

The Topological μ -Calculus

David Fernández-Duque

ICS of the Czech Academy of Sciences
Ghent University

TACL 2022



Topology, Algebra, and Categories in Logic

TACL 2011

MARSEILLES, FRANCE, JULY 26-30 2011

The Fifth International Conference on Topology, Algebra and Categories in Logic
is dedicated to the memory of **Leo Esakia** (1934-2010)

[TACL](#)

[Scientific Program](#)

[Registration](#)

[Venue](#)

[Contact](#)

Could not connect: Can't connect to MySQL server on 'serveurbd2' (4)

The perfect core

Theorem (Cantor-Bendixson)

Any closed subset of a Polish space is the disjoint union of a perfect set and a countable set.

The perfect core

Theorem (Cantor-Bendixson)

Any closed subset of a Polish space is the disjoint union of a perfect set and a countable set.

Theorem (General C-B)

*If X is any topological space and $A \subseteq X$, A has a maximal perfect subset, called its **perfect core**.*

The perfect core

Theorem (Cantor-Bendixson)

Any closed subset of a Polish space is the disjoint union of a perfect set and a countable set.

Theorem (General C-B)

*If X is any topological space and $A \subseteq X$, A has a maximal perfect subset, called its **perfect core**.*

The perfect core is an example of a **topological fixed point**.

Unimodal language

Modal language \mathcal{L}_\diamond :

$p \mid \neg\varphi \mid \varphi \wedge \psi \mid \Box\varphi$

Unimodal language

Modal language \mathcal{L}_\diamond :

$$p \mid \neg\varphi \mid \varphi \wedge \psi \mid \Box\varphi$$

Usual abbreviations:

▶ $\varphi \vee \psi := \neg(\neg\varphi \wedge \neg\psi)$

▶ $\Diamond\varphi := \neg\Box\neg\varphi$

Topological c -semantics of modal logic

If $X = (X, c)$ is a topological space, a topological model is a tuple $(X, \llbracket \cdot \rrbracket)$ where $\llbracket \cdot \rrbracket : \mathcal{L}_\mu^\diamond \rightarrow 2^X$ is such that

$$\llbracket \Box \varphi \rrbracket = i \llbracket \varphi \rrbracket$$

$$\llbracket \Diamond \varphi \rrbracket := c \llbracket \varphi \rrbracket$$

Topological c -semantics of modal logic

If $X = (X, c)$ is a topological space, a topological model is a tuple $(X, \llbracket \cdot \rrbracket)$ where $\llbracket \cdot \rrbracket : \mathcal{L}_\mu^\diamond \rightarrow 2^X$ is such that

$$\llbracket \Box \varphi \rrbracket = i \llbracket \varphi \rrbracket$$

$$\llbracket \Diamond \varphi \rrbracket := c \llbracket \varphi \rrbracket$$

Modal logic S4:

- ▶ $\Box(p \rightarrow q) \rightarrow (\Box p \rightarrow \Box q)$
- ▶ $\Box p \rightarrow p$
- ▶ $\Box p \rightarrow \Box \Box p$
- ▶
$$\frac{\varphi}{\Box \varphi}$$

Topological completeness of S4

If (W, \sqsubseteq) is a Kripke frame where \sqsubseteq is a preorder, the upwards-closed sets form a topology $\mathcal{U}_{\sqsubseteq}$ on W .

Topological completeness of S4

If (W, \sqsubseteq) is a Kripke frame where \sqsubseteq is a preorder, the upwards-closed sets form a topology $\mathcal{U}_{\sqsubseteq}$ on W .

- ▶ Kripke semantics for \sqsubseteq and topological semantics for $\mathcal{U}_{\sqsubseteq}$ coincide.

Topological completeness of S4

If (W, \sqsubseteq) is a Kripke frame where \sqsubseteq is a preorder, the upwards-closed sets form a topology $\mathcal{U}_{\sqsubseteq}$ on W .

- ▶ Kripke semantics for \sqsubseteq and topological semantics for $\mathcal{U}_{\sqsubseteq}$ coincide.
- ▶ Every finite topological space is of this form.

Topological completeness of S4

If (W, \sqsubseteq) is a Kripke frame where \sqsubseteq is a preorder, the upwards-closed sets form a topology $\mathcal{U}_{\sqsubseteq}$ on W .

- ▶ Kripke semantics for \sqsubseteq and topological semantics for $\mathcal{U}_{\sqsubseteq}$ coincide.
- ▶ Every finite topological space is of this form.

Theorem (McKinsey, Tarski, 1940's)

S4 is the logic of

1. *All topological spaces (with closure semantics)*

Topological completeness of S4

If (W, \sqsubseteq) is a Kripke frame where \sqsubseteq is a preorder, the upwards-closed sets form a topology $\mathcal{U}_{\sqsubseteq}$ on W .

- ▶ Kripke semantics for \sqsubseteq and topological semantics for $\mathcal{U}_{\sqsubseteq}$ coincide.
- ▶ Every finite topological space is of this form.

Theorem (McKinsey, Tarski, 1940's)

S4 is the logic of

- 1. All topological spaces (with closure semantics)*
- 2. Finite transitive reflexive frames*

Topological completeness of S4

If (W, \sqsubseteq) is a Kripke frame where \sqsubseteq is a preorder, the upwards-closed sets form a topology $\mathcal{U}_{\sqsubseteq}$ on W .

