te} TTX$$

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Theorem

(S. Fakir, 1970) $\Phi : \mathbf{C} \rightarrow \mathbf{C}$ underlies a monad $\mathbb{M} = (\Phi, m^\varphi, e^\varphi)$.

Natural equivalence

From the adjunction $F^{\mathbb{T}} \dashv G^{\mathbb{T}}$, consider the two functors

$$\theta^T : \mathbf{C} \rightarrow F^{\mathbb{T}}(\mathbf{C}) \text{ and } K : F^{\mathbb{T}}(\mathbf{C}) \rightarrow \mathbf{C}$$

defined by

- $\theta^T(f : X \rightarrow Y) = Tf : (TX, m_X) \rightarrow (TY, m_Y).$
- $K(Tf : (TX, m_X) \rightarrow (TY, m_Y)) = \Phi(f) : \Phi(X) \rightarrow \Phi(Y)$

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- $K(Tf : (TX, m_X) \rightarrow (TY, m_Y)) = \Phi(f) : \Phi(X) \rightarrow \Phi(Y)$

Proposition

If the monad $\mathbb{M} = (\Phi, m^\varphi, e^\varphi)$ is idempotent, then $\theta^T \cdot K \cong 1$. If Φ is the identity on \mathbf{C} , then $K \cdot \theta^T \cong 1$.

Natural equivalence

Proof.

There are two natural transformations given by $e_X^\varphi : X \rightarrow K(\theta^T(X))$ and $Te^\varphi : \theta^T(X) \rightarrow \theta^T(K(\theta^T(X)))$, with $K \cdot \theta^T = \Phi$ and $\theta^T \cdot K = F^T \cdot \Phi$.

- If \mathbb{M} is idempotent, then Te^φ is an isomorphism. (S. Fakir, 1970)
- If Φ fixes all objects in \mathbf{C} , then e_X^φ is an isomorphism.



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Theorem

If \mathbb{M} is the identity monad, then $\mathbf{C} \simeq F^\mathbb{T}(\mathbf{C})$.

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Theorem

If \mathbb{M} is the identity monad, then $\mathbf{C} \simeq F^T(\mathbf{C})$.

Theorem

Since \mathbb{M} restricts to the identity on $\mathbf{Alg}(\mathbb{T})$, then

$$\mathbf{Alg}(\mathbb{T}) \simeq F^T(\mathbf{Alg}(\mathbb{T})) \simeq \mathbf{Alg}(\mathbb{T}_1).$$

Distributive lattices and coherent frames

In the equalizer

$$E \xrightarrow{\phi} \mathfrak{I}D \underset{\downarrow \mathfrak{I}D}{\overset{\mathfrak{I}(\downarrow)}{\rightrightarrows}} \mathfrak{I}\mathfrak{I}D$$

- $\mathfrak{I}(\downarrow)(I) = \{J \mid J \ll I\}.$
- $\downarrow_{\mathfrak{I}D} (I) = \{J \mid J \subseteq I\}.$

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- $\downarrow \mathfrak{I}D(I) = \{J \mid J \subseteq I\}$.

Therefore $I \in E$ if and only if $I \ll I$, if and only if $I = \downarrow x$ for some $x \in D$.

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Theorem

DLat \simeq CohFrm.

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Theorem

DLat \simeq **CohFrm**.

Theorem

(B. Banaschewski and S. B. Niefield, 1991) The downset functor induces an equivalence between the category of meet-semilattices and that of supercoherent frames.