

Modality in worlds with different logics

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Problem

This is a joint work with Manuel António Martins - CIDMA University of Aveiro

In this presentation we would like to consider

- 1 non-trivial modal structures where worlds can operate in different logics.
- 2 the thesis that:
 - 1 (NAW) necessity means ‘true in all accessed worlds’ or the similar version.
 - 2 (PAW) possibility means ‘true in some accessed world’.
- 3 modal structures where NAW and PAW fails.
- 4 and frames that study the dynamic between logic systems.

Quine's Charge

Can we change logic and still preserve meaning?
“change of logic, change of subject” [Qui86, p. 80]

Kripke-Barcan Rigid designation

Even though Newton discovered the theory of universal gravitation, the phrase

- 1 'It is possible that Newton died in infancy'
- 2 'It is possible that the one who discovered the theory of universal gravitation died in infancy.'

Have different truth values...

Lattice semantics

Most studied logic systems have a semantics where **conjunction** and **disjunction** are **meet** and **join** in the lattice of truth values (with the appropriate designated values).

- 1 Classical: boolean algebras.
- 2 Intuitionism: Heytin algebras.
- 3 Logic of Paradox: glut lattice.

Base lattice L

Definition

A world w over L is a pair $\langle A, i \rangle$ such that

- 1 A is a complete sublattice of L and it is closed under $\neg(x)$ operation.
- 2 and i is a name for the world.

Note that L need not be complete.

Definition

A model M over L is a tuple $\langle W, R, v \rangle$ where W is a collection of worlds over L , R is a relation in $W \times W$ and v is a valuation for propositional values (i.e. $v : Var \rightarrow A_w$).

Valuation - in world

Definition

The valuation v_w of formulas in a world $w = \langle A, v \rangle$ is such that

- 1 $v_w(\neg\varphi) = -v_w(\varphi)$.
- 2 $v_w(\varphi \vee \psi) = v_w(\varphi) + v_w(\psi)$.
- 3 $v_w(\varphi \wedge \psi) = v_w(\varphi) \cdot v_w(\psi)$.
- 4 $v_w(\varphi \rightarrow \psi) = v_w(\varphi) \supset v_w(\psi)$

An important alternative for implication:

$$v_w(\varphi \rightarrow \psi) = \begin{cases} \bigcup A, & \text{if } v_w(\varphi) \leq v_w(\psi) \\ v_w(\psi), & \text{otherwise} \end{cases}$$

Valuation - modal

Definition

In a base lattice L , a value $a \in L$ is interpreted in a sublattice L' as:

- 1 **Down-interpretation** $a^{L'} = \bigcup_{L'} \{x \in L' \mid x \leq a\}$.
- 2 **Up-interpretation** $a^{L'} = \bigcap_{L'} \{x \in L' \mid x \geq a\}$.

Definition

The valuation v_w of formulas in a world $w = \langle A, v \rangle$ is such that

$$v_w(\Box\varphi) = \bigcap_A \{v_{w'}(\varphi)^A \mid wRw'\}$$

Truth in a world

Definition

For a filter F over L , we define the relation \vDash_F for a world w :

$$w \vDash_F \varphi \text{ if, and only if, } v_w(\varphi) \in F$$

Definition

For a filter F over L , we define the relation \vDash_F for a model M :

$$M \vDash_F \varphi \text{ if, and only if, } v_w(\varphi) \in F \text{ for all } w \in W_M$$

Axiom K

Theorem

Let M be a many-algebra modal model over L . Then any F -implicative world of M satisfies (under \models_F)

$$(\Box\varphi \vee \Box\psi) \rightarrow \Box(\varphi \vee \psi).$$

Theorem

Let M be a many-algebra modal model over L and F a filter. Let $w \in M$ be a F -implicative world such that A_w preserves order under $-$, and A_w has down-distribution over constants. Then w satisfies (\models_F) the axiom K, i.e. $\Box(\varphi \rightarrow \psi) \rightarrow (\Box\varphi \rightarrow \Box\psi)$.