- ▶ Kripke semantics for \sqsubseteq and topological semantics for $\mathcal{U}_{\sqsubseteq}$ coincide.
- ▶ Every finite topological space is of this form.

Theorem (McKinsey, Tarski, 1940's)

S4 is the logic of

1. *All topological spaces (with closure semantics)*
2. *Finite transitive reflexive frames*
3. *Any crowded metric space (Rasiowa and Sikorski, 1960's)*

The μ -calculus

Language \mathcal{L}_μ^\diamond :

Add expressions $\mu p.\varphi(p)$ to the modal language, where p appears only **positively** in φ .

The μ -calculus

Language \mathcal{L}_μ^\diamond :

Add expressions $\mu p.\varphi(p)$ to the modal language, where p appears only **positively** in φ .

- ▶ $\llbracket \mu p.\varphi(p) \rrbracket$ is the **least fixed point** of $A \mapsto \llbracket \varphi(A) \rrbracket$.

The μ -calculus

Language \mathcal{L}_μ^\diamond :

Add expressions $\mu p.\varphi(p)$ to the modal language, where p appears only **positively** in φ .

- ▶ $\llbracket \mu p.\varphi(p) \rrbracket$ is the **least fixed point** of $A \mapsto \llbracket \varphi(A) \rrbracket$.
- ▶ $\nu p.\varphi(p) := \neg \mu p.\neg \varphi(\neg p)$ is the **greatest fixed point** of $A \mapsto \llbracket \varphi(A) \rrbracket$.

The μ -calculus

Language \mathcal{L}_μ^\diamond :

Add expressions $\mu p.\varphi(p)$ to the modal language, where p appears only **positively** in φ .

- ▶ $\llbracket \mu p.\varphi(p) \rrbracket$ is the **least fixed point** of $A \mapsto \llbracket \varphi(A) \rrbracket$.
- ▶ $\nu p.\varphi(p) := \neg \mu p.\neg \varphi(\neg p)$ is the **greatest fixed point** of $A \mapsto \llbracket \varphi(A) \rrbracket$.

Example: Transitive closure:

$$\diamond^* \varphi := \mu p.(\varphi \vee \diamond p)$$

Results for the μ -calculus

Axiomatization of μ -K4: Extend the logic K by

▶ $\varphi(\mu p.\varphi(p)) \rightarrow \mu p.\varphi(p)$

▶
$$\frac{\varphi(\theta) \rightarrow \theta}{\mu p.\varphi(p) \rightarrow \theta}$$

Results for the μ -calculus

Axiomatization of μ -K4: Extend the logic K by

▶ $\varphi(\mu p.\varphi(p)) \rightarrow \mu p.\varphi(p)$

▶
$$\frac{\varphi(\theta) \rightarrow \theta}{\mu p.\varphi(p) \rightarrow \theta}$$

Theorem (Kozen 1982)

The μ -calculus has the finite model property over the class of Kripke frames.

Results for the μ -calculus

Axiomatization of μ -K4: Extend the logic K by

$$\blacktriangleright \varphi(\mu p.\varphi(p)) \rightarrow \mu p.\varphi(p)$$

$$\blacktriangleright \frac{\varphi(\theta) \rightarrow \theta}{\mu p.\varphi(p) \rightarrow \theta}$$

Theorem (Kozen 1982)

The μ -calculus has the finite model property over the class of Kripke frames.

Theorem (Walukiewicz 1995)

The μ -calculus is sound and complete for the class of Kripke frames.

The topological μ -calculus

The μ -calculus can also be defined on topological spaces.

The topological μ -calculus

The μ -calculus can also be defined on topological spaces.

Theorem

The μ -calculus over S4 (μ -S4) is sound and complete for the class of topological spaces.

The topological μ -calculus

The μ -calculus can also be defined on topological spaces.

Theorem

The μ -calculus over S4 (μ -S4) is sound and complete for the class of topological spaces.

Proof idea.

Let φ^* be the result of replacing \diamond by \diamond^* , and apply Walukiewicz's theorem. □

The topological μ -calculus

The μ -calculus can also be defined on topological spaces.

Theorem

The μ -calculus over S4 (μ -S4) is sound and complete for the class of topological spaces.

Proof idea.

Let φ^* be the result of replacing \diamond by \diamond^* , and apply Walukiewicz's theorem. □

Theorem

The μ -calculus has the finite model property over the class of topological spaces.

The topological μ -calculus

The μ -calculus can also be defined on topological spaces.

Theorem

The μ -calculus over S4 (μ -S4) is sound and complete for the class of topological spaces.

Proof idea.

Let φ^* be the result of replacing \diamond by \diamond^* , and apply Walukiewicz's theorem. □

Theorem

The μ -calculus has the finite model property over the class of topological spaces.

Theorem (Goldblatt and Hodkinson, 2016)

μ -S4 is sound and strongly complete for the class of topological spaces.

Hierarchy collapse

Theorem (D'Agostino and Lenzi, 2010)

Every formula of the μ -calculus is equivalent to one in the alternation-free fragment over the class of transitive frames.

Hierarchy collapse

Theorem (D'Agostino and Lenzi, 2010)

Every formula of the μ -calculus is equivalent to one in the alternation-free fragment over the class of transitive frames.

Corollary

Every formula of the μ -calculus is topologically equivalent to an alternation-free one.

