## Internal Factorisation Systems

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In a category  $\mathbb{C}$ , a morphism f is *orthogonal* to a morphism g (written  $f \downarrow g$ ) if for all morphisms u and v in  $\mathbb{C}$  such that vf = gu, there is a unique morphism z satisfying zf = u and gz = v. A factorisation system on a category  $\mathbb{C}$  may then be defined as a pair,  $(\mathcal{E}, \mathcal{M})$ , of classes of morphisms of  $\mathbb{C}$  which both contain all isomorphisms of  $\mathbb{C}$  and are closed under composition, such that for all e in  $\mathcal{E}$  and m in  $\mathcal{M}$ ,  $e \downarrow m$  and for all f in  $\mathbb{C}$  there exist e in  $\mathcal{E}$  and m in  $\mathcal{M}$  such that f = me.

We internalise this notion, that is, introduce internal factorisation systems for internal categories. Firstly, for an internal category, C (where  $C^{\leftarrow\leftarrow}$  is defined by the pullback on the right),

$$C_{0} \xrightarrow[]{e}{c} C_{1} \xleftarrow{m}{c} C^{\leftarrow\leftarrow} \qquad \begin{array}{c} C^{\leftarrow\leftarrow} & \xrightarrow{\pi_{2}}{c} C_{1} \\ \pi_{1} \downarrow & c \downarrow \\ C_{1} & \xrightarrow{d} C_{0} \end{array}$$

in a finitely complete category  $\mathbb{C}$ , we define the *object of points* and *object of isomorphisms* as the following pullbacks (where square brackets indicate the domain of the preceding pullback projection):

$$\begin{array}{ccc} \operatorname{Pt}(C) \xrightarrow{\pi_2} C^{\leftarrow \leftarrow} & \operatorname{Iso}(C) \xrightarrow{\pi_2} \operatorname{Pt}(C) \\ \pi_1 & \downarrow & \downarrow m & \pi_1 \\ C_0 \xrightarrow{e} C_1 & \operatorname{Pt}(C)^{\pi_2[C^{\leftarrow \leftarrow}]\pi_2[\operatorname{Pt}(C)]} C_1 \end{array} \end{array} \xrightarrow{\pi_2} \operatorname{Pt}(C)$$

and show that the two composites of projections  $\sigma = \pi_2[C^{\leftarrow\leftarrow}]\pi_2[\operatorname{Pt}(C)]\pi_1[\operatorname{Iso}(C)] : \operatorname{Iso}(C) \to C_1$  and  $\sigma' = \pi_1[C^{\leftarrow\leftarrow}]\pi_2[\operatorname{Pt}(C)]\pi_1[\operatorname{Iso}(C)] : \operatorname{Iso}(C) \to C_1$  are (equivalent) subobjects of  $C_1$ , which make the following left diagram a pullback (where  $C^{\leftarrow}$  is defined as the right pullback):

$$\begin{array}{cccc} \operatorname{Iso}(C) & \xrightarrow{\langle \langle \sigma', \sigma \rangle, \langle \sigma, \sigma' \rangle \rangle} & C^{\leftrightarrows} & & C^{\backsim} & & \\ (d,c)\sigma & & & & \\ (d,c)\sigma & & & & \\ C_0 \times C_0 & \xrightarrow{e \times e} & C_1 \times C_1 & & C^{\leftarrow \leftarrow} & \underbrace{(\pi_1[C^{\leftarrow \leftarrow}], \pi_2[C^{\leftarrow \leftarrow}])}_{C_1 \times C_1} & C_1 & \\ \end{array}$$

For two arbitrary subobjects  $\alpha : A \to C_1$  and  $\beta : B \to C_1$  of  $C_1$ , we call the following pullback the *object of composable morphisms* (of  $\alpha$  and  $\beta$ ):

$$\begin{array}{ccc} B^{\leftarrow} A^{\leftarrow} & \xrightarrow{\pi_2} & A \\ & & & \downarrow & & \\ \pi_1 & & & & c\alpha \\ & & & & & c\alpha \\ & & & & & & \\ B & \xrightarrow{d\beta} & & C_0 \end{array}$$

We write  $A^{\leftarrow} = A^{\leftarrow} A^{\leftarrow}$  if  $\alpha = \beta$  and define an *object of composable triples*, for three subobjects of  $C_1$ , similarly.

For some subobject,  $\alpha : A \to C_1$  of  $C_1$ ,  $\alpha$  is closed under composition (in C) if there exists a morphism  $m_\alpha : A^{\leftarrow \leftarrow} \to A$  making the following diagram commute:

$$\begin{array}{ccc} A & \stackrel{\leftarrow}{\leftarrow} & \stackrel{-m_{\alpha}}{- & - \rightarrow} & A \\ \alpha \times \alpha & & & \downarrow \alpha \\ C & \stackrel{\leftarrow}{\leftarrow} & \stackrel{-m}{\longrightarrow} & C_1 \end{array}$$

Then, a pair of subobjects of  $C_1$ ,  $(\varepsilon : E \to C_1, \mu : M \to C_1)$  form an *internal factorisation* system on C if  $\sigma \leq \varepsilon$  and  $\sigma \leq \mu$  (that is, if there exist  $\sigma_{\varepsilon} : \operatorname{Iso}(C) \to E$  and  $\sigma_{\mu} : \operatorname{Iso}(C) \to M$ such that  $\varepsilon \sigma_{\varepsilon} = \sigma$  and  $\mu \sigma_{\mu} = \sigma$ ),  $\varepsilon$  and  $\mu$  are both closed under composition, the following square is a pullback:

$$\begin{array}{ccc} M^{\leftarrow}C_1^{\leftarrow}E^{\leftarrow} & \xrightarrow{m(\mu \times 1) \times 1} & C_1^{\leftarrow}E^{\leftarrow} \\ \downarrow^{1 \times m(1 \times \varepsilon)} & \downarrow^{m(1 \times \varepsilon)} \\ M^{\leftarrow}C_1^{\leftarrow} & \xrightarrow{m(\mu \times 1)} & C_1 \end{array}$$

and there exists a morphism  $\tau: C_1 \to M^{\leftarrow} E^{\leftarrow}$  such that  $m(\mu \times \varepsilon)\tau = 1_{C_1}$ .

We internalise various properties of factorisation systems. Specifically, that the intersection of the classes is precisely the isomorphisms, that  $\mathcal{E}$  and  $\mathcal{M}$  respectively satisfy the right and left cancellation properties and that factorisations are unique up to isomorphism, which are respectively given by the fact that the following four squares are pullbacks:

and show that these properties are satisfied by every internal factorisation system. We then induce an order on the internal factorisation systems of an internal category and show that  $\varepsilon$ and  $\mu$  determine each other (up to equivalence of subobjects). We show that  $(\sigma, 1_{C_1})$  forms the *trivial internal factorisation system* on an internal category, which is the top element of the order. Finally, when the base category  $\mathbb{C}$  is the category **Grp**, we prove that every internal factorisation system is equivalent to the trivial one.