Necessitation

Theorem

Let L be a lattice and F a filter over L , then all L -models satisfy (\models_K) necessitation if, and only if, F is complete-closed and sublattice-uniform in L .

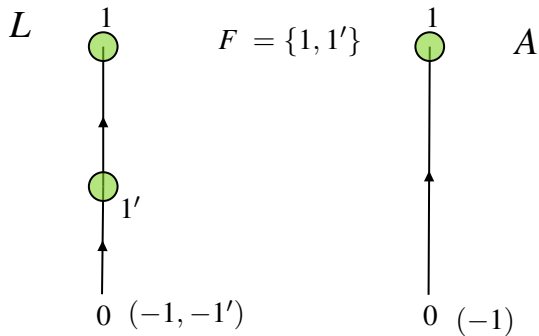
Notation

Let L be a lattice and F a filter in L , then we write L_{cu}^F for the set of all complete sublattices X of L such that F is complete-closed in X and F is X -uniform.

Theorem

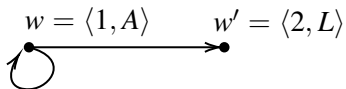
Let L be a lattice and F a filter over L , then all L -models satisfy (\models_F) L_{cu}^F -necessitation.

Consider the lattices:



Consider $w = \langle 1, A \rangle$ and $w' = \langle 2, L \rangle$ as in the figure.

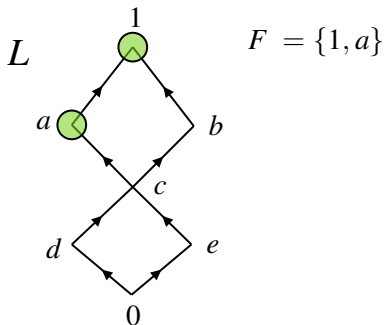
And take a valuation s such that $v_w(\alpha) = 1$ and $v_{w'}(\alpha) = 1'$ for a propositional variable α .



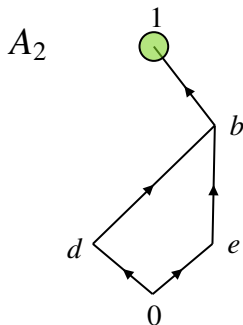
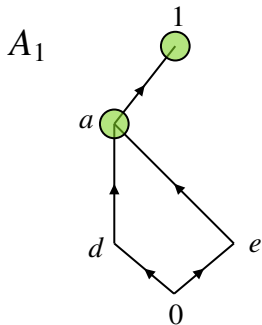
Now we observe that $w \models \alpha$, $w \not\models \neg\alpha$, $w' \models \alpha$, $w' \not\models \neg\alpha$.

And we have $w \not\models \Box\alpha$, for $v_w(\Box\alpha) = \bigcap_A \{1, 1'\} = 1 \cap 0 = 0$.

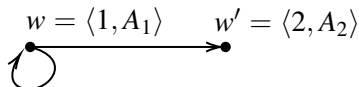
Now consider the lattice:



And the sublattices:



Consider the following frame:



And $v_w(\alpha) = 1$ and $v_{w'}(\alpha) = b$, for a propositional variable α . In this case we have $w \models \alpha$ and $w' \not\models \alpha$. Moreover, $w \models \Box\alpha$, since $v_w(\Box\alpha) = a$.

Extended notion of frame

We should now differentiate **frames** and **graphs**.

- 1 frame: including the lattice structure and world.
- 2 graph: only the structure for the accessibility relation.

We can now obtain interesting frames:

- 1 classically increasing frame: access worlds with sublattices where $a +_{w'} b \geq a +_w b$ and $a \cdot_{w'} b \leq a \cdot_w b$.
- 2 classically decreasing frame: $a +_{w'} b \leq a +_w b$ and $a \cdot_{w'} b \geq a \cdot_w b$.
- 3 dialectic frame (mix of logic and accessibility relation): two less classical worlds w_1 and w_2 accessible to w have a common accessible w' which is more classical than w .

Further discussion

- 1 up/down interpretations for negation.
- 2 recovering traditional approaches.
- 3 characterization of frames.
- 4 general properties that imply Finite Model Property.

Thank you.
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