The tangled closure

Define

$$\diamond^\infty\{\varphi_1, \dots, \varphi_n\} := \nu p. \bigwedge \diamond(p \wedge \varphi_i).$$

The tangled closure

Define

$$\diamond^\infty\{\varphi_1, \dots, \varphi_n\} := \nu p. \bigwedge \diamond(p \wedge \varphi_i).$$

1. If (W, \sqsubseteq) is a finite S4 frame and $A_1, \dots, A_n \subseteq W$,
 $w \in \llbracket \diamond^\infty\{A_1, \dots, A_n\} \rrbracket$ iff there is a cluster $C \sqsupseteq w$ such that
for each $i \leq n$, $A_i \cap C \neq \emptyset$.

The tangled closure

Define

$$\diamond^\infty\{\varphi_1, \dots, \varphi_n\} := \nu p. \bigwedge \diamond(p \wedge \varphi_i).$$

1. If (W, \sqsubseteq) is a finite S4 frame and $A_1, \dots, A_n \subseteq W$,
 $w \in \llbracket \diamond^\infty\{A_1, \dots, A_n\} \rrbracket$ iff there is a cluster $C \sqsupseteq w$ such that
for each $i \leq n$, $A_i \cap C \neq \emptyset$.
2. If (W, \sqsubseteq) is an arbitrary S4 frame, $w \in \llbracket \diamond^\infty\{A_1, \dots, A_n\} \rrbracket$ iff
there is a path

$$w_0 \sqsubseteq w_1 \sqsubseteq w_2 \sqsubseteq \dots$$

such that $w_j \in A_i$ for infinitely many j .

The tangled closure

Define

$$\diamond^\infty\{\varphi_1, \dots, \varphi_n\} := \nu p. \bigwedge \diamond(p \wedge \varphi_i).$$

1. If (W, \sqsubseteq) is a finite S4 frame and $A_1, \dots, A_n \subseteq W$,
 $w \in \llbracket \diamond^\infty\{A_1, \dots, A_n\} \rrbracket$ iff there is a cluster $C \sqsupseteq w$ such that
for each $i \leq n$, $A_i \cap C \neq \emptyset$.

2. If (W, \sqsubseteq) is an arbitrary S4 frame, $w \in \llbracket \diamond^\infty\{A_1, \dots, A_n\} \rrbracket$ iff
there is a path

$$w_0 \sqsubseteq w_1 \sqsubseteq w_2 \sqsubseteq \dots$$

such that $w_j \in A_i$ for infinitely many j .

3. Topologically, $\llbracket \diamond^\infty\{A_1, \dots, A_n\} \rrbracket$ is the largest subspace in
which every A_i is dense.

Universality of tangle

Theorem (Dawar and Otto 2009)

Every formula of the μ -calculus is equivalent to a formula in $\mathcal{L}_{\diamond\infty}^{\diamond}$ over the class of transitive frames.

Universality of tangle

Theorem (Dawar and Otto 2009)

Every formula of the μ -calculus is equivalent to a formula in $\mathcal{L}_{\diamond\infty}^{\diamond}$ over the class of transitive frames.

Corollary

Every formula of the μ -calculus is equivalent to a formula in $\mathcal{L}_{\diamond\infty}^{\diamond}$ over the class of topological spaces.

Tangled modal logic

Define $S4^\infty$ by adding, for $\Phi = \{\varphi_1, \dots, \varphi_n\}$,

$$\diamond^\infty \Phi \rightarrow \bigwedge_i \diamond (\varphi_i \wedge \diamond^\infty \Phi) \qquad \frac{\theta \rightarrow \bigwedge_i \diamond (\varphi_i \wedge \theta)}{\theta \rightarrow \diamond^\infty \Phi}$$

Tangled modal logic

Define $S4^\infty$ by adding, for $\Phi = \{\varphi_1, \dots, \varphi_n\}$,

$$\diamond^\infty \Phi \rightarrow \bigwedge_i \diamond(\varphi_i \wedge \diamond^\infty \Phi) \qquad \frac{\theta \rightarrow \bigwedge_i \diamond(\varphi_i \wedge \theta)}{\theta \rightarrow \diamond^\infty \Phi}$$

Theorem (F-D, 2011)

$S4^\infty$ is sound and complete for

- ▶ *the class of all topological spaces*
- ▶ *the class of all finite topological spaces*

Topological d -semantics

If X is a topological space and $A \subseteq X$, define the **Cantor derivative** or **set of limit points of A** by

$$dA = \{x \in X : x \in c(A \setminus \{x\})\}.$$

Topological d -semantics

If X is a topological space and $A \subseteq X$, define the **Cantor derivative** or **set of limit points of A** by

$$dA = \{x \in X : x \in c(A \setminus \{x\})\}.$$

d -Semantics: $\llbracket \diamond \varphi \rrbracket := d \llbracket \varphi \rrbracket$.

Topological d -semantics

If X is a topological space and $A \subseteq X$, define the **Cantor derivative** or **set of limit points of A** by

$$dA = \{x \in X : x \in c(A \setminus \{x\})\}.$$

d -Semantics: $\llbracket \diamond \varphi \rrbracket := d \llbracket \varphi \rrbracket$.

Weak transitivity axiom: $\varphi \wedge \Box \varphi \rightarrow \Box \Box \varphi$.

Topological d -semantics

If X is a topological space and $A \subseteq X$, define the **Cantor derivative** or **set of limit points of A** by

$$dA = \{x \in X : x \in c(A \setminus \{x\})\}.$$

d -Semantics: $\llbracket \diamond \varphi \rrbracket := d \llbracket \varphi \rrbracket$.

Weak transitivity axiom: $\varphi \wedge \square \varphi \rightarrow \square \square \varphi$.

Topological interior: Definable by $\square \varphi := \varphi \wedge \square \varphi$.

Topological d -semantics

If X is a topological space and $A \subseteq X$, define the **Cantor derivative** or **set of limit points of A** by

$$dA = \{x \in X : x \in c(A \setminus \{x\})\}.$$

d -Semantics: $\llbracket \diamond \varphi \rrbracket := d \llbracket \varphi \rrbracket$.

Weak transitivity axiom: $\varphi \wedge \square \varphi \rightarrow \square \square \varphi$.

Topological interior: Definable by $\square \varphi := \varphi \wedge \square \varphi$.

Warning: From now on, \diamond is d , \diamond is c .

Expressivity of d -semantics

Kuratowski 1920s: The d -semantics are more expressive than the c -semantics.

Expressivity of d -semantics

Kuratowski 1920s: The d -semantics are more expressive than the c -semantics.

Let $Kur := \Box(\Box p \vee \Box \neg p) \rightarrow (\Box p \vee \Box \neg p)$.

Expressivity of d -semantics

Kuratowski 1920s: The d -semantics are more expressive than the c -semantics.

Let $Kur := \Box(\Box p \vee \Box \neg p) \rightarrow (\Box p \vee \Box \neg p)$.

Then, $\mathbb{R}^2 \models Kur$ but $\mathbb{R} \not\models Kur$.

Expressivity of d -semantics

Kuratowski 1920s: The d -semantics are more expressive than the c -semantics.

Let $Kur := \Box(\Box p \vee \Box \neg p) \rightarrow (\Box p \vee \Box \neg p)$.

Then, $\mathbb{R}^2 \models Kur$ but $\mathbb{R} \not\models Kur$.

Theorem (Shehtman, 1990)

$K4 + Kur$ is the logic of \mathbb{R}^n if $n \geq 2$.

Expressivity of d -semantics

Kuratowski 1920s: The d -semantics are more expressive than the c -semantics.

Let $Kur := \Box(\Box p \vee \Box \neg p) \rightarrow (\Box p \vee \Box \neg p)$.

Then, $\mathbb{R}^2 \models Kur$ but $\mathbb{R} \not\models Kur$.

Theorem (Shehtman, 1990)

$K4 + Kur$ is the logic of \mathbb{R}^n if $n \geq 2$.

Theorem (F-D and Iliev, 2018)

\mathcal{L}^\diamond and \mathcal{L}_μ^\diamond are both exponentially more succinct than \mathcal{L}^\diamond

T_D spaces

Fact: K4 is not sound for the class of topological spaces!

T_D spaces

Fact: K4 is not sound for the class of topological spaces!

Definition

A space X is T_D if it satisfies $ddA \subseteq dA$ for all $A \subseteq X$.

T_D spaces

Fact: K4 is not sound for the class of topological spaces!

Definition

A space X is T_D if it satisfies $ddA \subseteq dA$ for all $A \subseteq X$.

- ▶ A space is T_D iff every singleton is the intersection of a closed set and an open set.

T_D spaces

Fact: K4 is not sound for the class of topological spaces!

Definition

A space X is T_D if it satisfies $ddA \subseteq dA$ for all $A \subseteq X$.

- ▶ A space is T_D iff every singleton is the intersection of a closed set and an open set.

- ▶ $T_0 \supseteq T_D \supseteq T_1 \supseteq T_2$

T_D spaces

Fact: K4 is not sound for the class of topological spaces!

Definition

A space X is T_D if it satisfies $ddA \subseteq dA$ for all $A \subseteq X$.

- ▶ A space is T_D iff every singleton is the intersection of a closed set and an open set.
- ▶ $T_0 \supseteq T_D \supseteq T_1 \supseteq T_2$

So, \mathbb{R} , \mathbb{Q} , the Cantor space, \dots are all T_D .

K4 frames as T_D spaces

Transitive frames do not necessarily coincide with their c -semantics/ d -semantics.

K4 frames as T_D spaces

Transitive frames do not necessarily coincide with their c -semantics/ d -semantics.

If (W, \sqsubset) is **transitive** and **irreflexive** then the Cantor derivative semantics on \mathcal{U}_{\sqsubset} coincides with the Kripke semantics.

K4 frames as T_D spaces

Transitive frames do not necessarily coincide with their c -semantics/ d -semantics.

If (W, \sqsubset) is **transitive** and **irreflexive** then the Cantor derivative semantics on \mathcal{U}_{\sqsubset} coincides with the Kripke semantics.

If (W, \sqsubset) is any Kripke frame, its tree unwinding can thus be seen as a T_D space.

Completeness of K4

Theorem

K4 is the logic of

- ▶ *All transitive Kripke frames.*

Completeness of K4

Theorem

K4 is the logic of

- ▶ *All transitive Kripke frames.*
- ▶ *All **finite** transitive Kripke frames.*

Completeness of K4

Theorem

K4 is the logic of

- ▶ *All transitive Kripke frames.*
- ▶ *All **finite** transitive Kripke frames.*
- ▶ *All T_D spaces.*

Completeness of K4

Theorem

K4 is the logic of

- ▶ *All transitive Kripke frames.*
- ▶ *All **finite** transitive Kripke frames.*
- ▶ *All T_D spaces.*

No topological FMP!

Completeness of K4

Theorem

K4 is the logic of

- ▶ *All transitive Kripke frames.*
- ▶ *All **finite** transitive Kripke frames.*
- ▶ *All T_D spaces.*

No topological FMP!

Theorem

GL is the logic of finite T_D spaces:

$$\Box(\Box p \rightarrow p) \rightarrow \Box p$$

Basic results

Theorem

μ -K4 is the logic of

- ▶ All transitive frames.
- ▶ All finite, transitive frames.
- ▶ All T_D spaces.

Basic results

Theorem

μ -K4 is the logic of

- ▶ *All transitive frames.*
- ▶ *All finite, transitive frames.*
- ▶ *All T_D spaces.*

Theorem

Any μ -calculus formula is equivalent to an alternation-free formula over the class of T_D spaces.

The tangled derivative

Recall:

$$\diamond^\infty \{\varphi_1, \dots, \varphi_n\} := \nu p. \bigwedge \diamond(p \wedge \varphi_i).$$

The tangled derivative

Recall:

$$\diamond^\infty \{\varphi_1, \dots, \varphi_n\} := \nu p. \bigwedge \diamond(p \wedge \varphi_i).$$

Theorem

Every formula of the μ -calculus is equivalent to a formula in $\mathcal{L}_{\diamond^\infty}^{\diamond}$ over the class of T_D spaces with Cantor derivative.

Tangled derivative in K4 frames

A finite K4 frame (W, \sqsubset) has two types of clusters:

- ▶ reflexive clusters
- ▶ irreflexive singletons

Tangled derivative in K4 frames

A finite K4 frame (W, \sqsubset) has two types of clusters:

- ▶ reflexive clusters
- ▶ irreflexive singletons

$\diamond^\infty\{A_1, \dots, A_n\}$ holds in w iff there is a **reflexive** cluster $\mathcal{C} \supseteq w$ such that for each $i \leq n$, $A_i \cap \mathcal{C} \neq \emptyset$.

Tangled derivative in K4 frames

A finite K4 frame (W, \sqsubset) has two types of clusters:

- ▶ reflexive clusters
- ▶ irreflexive singletons

$\diamond^\infty\{A_1, \dots, A_n\}$ holds in w iff there is a **reflexive** cluster $\mathcal{C} \supseteq w$ such that for each $i \leq n$, $A_i \cap \mathcal{C} \neq \emptyset$.

Topologically, $\llbracket \diamond^\infty\{A_1, \dots, A_n\} \rrbracket$ is the largest subspace S in which every A_i is **strictly** dense:

$$S \subseteq d(S \cap A_i)$$

Tangled closure vs. derivative

Not equivalent in general:

$$\diamond^\infty \{(-\infty, 0], [0, \infty)\} \neq \diamond^\infty \{(-\infty, 0], [0, \infty)\}$$

Tangled closure vs. derivative

Not equivalent in general:

$$\diamond^\infty \{(-\infty, 0], [0, \infty)\} \neq \diamond^\infty \{(-\infty, 0], [0, \infty)\}$$

Goldblatt and Hodkinson, 2016: Over T_D spaces,

$$\diamond^\infty \Phi \equiv \bigwedge \Phi \vee \diamond \bigwedge \Phi \vee \diamond^\infty \Phi$$

A McKinsey-Tarski theorem

Theorem (Goldblatt, Hodkinson 2016)

If X is any crowded metric space, the logics

1. $K4^\infty + Kur$

A McKinsey-Tarski theorem

Theorem (Goldblatt, Hodkinson 2016)

If X is any crowded metric space, the logics

1. $K4^\infty + Kur$
2. $\mu\text{-}K4 + Kur$

are strongly complete for X .

A McKinsey-Tarski theorem

Theorem (Goldblatt, Hodkinson 2016)

If X is any crowded metric space, the logics

1. $K4^\infty + Kur$

2. $\mu\text{-}K4 + Kur$

are strongly complete for X .

Note that *Kur* need not be **sound** for X !

Arbitrary spaces

A relation $\sqsubseteq \subseteq W \times W$ is **weakly transitive** if $w \sqsubseteq v \sqsubseteq u$ implies that $w \sqsubseteq u$.

Arbitrary spaces

A relation $\sqsubseteq \subseteq W \times W$ is **weakly transitive** if $w \sqsubseteq v \sqsubseteq u$ implies that $w \sqsubseteq u$.

Theorem (Esakia, 2000's)

The logic $wK4$ is sound and complete for the class of

Arbitrary spaces

A relation $\sqsubseteq \subseteq W \times W$ is **weakly transitive** if $w \sqsubseteq v \sqsubseteq u$ implies that $w \sqsubseteq u$.

Theorem (Esakia, 2000's)

The logic $wK4$ is sound and complete for the class of

- ▶ *weakly transitive frames*

Arbitrary spaces

A relation $\sqsubseteq \subseteq W \times W$ is **weakly transitive** if $w \sqsubseteq v \sqsubseteq u$ implies that $w \sqsubseteq u$.

Theorem (Esakia, 2000's)

The logic $wK4$ is sound and complete for the class of

- ▶ *weakly transitive frames*
- ▶ *finite **irreflexive** $wK4$ frames*

Arbitrary spaces

A relation $\sqsubseteq \subseteq W \times W$ is **weakly transitive** if $w \sqsubseteq v \sqsubseteq u$ implies that $w \sqsubseteq u$.

Theorem (Esakia, 2000's)

The logic $wK4$ is sound and complete for the class of

- ▶ *weakly transitive frames*
- ▶ *finite **irreflexive** $wK4$ frames*
- ▶ *finite topological spaces*

Arbitrary spaces

A relation $\sqsubseteq \subseteq W \times W$ is **weakly transitive** if $w \sqsubseteq v \sqsubseteq u$ implies that $w \sqsubseteq u$.

Theorem (Esakia, 2000's)

The logic $wK4$ is sound and complete for the class of

- ▶ *weakly transitive frames*
- ▶ *finite **irreflexive** $wK4$ frames*
- ▶ *finite topological spaces*
- ▶ *all topological spaces*

No (more) free lunch

The weakly transitive closure is **not** μ -**definable**

No (more) free lunch

The weakly transitive closure is **not** μ -**definable**

- ▶ Completeness

No (more) free lunch

The weakly transitive closure is **not** μ -**definable**

- ▶ Completeness
- ▶ FMP

No (more) free lunch

The weakly transitive closure is **not** μ -**definable**

- ▶ Completeness
- ▶ FMP
- ▶ Alternation-elimination

No (more) free lunch

The weakly transitive closure is **not** μ -**definable**

- ▶ Completeness
- ▶ FMP
- ▶ Alternation-elimination
- ▶ Expressive completeness of tangle

No (more) free lunch

The weakly transitive closure is **not** μ -**definable**

- ▶ Completeness
- ▶ FMP
- ▶ Alternation-elimination
- ▶ Expressive completeness of tangle

do not follow from classic μ -calculus results!

The final submodel

Let $\mathcal{M}_c = (W_c, \sqsubset_c, \llbracket \cdot \rrbracket_c)$ be the canonical model for μ -wK4. This model is based on a wK4 frame.

The final submodel

Let $\mathcal{M}_c = (W_c, \sqsubset_c, \llbracket \cdot \rrbracket_c)$ be the canonical model for μ -wK4. This model is based on a wK4 frame.

But: The **truth lemma** fails for \mathcal{M}_c over the μ -calculus: it may be that $\mu p.\varphi(p) \in T$ but $T \notin \llbracket \mu p.\varphi(p) \rrbracket_c$

The final submodel

Let $\mathcal{M}_c = (W_c, \sqsubset_c, \llbracket \cdot \rrbracket_c)$ be the canonical model for μ -wK4. This model is based on a wK4 frame.

But: The **truth lemma** fails for \mathcal{M}_c over the μ -calculus: it may be that $\mu p.\varphi(p) \in T$ but $T \notin \llbracket \mu p.\varphi(p) \rrbracket_c$

Definition (Fine 1985)

Say that T is φ -**final** if $\varphi \in T$ and whenever $S \sqsupseteq T$ and $\varphi \in S$, it follows that $T \sqsupseteq S$.

The final submodel

Let $\mathcal{M}_c = (W_c, \sqsubset_c, \llbracket \cdot \rrbracket_c)$ be the canonical model for μ -wK4. This model is based on a wK4 frame.

But: The **truth lemma** fails for \mathcal{M}_c over the μ -calculus: it may be that $\mu p.\varphi(p) \in T$ but $T \notin \llbracket \mu p.\varphi(p) \rrbracket_c$

Definition (Fine 1985)

Say that T is φ -**final** if $\varphi \in T$ and whenever $S \sqsupseteq T$ and $\varphi \in S$, it follows that $T \sqsupseteq S$.

Say that T is Σ -**final** if T is φ -final for some $\varphi \in \Sigma$.

The final submodel

Let $\mathcal{M}_c = (W_c, \sqsubset_c, \llbracket \cdot \rrbracket_c)$ be the canonical model for μ -wK4. This model is based on a wK4 frame.

But: The **truth lemma** fails for \mathcal{M}_c over the μ -calculus: it may be that $\mu p.\varphi(p) \in T$ but $T \notin \llbracket \mu p.\varphi(p) \rrbracket_c$

Definition (Fine 1985)

Say that T is **φ -final** if $\varphi \in T$ and whenever $S \sqsupseteq T$ and $\varphi \in S$, it follows that $T \sqsupseteq S$.

Say that T is **Σ -final** if T is φ -final for some $\varphi \in \Sigma$.

Final submodel: $\mathcal{M}_c^\Sigma = (W_c^\Sigma, \sqsubset_c^\Sigma, \llbracket \cdot \rrbracket_c^\Sigma)$ is the submodel of Σ -final theories.

Truth lemma for the final submodel

Lemma (Σ -Final Truth Lemma)

Let Σ be finite and closed under subformulas (and a few other operations, such as single negation).

Then, for $T \in W_c^\Sigma$ and $\varphi \in \Sigma$, $T \in \llbracket \varphi \rrbracket_c^\Sigma$ iff $\varphi \in W$.

Truth lemma for the final submodel

Lemma (Σ -Final Truth Lemma)

Let Σ be finite and closed under subformulas (and a few other operations, such as single negation).

Then, for $T \in W_c^\Sigma$ and $\varphi \in \Sigma$, $T \in \llbracket \varphi \rrbracket_c^\Sigma$ iff $\varphi \in W$.

Theorem (Baltag, Bezhanishvili, F-D, 2021)

1. *The logic μ -wK4 is sound and complete for the class of wK4 frames.*

Truth lemma for the final submodel

Lemma (Σ -Final Truth Lemma)

Let Σ be finite and closed under subformulas (and a few other operations, such as single negation).

Then, for $T \in W_c^\Sigma$ and $\varphi \in \Sigma$, $T \in \llbracket \varphi \rrbracket_c^\Sigma$ iff $\varphi \in W$.

Theorem (Baltag, Bezhanishvili, F-D, 2021)

- 1. The logic μ -wK4 is sound and complete for the class of wK4 frames.*
- 2. The μ -calculus has the FMP over the class of wK4 frames.*

Cofinal subframe logics

A **cofinal subframe** of (W, \sqsubset) is a subframe based on **unbounded** $U \subseteq W$.

Cofinal subframe logics

A **cofinal subframe** of (W, \sqsubseteq) is a subframe based on **unbounded** $U \subseteq W$.

A logic is cofinal if any cofinal subframe of a Λ -frame is a Λ -frame.

Cofinal subframe logics

A **cofinal subframe** of (W, \sqsubset) is a subframe based on **unbounded** $U \subseteq W$.

A logic is cofinal if any cofinal subframe of a Λ -frame is a Λ -frame.

Theorem (Baltag, Bezhanishvili, F-D)

If Λ is a canonical, cofinal subframe extension of $wK4$, then $\mu\text{-}\Lambda$ is sound and complete for the class of finite Λ frames.

Cofinal subframe logics

A **cofinal subframe** of (W, \sqsubset) is a subframe based on **unbounded** $U \subseteq W$.

A logic is cofinal if any cofinal subframe of a Λ -frame is a Λ -frame.

Theorem (Baltag, Bezhanishvili, F-D)

If Λ is a canonical, cofinal subframe extension of $wK4$, then $\mu\text{-}\Lambda$ is sound and complete for the class of finite Λ frames.

This includes $\mu\text{-S4}$, $\mu\text{-K4}$, and many other examples.

Topological completeness

Theorem (Baltag, Bezhanishvili, F-D)

1. *The logic μ -wK4 is sound and complete for the class of topological spaces with Cantor derivative.*

Topological completeness

Theorem (Baltag, Bezhanishvili, F-D)

1. *The logic μ -wK4 is sound and complete for the class of topological spaces with Cantor derivative.*
2. *The logic μ -K4 is sound and complete for the class of T_D spaces with Cantor derivative.*

Topological completeness

Theorem (Baltag, Bezhanishvili, F-D)

1. *The logic μ -wK4 is sound and complete for the class of topological spaces with Cantor derivative.*
2. *The logic μ -K4 is sound and complete for the class of T_D spaces with Cantor derivative.*
3. *The logic μ -S4 is sound and complete for the class of T_D spaces with topological closure.*

Topological completeness

Theorem (Baltag, Bezhanishvili, F-D)

1. *The logic μ -wK4 is sound and complete for the class of topological spaces with Cantor derivative.*
2. *The logic μ -K4 is sound and complete for the class of T_D spaces with Cantor derivative.*
3. *The logic μ -S4 is sound and complete for the class of T_D spaces with topological closure.*
4. *The logic μ -wK4 T_0 (which I won't define here) is sound and complete for the class of T_0 spaces with topological derivative.*

Alternation elimination

Theorem (Pacheco and Tanaka, 2022)

Every formula of the μ -calculus is equivalent to an alternation-free formula over the class of topological spaces and over the class of wK4-frames.

Alternation elimination

Theorem (Pacheco and Tanaka, 2022)

Every formula of the μ -calculus is equivalent to an alternation-free formula over the class of topological spaces and over the class of wK4-frames.

D'Agostino and Lenzi's result does not apply, but the **proof** does (with some care).

Alternation elimination

Theorem (Pacheco and Tanaka, 2022)

Every formula of the μ -calculus is equivalent to an alternation-free formula over the class of topological spaces and over the class of wK4-frames.

D'Agostino and Lenzi's result does not apply, but the **proof** does (with some care).

Question: Is the tangled fragment expressively complete?

Expressive incompleteness

- ▶ Tangled closure: \diamond^∞
- ▶ Tangled derivative: \diamond^∞

Expressive incompleteness

- ▶ Tangled closure: \diamond^∞
- ▶ Tangled derivative: \diamond^∞

Theorem (F-D, Gougeon)

1. \diamond^∞ is not definable in $\mathcal{L}_{\diamond, \diamond^\infty}$ over the class of T_D spaces

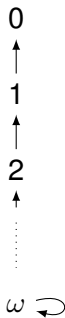
Expressive incompleteness

- ▶ Tangled closure: \diamond^∞
- ▶ Tangled derivative: \diamond^∞

Theorem (F-D, Gougeon)

1. \diamond^∞ is not definable in $\mathcal{L}_{\diamond\diamond^\infty}$ over the class of T_D spaces
2. \diamond^∞ is not definable in $\mathcal{L}_{\diamond\diamond^\infty}$ over the class of T_0 spaces

Undefinability of $\diamond^\infty\{\top\}$ in $\mathcal{L}_{\diamond^\infty}$



The hybrid tangle

Given $\Phi = (\varphi_1, \dots, \varphi_n)$,

$$\blacklozenge^\infty \Phi := \nu p. \bigvee_{i \leq n} \left(\blacklozenge (\varphi_i \wedge \blacklozenge^\infty \Phi) \wedge \bigwedge_{j \neq i} \blacklozenge (\varphi_j \wedge \blacklozenge^\infty \Phi) \right)$$

The hybrid tangle

Given $\Phi = (\varphi_1, \dots, \varphi_n)$,

$$\blacklozenge^\infty \Phi := \nu p. \bigvee_{i \leq n} \left(\blacklozenge (\varphi_i \wedge \blacklozenge^\infty \Phi) \wedge \bigwedge_{j \neq i} \blacklozenge (\varphi_j \wedge \blacklozenge^\infty \Phi) \right)$$

Theorem (F-D, Gougeon)

\blacklozenge^∞ and \blacklozenge^∞ are definable in $\mathcal{L}_{\blacklozenge \blacklozenge^\infty \blacklozenge^\infty}$ over the class of topological spaces.

The hybrid tangle

Given $\Phi = (\varphi_1, \dots, \varphi_n)$,

$$\blacklozenge^\infty \Phi := \nu p. \bigvee_{i \leq n} \left(\lozenge (\varphi_i \wedge \blacklozenge^\infty \Phi) \wedge \bigwedge_{j \neq i} \lozenge (\varphi_j \wedge \blacklozenge^\infty \Phi) \right)$$

Theorem (F-D, Gougeon)

\lozenge^∞ and \blacklozenge^∞ are definable in $\mathcal{L}_{\lozenge \blacklozenge^\infty \lozenge^\infty}$ over the class of topological spaces.

$$\lozenge^\infty \{\varphi_1, \dots, \varphi_n\} := \blacklozenge^\infty (\varphi_1, \varphi_1, \dots, \varphi_n, \varphi_n)$$

The hybrid tangle

Given $\Phi = (\varphi_1, \dots, \varphi_n)$,

$$\blacklozenge^\infty \Phi := \nu p. \bigvee_{i \leq n} \left(\lozenge (\varphi_i \wedge \blacklozenge^\infty \Phi) \wedge \bigwedge_{j \neq i} \lozenge (\varphi_j \wedge \blacklozenge^\infty \Phi) \right)$$

Theorem (F-D, Gougeon)

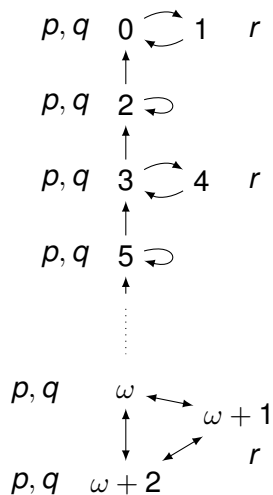
\lozenge^∞ and \blacklozenge^∞ are definable in $\mathcal{L}_{\lozenge \blacklozenge^\infty \lozenge^\infty}$ over the class of topological spaces.

$$\lozenge^\infty \{\varphi_1, \dots, \varphi_n\} := \blacklozenge^\infty (\varphi_1, \varphi_1, \dots, \varphi_n, \varphi_n)$$

Theorem (F-D, Gougeon)

\blacklozenge^∞ is not definable in $\mathcal{L}_{\lozenge \blacklozenge^\infty \lozenge^\infty}$ over the class of topological spaces.

Undefinability of $\diamond^\infty(p, q, r)$ in $\mathcal{L}_{\diamond^\infty}^{\diamond^\infty}$



Concluding remarks

- ▶ Much of the classic work on the μ -calculus largely carries over to its topological variant, but breaks down when dropping the T_D assumption.

Concluding remarks

- ▶ Much of the classic work on the μ -calculus largely carries over to its topological variant, but breaks down when dropping the T_D assumption.
- ▶ However, proofs from scratch are simpler than their Kripke analogues.

Concluding remarks

- ▶ Much of the classic work on the μ -calculus largely carries over to its topological variant, but breaks down when dropping the T_D assumption.
- ▶ However, proofs from scratch are simpler than their Kripke analogues.
- ▶ Well-studied tangled fragments are no longer expressively complete over arbitrary spaces.

Concluding remarks

- ▶ Much of the classic work on the μ -calculus largely carries over to its topological variant, but breaks down when dropping the T_D assumption.
- ▶ However, proofs from scratch are simpler than their Kripke analogues.
- ▶ Well-studied tangled fragments are no longer expressively complete over arbitrary spaces.

Question: Is there also a simple, expressively complete fragment for all topological spaces?

Concluding remarks

- ▶ Much of the classic work on the μ -calculus largely carries over to its topological variant, but breaks down when dropping the T_D assumption.
- ▶ However, proofs from scratch are simpler than their Kripke analogues.
- ▶ Well-studied tangled fragments are no longer expressively complete over arbitrary spaces.

Question: Is there also a simple, expressively complete fragment for all topological spaces?

- ▶ Open problems abound! (Connectedness, polytopological μ -calculus, definable classes. . .)

Concluding remarks

- ▶ Much of the classic work on the μ -calculus largely carries over to its topological variant, but breaks down when dropping the T_D assumption.
- ▶ However, proofs from scratch are simpler than their Kripke analogues.
- ▶ Well-studied tangled fragments are no longer expressively complete over arbitrary spaces.

Question: Is there also a simple, expressively complete fragment for all topological spaces?

- ▶ Open problems abound! (Connectedness, polytopological μ -calculus, definable classes. . .)

Obrigado!